

# INTERMEDIATE PROGRAMMING FOR THE TRS-80® (MODEL I)

David L. Heiserman

FD C6  
FD C9  
FD CC  
FD CD  
FD D0  
FD D2  
FD D4  
FD D7  
FD D8  
FD D9  
FD DA  
FD DB  
FD DC  
FD DD  
FD DE  
FD DF  
FD E0  
FD E1  
FD E2  
FD E3  
FD E4  
FD E5  
FD E6  
FD E7  
FD E8  
FD E9  
FD EA  
FD EB  
FD EC  
FD ED  
FD EE  
FD EF  
FD F0  
FD F1  
FD F2  
FD F3  
FD F4  
FD F5  
FD F6  
FD F7  
FD F8  
FD F9  
FD FA  
FD FB  
FD FC  
FD FD  
FD FE  
FD FF

21 E9 FD  
22 18 40  
AF  
21 C5 FD  
ED 52  
28 0F  
2A 11 44  
11 C0 FD  
52



LD HL, 0FDE9H  
LD HL, (4015H), HL  
LD A  
LD HL, 0FDC5H  
SBC HL, DE  
JR Z, \$+11H ; FDE3H  
LD HL, (4411H)  
LD DE, 0FCC0H  
XOR A  
SBC HL, C  
RET



LD HL, 0FDE9H  
LD HL, (4015H), HL  
LD A  
LD HL, 0FDC5H  
LD DE  
LD HL, \$+11H ; FDE3H  
LD HL, (4411H)  
LD DE, 0FCC0H  
LD DE  
LD HL



# **Intermediate Programming for the TRS-80<sup>®</sup> (Model I)**

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# **Intermediate Programming for the TRS-80<sup>®</sup> (Model I)**

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## Preface

I don't think there is such an animal as a perfectly contented computer programmer and operator. Such people are chronically afflicted with a desire to acquire more equipment or find ways to use existing schemes more effectively.

The matter of acquiring more equipment—more memory, more peripherals, and so on—can be an expensive route to doing more and better things with a computer system. In fact, it is often a prohibitively expensive route.

A much less expensive alternative is to learn how to use a system's existing features more effectively. It is quite surprising in many instances to discover that your own system has more computing power than you imagined. Tucked away in those ROMs and available utility tapes are a lot of features just crying to be used. Unfortunately, those bits and pieces of computing power are often overlooked, by both the user and the manufacturer who wrote the instruction manual.

This book helps you uncover some of those “hidden” features of the TRS-80 personal computer system and then shows you how to exploit them to your programming advantage. If you already have a standard Model I TRS-80 with 16K of RAM and Level II capability, your only additional expense will be tied up in two Radio Shack cassette tapes: T-BUG and Editor/Assembler. A line printer can be very helpful in the later stages of this exploration of the TRS-80, but it isn't absolutely essential.

The book is designed to stimulate your own sense of programming creativity. It shows how something can be done in a general way, illustrating the point with specific examples. It is then up to you to use the information to work out programs of your own invention.

DAVID L. HEISERMAN

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## CHAPTER 1

# You, Your TRS-80, and This Book

Do you remember the first day you fired up your brand-new TRS-80? You probably do. For most of us that was an exciting and rewarding experience.

One of the nice things about owning your own home computer is that the feeling of having first-time adventures doesn't have to wear off; there is always something new to learn and try. Learning something new, trying it, and making it work can be just as much fun as turning on the computer the first time.

Oh sure, there are times when things don't go right and you feel like throwing the whole system across the room. Things can go wrong, they don't always work out as expected, and the frustration level can grow to disheartening proportions. But that happens to everyone who works with computers and computer programming, no matter how much or how little experience they have and no matter how sophisticated or modest the computer system is.

Home computer programming, however, still retains all the potential for being a continuously rewarding experience. All you have to do is learn what you need to know as you go along, applying the newfound knowledge until it becomes second nature to you. Then you are ready to learn something else. There is really no end to it. And it's great fun.

The key to maintaining an ongoing love affair with your computer is thus a matter of learning to do new things or to do old things in a new and better way. Doing the same old things in the same old fashion can become boring or tedious, no matter how well they work out. But in the process of trying new things, there is always the risk of failure, disappointment, and frustration.

This sort of frustration and failure most often arises from ig-

norance—having an unclear grasp of certain operating details, misunderstanding some of them, or, perhaps, being unaware of them altogether. The more you know about the workings of your computer, the more likely are your chances of success with new ideas.

The primary objective of this book is to help you get more fun out of creating computer programs on your TRS-80. The idea is to make it possible for you to engage in that unending and rewarding adventure mentioned in the opening paragraphs of this chapter.

How does this book help you? It describes some powerful operating details that are usually mentioned too briefly or overlooked altogether in the current Radio Shack literature. These are the sort of operating details that make it possible for you to do new things and to do old things in new and better ways. The descriptions give you information that provides a logical next step in your computer programming experiences.

There are plenty of examples and demonstrations to illustrate the operating details under discussion, but it will be up to you to transfer the ideas to your own programming plans. In a manner of speaking, this book points the way to new programming possibilities, but you get all the fun of seeing the ideas working within the context of your programs.

## **WHICH SYSTEM DO YOU NEED?**

TRS-80s are now available in such a wide variety of configurations that it is impractical to attempt writing a book that suits all of them equally well. It is thus necessary to draw some lines, meeting the needs of the largest number of readers and hoping others will find information that is useful to them and applicable to their system configurations.

The examples and demonstrations in this book have been worked around a Model I TRS-80 having 16K of RAM and Level II BASIC. All information provided here applies to that system configuration.

Now, that doesn't mean every other system is completely counted out. Indeed, things are going to be a bit tough for anyone working with a 4K system or one equipped with Level I BASIC. Readers in this category ought to give serious thought to upgrading their systems—not only to get the most from this book but also to get a lot more computing opportunities.

Readers with bigger systems will find that virtually all the material in this book applies to their machines. Some of the numbers will be different for Model II and Model III machines, and you will have to work out the conversions for yourself. The spirit of the presentations, however, is almost universally applicable.

In keeping with the notion of working toward the most popular

TRS-80 configuration, you will need a cassette tape machine. DOS users will be able to work all the demonstrations, too, but there is no DOS-related material presented in this book.

A line printer can be quite helpful in some of the discussions, but it is never a critical requirement. Also, an expansion interface can be used, but it isn't necessary.

Looking ahead to work in the latter part of the book, you will need a couple of Radio Shack program tapes: T-BUG (catalog No. 26-2001) and Editor/Assembler (catalog No. 26-2002).

You will also need a Z-80 programming reference book. The Z-80 instruction set is offered in Appendix B of this book, but there is no room here for spelling out the details of machine-language programming in a general sense. If you need such a book, try Radio Shack's *How to Program the Z-80* (catalog No. 62-2066).

## HOW THIS BOOK IS ORGANIZED

When most people are introduced to their first home computer, they get a lot of delight from running BASIC programs listed in the programming manual and, maybe, running a few "canned" programs—usually game programs.

But we have to be honest here: running simple BASIC programs and prepared cassette tapes can wear thin after a while. Of course you can buy more elaborate and expensive prepared programs, or start entering some programs from published BASIC or machine-language listings; but even if they meet your expectations (and many of them won't), they also become old hat after a while.

One way to overcome this sagging enthusiasm for your computer is to begin writing custom programs—cutting and trying a few of your own invention. That can be a lot of fun, especially if you know what you're doing. Learning to use conventional BASIC can keep you going for a long time.

Indeed, writing custom programs in BASIC can serve a lot of personal needs; however, it usually doesn't take long to become dissatisfied with certain limitations of it. Experienced TRS-80 BASIC programmers often begin feeling straightjacketed by some of the built-in procedures. It is quite possible to know exactly what you want to do but find that BASIC—whether in the TRS-80 or not—cannot handle the job as effectively or adequately as you'd like.

Animated graphics, for instance, can fall flat in BASIC because of the long execution times for the statements involved. Or maybe you would like to assign something more than 240 characters to a single string variable. In such instances the built-in relationships between BASIC and the operating system stand in your way.

There are untold instances where the structure of the system

and the way BASIC is meant to work serve as roadblocks to effective programming.

One way to tackle this particular sort of problem is by digging through the avalanche of books and magazines written for people who want to go beyond the limitations of their present know-how. If you have tried that route, you have probably been disappointed more than once. It isn't that there is anything necessarily wrong with all that available information; it can prove quite valuable in many ways.

But most of the literature dealing with new TRS-80 tricks and techniques applies only to the particular situation the author is describing. And what's more, you always run the risk of applying the idea without learning anything you can use on your own at some later time.

More often than not, the real value of a book or article lies in the gems of wisdom tucked away in the program listings or accompanying text. Specific solutions for specific problems may or may not be truly helpful, but the methods and ideas behind them can be invaluable. Therein lies the virtue of most computer books and magazine articles.

For example, an article describing how to move a spot of light across the screen by depressing a certain key might not seem all that useful or exciting to you; but the technique for sensing the key depression or drawing the moving spot can be applied in countless ways—once you grasp the main principles behind those actions.

This is a book about main principles. You won't have to dig through the program listings to uncover important ideas; they are clearly spelled out.

Yes, indeed, there are a lot of program listings in this book, but they are intended only to illustrate the workings of the principle at hand. The programs, as such, aren't all that useful or exciting; they are to-the-point illustrations and not highly refined and fully developed programs. The programs are trimmed to the bare bones so that the point they illustrate will stand out as clearly as possible. Finished programs tend to be cluttered with a lot of "whistles and bells" that obscure the finer, more important details.

It is up to you to grasp the essence of an idea presented in this book, give it a first try with an accompanying program listing, and then fit it into some programming schemes of your own.

The author hopes that you will find this book rich in ideas and short on razzle-dazzle.

Finally, you should be assured at this point that the book is not devoted exclusively to machine-language programming. Many people who feel the itch to go beyond basic BASIC are told—or at

least get the impression—that the next step in their programming experience must be in the direction of machine-language programming.

That is not true. Growing up in the business of computer programming is an evolutionary process. Your knowledge ought to develop gradually and smoothly, and, for most people, moving directly from basic BASIC to pure machine-language programming is hardly a gradual and smooth process. In fact, it is a terrible mistake to drop BASIC and move fully to machine-language programming if you've had no training or experience with it. The change is too big and too abrupt.

No, machine-language programming is not the first topic offered in this book. It turns out that familiar old BASIC can become exciting again, once you know more about the internal workings of the TRS-80. You can access some very useful inner workings from BASIC and do a lot of things that will tear down some of the usual programming limitations.

And that's where this book starts.

Once you know more about the system from a BASIC viewpoint, you are ready to begin some machine-language programming. But even then it is possible to work into that sort of programming from BASIC. The idea is to let you wade into the deeper waters of machine language, while keeping a tight grasp on a familiar BASIC handle.

Toward the last part of the book, we will finally get to purely machine-language programming. By that time, though, you will be well grounded in the TRS-80's internal workings and better prepared to write successful machine-language programs of your own.

So this book is for you if you want to do more with BASIC, learn more about the internal workings of the TRS-80, ease your way gradually into machine-language programming, and learn to devise new and exciting programs of your own.

In short, if you are getting a little tired of your TRS-80 system, this book ought to serve as a shot of adrenalin.

## **HOW TO GET THE MOST FROM THIS BOOK**

This is not really a reference book, although it might have that general appearance. The book represents a step-by-step process, and, as such, you will get more from it by working through it from beginning to end, as opposed to dipping in at some point that sounds interesting to you.

Besides working through the material in sequence and from the beginning, it is also helpful to study the material with your TRS-80 system right at hand. A program listing accompanies every new idea;

so as you read about a new idea, you should be able to try it for yourself, right on the spot. You probably already know that the learning-by-doing process is the most effective one for working with computer programs.

But as mentioned earlier, the program listings are meant only to illustrate—to emphasize—one or two key points. So not many of the programs presented here are, in themselves, very useful. The idea is to encourage you to try the new ideas in a bare-bones fashion. Once you understand a new idea and have an opportunity to see it work, you should try fitting it into a more elaborate or useful program of your own design. If things don't work out the way you think they should, go back to the book and see whether or not you have overlooked or misunderstood something.

This book is a guide—a self-teaching guide. You will get the most from it by attempting to apply the new ideas in your fashion. It is really based on the old notion that you can keep a man alive for a day if you give him a fish, but you can keep him alive for a good many years if you give him some fishing equipment and show him how to use it.

You have all the equipment, and here come some ideas about how to use it in more effective ways.

## CHAPTER 2

# The Video Environment

The crt screen is the most-often used output device for most TRS-80 systems, and it turns out that this video feature has a level of versatility that matches up fairly well with its importance. Understanding how the video system is set up and knowing how to use it can go a long way toward building some fine video-oriented programs.

The purpose of this chapter is to explore the main features of the TRS-80 video system in detail, expanding on the points that are merely suggested or implied in the *Level II BASIC Reference Manual*. Some of the video features are better implemented in BASIC, while others are more suited to machine-language treatments. The format used throughout this chapter is BASIC oriented; later chapters illustrate some of the same principles in machine language.

A few of the principles are of questionable value, appearing to be rather awkward in BASIC and machine language. These principles are described here for the sake of exploring some unusual ideas. Even these points of questionable value will at least serve to build your understanding of the TRS-80 video system and its potential applications.

## VIDEO DATA

The information flowing around inside a computer consists of two parts: addresses and data. The address portion of that information specifies *where* something is to be placed, and the data portion specifies *what* is to be placed there.

In the context of the TRS-80 video system the addressing is usually



related to some position on the crt screen, and the data concerns what is to be done at that point. This section deals with the video data. The next section of this chapter introduces some details of the video memory addressing.

The video data is carried as a 1-byte (8-bit) binary code, which means there are 256 possible combinations of things that can be done at a given point on the crt screen.

In the TRS-80 system those 256 video-related data codes can be divided into four categories:

- control codes—data codes 0 through 31
- alphanumeric character (ASCII) codes—32 through 127
- TRS-80 graphic codes—128 through 191
- TRS-80 tab codes—192 through 255

This list accounts for all 256 video data codes. Some of the functions are redundant (do the same thing as others), and a few of the codes do nothing at all. But, even so, there is a lot that can be done with the video data set.

It would be nice if we could investigate the video data set in a systematic fashion, beginning with control code 0 and progressing, one step at a time, through TRS-80 tab code 255. Unfortunately, the control codes and tab codes do not lend themselves to very convincing demonstrations at this point in the discussion.

So, for the sake of avoiding undue confusion in the early going, those control and tab codes will be explained later in this chapter. Of immediate interest, then, are the alphanumeric character codes and the TRS-80 graphic codes.

### **Alphanumeric Character Codes**

The entire TRS-80 alphanumeric character set is represented by data codes 32 through 127. For the most part this code follows the well-established convention developed as the American Standard Code for Information Interchange (ASCII). The exceptions are the lack of special line, brace, and bracket symbols that appear on ASCII code charts but not on the TRS-80 keyboard.

Table 2-1 summarizes the TRS-80 alphanumeric character set, and if you want to see the characters on your own screen, try this short program:

```
10 REM ** ALPHANUMERIC CHARACTER SET DEMO **
20 CLS
30 FOR C=32 TO 127
40 PRINT CHR$(C);
50 NEXT C
60 END
```

This program prints the entire alphanumeric character set (charac-

**Table 2-1. TRS-80 ASCII Character Set**

Code (Decimal)	Character	Code (Decimal)	Character
32	space	80	P
33	!	81	Q
34	"	82	R
35	#	83	S
36	\$	84	T
37	%	85	U
38	&	86	V
39	'	87	W
40	(	88	X
41	)	89	Y
42	*	90	Z
43	+	91	↑
44	,	92	↓
45	-	93	←
46	.	94	→
47	/	95	—
48	0	96	@
49	1	97	a
50	2	98	b
51	3	99	c
52	4	100	d
53	5	101	e
54	6	102	f
55	7	103	g
56	8	104	h
57	9	105	i
58	:	106	j
59	;	107	k
60	<	108	l
61	=	109	m
62	>	110	n
63	?	111	o
64	@	112	p
65	A	113	q
66	B	114	r
67	C	115	s
68	D	116	t
69	E	117	u
70	F	118	v
71	G	119	w
72	H	120	x
73	I	121	y
74	J	122	z
75	K	123	↑
76	L	124	↓
77	M	125	←
78	N	126	→
79	O	127	—

ter codes 32 through 127) in about 1½ lines across the top of the screen.

If you do not have the lowercase modification for your system, you will see the alphabetical characters (A-Z), four arrows, and a dash printed at two places in the display. The standard system cannot distinguish uppercase and lowercase characters, and codes 64 through 95 do the same thing as codes 96 through 127. With the lowercase modification, the alphabet appears in a lowercase format the second time through. So there is some redundancy here, unless you happen to have the lowercase character generator installed.

The CHR\$(n) statement gives you access to the TRS-80 character generator in BASIC and with decimal code numbers. If you would rather tinker around with the character generator on a one-at-a-time basis, try this program:

```
10 REM ** CALL A CHARACTER **
20 CLS
30 INPUT "WHAT ASCII CHARACTER CODE (32-127)";C
40 IF C<32 OR C>127 THEN 30
50 PRINT CHR$(C)
60 PRINT
70 GOTO 30
```

Line 30 in this program requests an ASCII character code in the range 32 through 127. If you happen to respond with a number outside that range, line 40 returns the program to line 30, thus preventing you from tinkering with character or control codes outside the range of immediate interest here.

If you are still in the mood for playing around with this character set, it can be instructive to reverse the operation: writing a program that lets you specify an alphanumeric or any other ASCII character by typing it on the keyboard, and letting the computer print the character code, itself.

```
10 REM ** CALL-A-CODE **
20 CLS
30 C$=INKEY$:IF C$="" THEN 30
40 C=ASC(C$)
50 IF C>=32 AND C<=127 THEN 70
60 PRINT "THAT IS A CONTROL CHARACTER":GOTO 30
70 PRINT C$;"  ";C
80 GOTO 30
```

On running CALL-A-CODE, you can type away on the keyboard, and with every valid character entry you will see the character linked with its ASCII code number. Since some of the keys generate control codes, line 50 is necessary to avoid some FC ERROR messages that stop the program. Whenever you strike one of those

control keys, this program responds by printing THAT IS A CONTROL CHARACTER and letting you try another key.

While you are studying the alphanumeric character set with the CALL-A-CODE program, be sure to try it while the SHIFT key is depressed—that will give you the uppercase version of some of the special symbols and the lowercase codes for the alphabetical characters. Striking the A key with the SHIFT key *not* depressed, the program shows the letter A linked to character code 65. Striking that same key while the SHIFT key *is* depressed prints an A, followed by character code 97.

Also try the arrow keys with the CALL-A-CODE routine. This gives you code numbers for arrow characters that can be useful for writing graphics programs having arrow figures in them.

### TRS-80 Graphic Codes

Video data codes 128 through 191 hold graphic patterns that are unique to the TRS-80. There are 64 of them in that family, and they are all different from one another. See Fig. 2-1.

The rationale behind the structure of these figures is not really very obvious until they are described in terms of the binary versions of their code numbers—a task that will be handled later. But with a table of these figures available, you don't have to be concerned about any discernible relationship between the code numbers and the graphic figures they call to the screen. It is entirely possible to build up some elaborate graphics from the decimal versions of the graphic-code numbers.

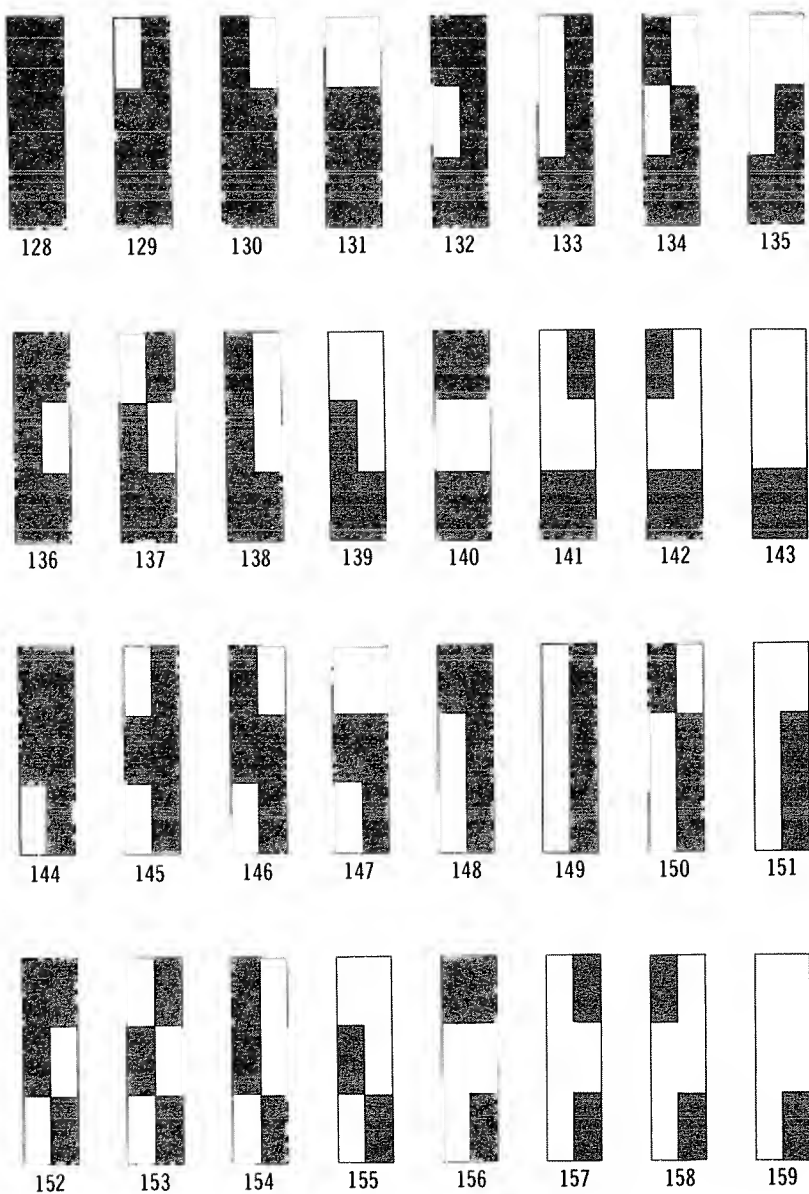
The graphic character set can be called to the screen just as the ASCII characters are, using PRINT CHR\$(*n*) statements, where *n* is the character code number between 128 and 191, inclusively.

To illustrate the point, try this program:

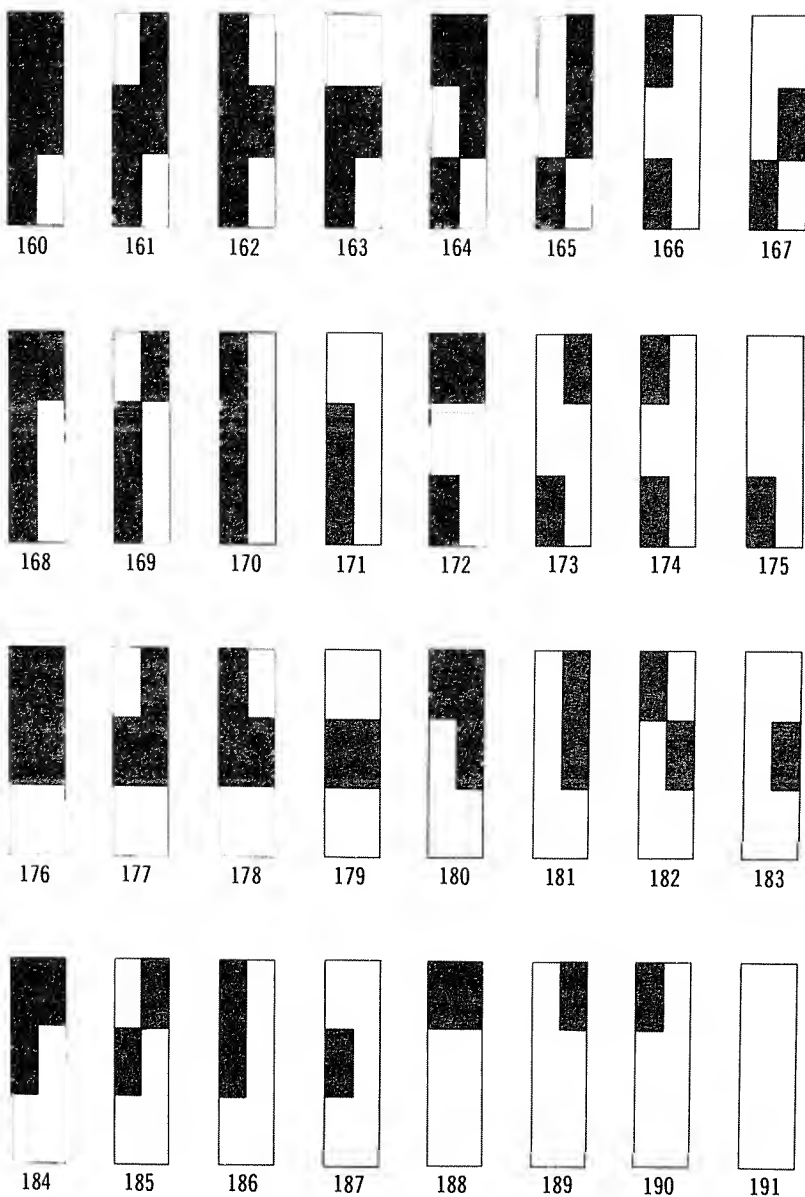
```
10 REM ** GRAPHIC CHARACTER SET DEMO **
20 CLS
30 FOR C=128 TO 191
40 PRINT CHR$(C);CHR$(128);
50 NEXT C
60 END
```

This program prints out the entire graphic character set, separating each character with a space that is called by the CHR\$(128) statement in line 40. Thus the graphics can be called to the screen just as the ASCII characters are, and that should come as no surprise, considering the graphics are generated by the same sort of read-only-memory (ROM) device.

If you want to link the graphic characters with their decimal code numbers in a clearer fashion, try this selection program:



**Fig. 2-1. The family of TRS-80 graphic characters**



and graphic codes (decimal).

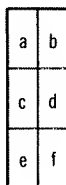
```

10 REM ** PICK A GRAPHIC **
20 CLS
30 INPUT "SPECIFY A GRAPHIC CODE (128-191)";C
40 IF C>=128 AND C<=191 THEN 60
50 PRINT "NOT A GRAPHIC";GOTO 30
60 PRINT CHR$(C)
70 PRINT;GOTO 30

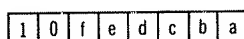
```

Now you can specify and view the graphic characters one at a time.

To get a better appreciation for what is going on here, compare Fig. 2-1 and Fig. 2-2. For graphic purposes each character space on the screen is divided into a set of six smaller rectangles, labeled *a* through *f* in Fig. 2-2. A given rectangle, or *pixel*, within the character space is either lighted or darkened, depending on the code number being used.



(A) Geometry of the graphic character space.



- A 0 IN A GRAPHIC CODE BYTE DARKENS ITS PLACE IN CHARACTER SPACE.
- A 1 IN A GRAPHIC CODE BYTE LIGHTS ITS PLACE IN THE CHARACTER SPACE.

(B) Graphic code byte.

**Fig. 2-2. Graphic character format and character-code layout.**

Graphic code 129, for example, lights the little rectangle in the upper left-hand corner of the character space. In Fig. 2-2 that means pixel *a* is lighted by a binary 1 in the *a*-bit space of the character code. Code 130, on the other hand, lights the pixel in the upper right-hand corner of the current character space, and code 131 lights both of the top pixels (both *a* and *b*).

This process goes on through the entire 64-character graphic set, incrementing bits *a* through *f* in a binary counting sequence.

As mentioned previously, it is possible to use CHR\$ statements to call the graphic characters, and thus compose some fairly elaborate graphic images on the crt. Of course, you have to know more about the video memory in order to position the individual graphic characters on the screen, and that is the next topic of discussion.

If you plan to compose any special graphic images, you will do well to pick up a pad of Radio Shack's Video/Programming Worksheets (catalog No. 26-2105). Those worksheets divide the screen addresses into character spaces and graphic pixels for you. All you have to do is sketch your graphics on the worksheet, then use

Fig. 2-1 to select the graphic character code required for each character space in the figure.

## VIDEO ADDRESSING

If you consult a memory map for the TRS-80 system, you will find that the video memory occupies addresses 15360 through 16383, inclusively. That means 1024 different memory locations.

It is important to realize that those numbers represent address locations for random-access memory (RAM) built deep inside the TRS-80 keyboard/computer unit. The addresses refer to locations of data within a set of those RAM chips.

It just so happens that Radio Shack engineers linked each of those video memory addresses to a specific point on the crt screen. Tinkering around with the data in the video memory is thus tantamount to tinkering around with the video display. In fact, you cannot print any sort of character on the screen without placing the appropriate character data into some address location in the video memory.

The correspondence between video memory addresses and geometric locations on the video screen is very systematic and straightforward. Not all home computers use schemes that are so systematic and straightforward.

Whenever a program places some character data into video memory location 15360 (the lowest video memory address), the character appears at the extreme upper left-hand corner of the screen. And by putting some valid character data into memory location 16383 (the highest video memory address) the character appears in the extreme lower right-hand corner of the screen.

The video memory occupies 1024 RAM addresses, running sequentially from 15360 through 16383. The video screen is arranged with a 16-line, 64-characters-per-line format—which also is 1024 character locations. There is a nice one-for-one correspondence between those addresses and character positions on the screen.

The first 64 video memory addresses are linked to the first line on the screen; the second 64 memory addresses are linked to the second line on the screen; and so on through all 1024 addresses/locations.

Fig. 2-3 is a rough representation of the TRS-80 video screen format: the 16-line, 64-characters-per-line format. The numbers down the left-hand column show the video memory addresses for the first character position in each line, and the numbers down the right-hand column show the memory addresses for the last character position in each line.

Suppose that you want to print an X character near the middle



15360	-----> /-----	15423
15424	-----> /-----	15487
15488	-----> /-----	15551
15552	-----> /-----	15615
15616	-----> /-----	15679
15680	-----> /-----	15743
15744	-----> /-----	15807
15808	-----> /-----	15871
15872	-----> /-----	15935
15936	-----> /-----	15999
16000	-----> /-----	16063
16064	-----> /-----	16127
16128	-----> /-----	16191
16192	-----> /-----	16255
16256	-----> /-----	16319
16320	-----> /-----	16383

Fig. 2-3. POKE screen address format.

of the screen, which means putting the ASCII data for character X into a video memory location that is halfway between 15808 and 15871—the extreme ends of the middle line on the screen. Since there are 64 character spaces on the line, the midpoint is between 31 and 32 character spaces to the right, or that same number of address locations greater than 15808.

So if you POKE an X into video memory location 15808+32, or 15840, you will come very close to hitting the middle of the screen. Try it. Do a POKE 15840,88. That POKEs an 88 (ASCII code for X) into video memory address 15840.

You can POKE any ASCII character or TRS-80 graphic into any video memory address and see the character appearing at the addressed location on the crt screen.

This represents the most direct way of dealing with the video memory from BASIC. It takes advantage of the fact that each of the 1024 video memory addresses has a well-defined position on the crt screen. The same things can be done in the same general fashion in machine language, using LOAD instructions and hexadecimal addresses and data.

The technique can be summarized this way:

**POKE address, data**

where

*address* is an address in video memory (an integer between 15360 and 16383),

*data* is a character code (an integer between 32 and 191).

But that isn't the only way to address the video memory; putting some character data into that address and having the character

appear at a well-defined spot on the screen. It's the most direct way to do the job, but Radio Shack engineers saw fit to work out some other methods for doing the same thing in a less direct but generally easier fashion.

### **PRINT @ Access to Video Memory**

The PRINT @ statement accesses video memory very much the same way a POKE statement does. Like the POKE method just described, the PRINT @ statement includes both address and character data information, but there are some differences in the ways the addresses and data are presented.

Recall that the video memory occupies 1024 contiguous RAM locations between 15360 and 16383. When accessing the video memory via a POKE statement, you must specify an address with some integer in that address range. The PRINT @ method, however, calls for addresses between 0 and 1023.

The addressing portion of a PRINT @ statement does exactly the same thing as the addressing portion of a POKE statement, addressing the same set of 1024 locations in video memory. The only difference between the two portions is that the PRINT @ addresses have lower numbers.

0	-----> /----->	63
64	-----> /----->	127
128	-----> /----->	191
192	-----> /----->	255
256	-----> /----->	319
320	-----> /----->	383
384	-----> /----->	447
448	-----> /----->	511
512	-----> /----->	575
576	-----> /----->	639
640	-----> /----->	703
704	-----> /----->	767
768	-----> /----->	831
832	-----> /----->	895
896	-----> /----->	959
960	-----> /----->	1023

**Fig. 2-4. PRINT @ screen address format.**

So a PRINT @ address of 0 actually accesses video memory address 15360—the first, lowest-numbered one; and PRINT @ address 1023 actually addresses video memory location 16383 (the highest address location). Strictly speaking, then, a *PRINT @ address is simply a video memory address minus 15383*. The first line of 64 character spaces on the crt screen are thus accessed by PRINT @ addresses 0 through 63, the second line by 64 through 127, and so on.

See the screen format and corresponding PRINT @ addresses, for the extreme ends of the lines, in Fig. 2-4.

When a BASIC interpreter executes a PRINT @ address, it must add 15360 to the PRINT @ address in order to get to the actual video memory address. That bit of arithmetic, many believe, is a small price to pay for working with video address numbers that are smaller than the actual address numbers.

In all fairness it must be pointed out that there is more than the use of smaller address numbers involved in the PRINT @ statement. The data part of the statement can be much more complicated than is possible with a POKE statement.

The data portion of a PRINT @ statement must have a string character to it. POKE data must be offered in a character-code format, while PRINT @ data has to be equivalent to a string.

To illustrate this point, suppose you want to print an A character at the first character location on the first line of the screen. A POKE 15360,65 will do the job, and so will a PRINT @ 0,"A". So will a PRINT @ 0, ASC(65), letting the ASC function convert the character code into its string version. When the BASIC interpreter sees the data portion of a PRINT @ statement, it must convert it from a string version into a character-code version that can be effectively POKEd into video memory. This conversion, like the matter of adding 15360 to the PRINT @ address, eats up some execution time.

So the PRINT @ statement, although it has the same general format as POKE, runs much more slowly than a direct POKE does. But the PRINT @ statement is justified by the fact that it can deal with a long string (up to 255) of alphanumeric ASCII characters. The first character in the string is stored in the address indicated by the PRINT @ statement. Successive characters are then stored in successive video memory locations.

Suppose that you do a PRINT @ 0,"WHAT IS GOING ON HERE?" That string message will appear in the upper left-hand corner of the screen, because the first W in the message is loaded into PRINT @ address 0. The remaining characters in the message, if you care to count them, occupy PRINT @ addresses 1 through 21. The BASIC interpreter automatically increments an address counter to deposit the ASCII version of the characters into successive video memory locations. If you tried to print out that same message, using a POKE method, you would have to increment the memory addresses yourself. Look at this:

```
10 REM ** POKE STRING DEMO **
20 A0=15360:A1=16383
30 CLS
40 FOR A=A0 TO A1
50 READ C$
```

```

60 IF C$="#" THEN 60
70 POKE A,ASC(C$)
80 NEXT A
100 DATA W,H,A,T," " ,I,S," " ,G,O,I,N,G," " ,O,N," " ,H,E,R,E,?,#

```

This isn't the only way to go about writing string messages with POKE statements, but the alternatives aren't significantly simpler. Indeed, PRINT @ statements show some real promise when it comes to writing strings of characters.

Incidentally, the PRINT @ data is not limited to the ASCII character set; it can be used for displaying the TRS-80 graphic set as well. All you do is specify the graphic code (128–191) as a "string," by applying the CHR\$ function. For example, PRINT @ 543,CHR\$(191) prints a white rectangle near the middle of the screen—PRINT @ address 543, graphic code 191.

So as a point of interest, at least, you can see that it is possible to build some nice screen graphics using the PRINT @ statement.

Finally, the PRINT @ statement has one further advantage over its POKE counterpart: a programmer cannot inadvertently poke data into memory locations outside the video memory range. PRINT @ statements do not allow addresses outside the range of 0 to 1023. Attempting to do a PRINT @ 2000, *data*, for example, will simply result in an FC ERROR message. POKEing outside the video memory range risks a blowup of the program.

When it comes to doing video graphics and printing messages on the crt, the PRINT @ statement has a lot of advantages over its POKE counterpart. The only trade-off is the relatively slow speed of PRINT @—a feature made necessary by all the number crunching that takes place when compiling and executing a PRINT @. The following program clearly illustrates the difference in execution speed for POKE and PRINT @:

```

10 REM ** PRINT @/POKE GRAPHICS SPEED CONTEST **
20 CLS
30 PRINT "POKE RUN"
40 FOR A=15424 TO 15487
50 POKE A,191
60 NEXT A
70 PRINT:PRINT "PRINT @ RUN"
80 FOR A=192 TO 255
90 PRINT @ A,CHR$(191)
100 NEXT A
110 END

```

The first part of the program, lines 30 through 60, use POKE graphics to fill a line with solid white rectangles (character code 191). The second part, lines 70 through 100, do the same thing, but using PRINT @ graphics.

The two horizontal white bars are drawn in succession, but the differences in drawing speeds are fairly obvious.

The PRINT @ statement can be summarized this way:

```
PRINT @ address, data
```

where

*address* is a PRINT @ address between 0 and 1024,

*data* is a string-related character, variable, or string of characters.

### Relative Video Memory Addressing With Primitive PRINTs

POKE and PRINT @ statements access the video memory in an absolute fashion. That is to say, the addresses refer directly to some specific address in the video memory. The remaining techniques for working with the video memory use relative addressing.

Using relative addressing often simplifies the task of programming the video format, but it separates the program even further from the actual video memory system, and the execution time is lengthened as a result.

For our immediate purposes here, a primitive PRINT statement is a PRINT followed by " " (quote-space-quote) or nothing at all. A primitive PRINT changes the position of the cursor on the screen; and by doing a proper selection of PRINT and PRINT " " statements, it is possible to place the cursor at any desired point on the screen.

What does this have to do with video memory addressing? Recall that every character position on the screen has a well-defined video memory address. Doing a PRINT statement changes the setting of an address counter, incrementing it to deposit data into successive video address locations. The location of the cursor on the screen is a reflection of the next video memory address to be used.

As a specific example, suppose you execute this program:

```
10 CLS
20 PRINT " "; "X"
```

The CLS operation clears the screen and homes the cursor (sends it to the extreme upper left-hand corner of the screen—to video memory address 15360). Printing two spaces in succession, separated by a semicolon, moves the cursor two spaces to the right along the top line, and the X appears in the third space on the line. The character code for X ends up in video memory address 15362.

In this case you do not have to specify the exact address for the X character code—15362 for the POKE method, or 2 for the PRINT @ method. The two successive PRINT-space operations simply move the cursor two spaces from its starting position. An

internal cursor counter takes care of converting this relative address change into absolute direct memory addresses.

Here is a trick that is used quite often in BASIC programs aimed at doing some screen formatting:

```
10 CLS
20 PRINT:PRINT
30 PRINT "X"
```

As in the previous example the opening CLS statement sends the cursor to the lowest video memory address location. The two successive PRINT statements in line 20 cause the cursor to skip down two lines on the screen, and the X ends up at the beginning of the third line on the screen.

Unless instructed to do otherwise, every sort of PRINT statement in BASIC is terminated by a line-feed/carriage-return operation. In terms of the video memory addresses and the screen format (see Fig. 2-3) this means setting the video memory address counter to the next address representing the beginning of a new line. Refer to the addresses in the left-hand column in Fig. 2-3.

So the CLS statement sends the cursor to video memory address 15360. The first of the two PRINT statements in line 20 sends the cursor to address 15424, and the second one sends it to address 15488. The X is deposited at that location by line 30, and that operation (being a PRINT operation) ends by setting the cursor to address location 15552. You must refer to Fig. 2-3 to appreciate this explanation.

Primitive PRINT addressing is relative addressing in the sense that you do not have to specify an exact address anywhere along the line. The cursor is moved the specified number of lines or spaces from its initial position, and that initial position can be anywhere on the screen. The addressing began at 15360 in the previous examples only because the programs began with a CLS operation—one that takes the cursor to 15360. You can actually begin anywhere you want.

Not having to reckon with absolute video memory addresses often simplifies video-oriented programming. If, for instance, you want to skip two spaces on a line before printing another character, simply do a PRINT " "; " ". That will do the trick—the BASIC compiler takes care of figuring out the absolute memory addresses to be loaded with the space code (ASCII 32).

Now, it was mentioned earlier in this discussion that every PRINT statement is followed, automatically, by a line-feed/carriage-return operation—unless instructed to do otherwise. Instructing the system to inhibit that automatic line feed/carriage return is a matter of terminating the PRINT statement with a semicolon. That is how

it is possible to skip two character spaces on the same line by doing `PRINT " "; " "`.

The primitive `PRINT` scheme allows you to view the video screen as a two-dimensional coordinate system. You can, in other words, locate a character on the screen by thinking in terms of the number of lines down and the number of character spaces to the right. So if you want to print an `X` near the middle of the screen, you can do that by first homing the cursor (for convenient reference), skipping down eight lines (in a 16-line format) and 31 character spaces to the right along that line (in a 64-characters-per-line format). The program looks something like this:

```
10 CLS
20 FOR L=1 TO 8:PRINT:NEXT L
30 FOR S=1 TO 31:PRINT " ";:NEXT S
40 PRINT "X"
```

More compelling applications of the primitive-`PRINT` scheme are offered in the next chapter. The important point here is that `PRINT` statements offer a relative video memory addressing scheme and the opportunity to build displays based on a vertical/horizontal coordinate system.

To be sure, this scheme eats up more BASIC programming memory and execution time than `POKE` and `PRINT @` methods do, but, used on a small scale, the relative addressing feature of primitive `PRINT` formatting can simplify the original programming task.

### Horizontal Addressing With `PRINT TAB`

`PRINT TAB` statements allow you to simplify the task of specifying a printing location along a given line. If you do a statement such as `PRINT TAB(n)"X"`, where  $n$  is any integer between 0 and 63, you can place that `X` at any desired space location (0 through 63) along a given line on the screen.

```
10 CLS
20 PRINT TAB(31)"X"
```

This program prints the `X` very close to the middle of the top line on the screen. And if you want to put the `X` near the middle of the screen, try this:

```
10 CLS
20 FOR L=1 TO 8:PRINT:NEXT L
30 PRINT TAB(31)"X"
```

The first part of the program, lines 10 and 20, uses the line-skipping routine, described in the previous section of this chapter, to find the beginning of the middle line on the screen. The `PRINT TAB`

statement in line 30 both runs the cursor to the middle of that line and prints the *X* character. Using the PRINT TAB statement represents an improvement over doing the same thing with a series of PRINT " "; primitives.

Behind the scenes, a PRINT TAB(*n*) statement causes an absolute video memory address counter to *increment n* counts, effectively placing the cursor *n* places to the right of its starting position. The data specified by the PRINT TAB statement is automatically POKed into that newly counted video memory address location.

Incidentally, it is possible to specify PRINT TAB numbers up to 255, but it turns out that any over 63 are redundant. For instance, PRINT TAB(0), PRINT TAB(64) and PRINT TAB(128) set the character to the first character space on a line. So using PRINT TAB numbers larger than 63 changes nothing, and only adds to possible confusion on the part of the programmer.

As in the case of any PRINT statement, a PRINT TAB is automatically followed by a line-feed/carriage-return operation, unless followed by a semicolon. So if you do a PRINT TAB(31)"X", the *X* will appear in the middle of the current line, but the cursor will be resting at the beginning of the next line. By contrast, a PRINT TAB(31)"X"; prints the *X* at the middle of the current line and forces the cursor to remain at the next character space on that same line.

The fact that you can suppress the line-feed/carriage-return operation makes it possible to pull off some useful one-line tricks. Try this:

```
10 CLS
20 PRINT TAB(31)"X";
30 PRINT TAB(40)"Y"
```

This program prints the *X* near the middle of the first line on the screen. The semicolon at the end of line 20 inhibits the normal line-feed/carriage-return operation, so the *Y* appears in space 40 on the same line—nine spaces to the right of the *X*.

Unfortunately, it is not possible to set the absolute memory address counter backwards. It counts forward only. So, in the previous example, you will have some trouble if you print the *Y* at space position 40 first, and then try to print the *X* at space position 31. The *Y* will end up at position 40, but the *X* will appear at position 41, in spite of the fact you specify a PRINT TAB(31) for it.

In a manner of speaking, the *n* values in a PRINT TAB(*n*) statement represent a form of absolute addressing—absolute addressing with reference to character spaces on a given line on the screen. If you view a line of text as made up of locations 0 through 63, the TAB values refer absolutely to one of those 64 locations.



PRINT TAB statements can handle TRS-80 graphics as well as the standard ASCII code. If you want to print a rectangle of light at the middle of a line, just do a `PRINT TAB(31)CHR$(191)`. Any of the graphic-code numbers (128–191) can be worked into that `CHR$` data statement. You can thus use combinations of primitive PRINT statements (to select lines) and PRINT TAB statements (to select character spaces on a line) to do some graphics that combine alphanumerics and graphic symbols.

### **Relative Character-Space Addressing With TAB Control Codes**

The upper end of the TRS-80 character/control video set includes some numbers that Radio Shack calls “space compression codes.” From BASIC, these codes can be executed by a `PRINT CHR$(n)` statement, where  $n$  is an integer between 192 and 255, inclusively. But what do those so-called space compression codes do?

Used with `PRINT CHR$(n)`, the space compression codes function in a manner similar to `PRINT TAB(n)`. The codes indicate a number of spaces to be skipped (to the right) on a given line on the screen. Unlike PRINT TAB, however, the space compression codes have a relative addressing quality: the cursor is moved a given number of spaces, beginning from its current position, and not always from the beginning of a line.

There are 64 TAB control codes in the range of 192 through 255, and there are 64 character positions on each line of text on the video screen. The match of these two numbers is hardly incidental.

To see how the TAB control codes work, suppose the cursor is resting at space number 31 on a particular line on the screen—that’s near the middle of the line. But then you want to skip two more spaces to the right. There are two ways to go about it: you can do a `PRINT TAB(33)` or a `PRINT CHR$(194)`. The PRINT TAB version specifies an absolute character space that turns out to be two character units to the right of the original one—really, the 33rd space from the beginning of the line. `PRINT CHR$(194)`, on the other hand, moves the cursor two spaces to the right of its current position—there is no reference to the beginning of the line. But where did the 194 come from?

The TAB control codes allow you to move the cursor  $s$  spaces to the right, where  $s$  is an integer between 0 and 63, inclusively. The actual code number to be specified, however, is always equal to 192 (the smallest TAB control number) plus  $s$ . To move the cursor two spaces to the right of its current position,  $s=2$  and the TAB control code is  $192+2$ , or 194. That’s where the 194 comes from in the previous example.

The operation can be summarized this way:

$s$  is the number of character spaces the cursor is moved to the right of its current position.

Suppose you want to fill the top line on the screen with asterisks, each separated from the previous one by a space. The job could be done this way:

The PRINT asterisk-space combination in line 30 does the job. Or, if you are willing to sweat out some math, you could do the job with absolute PRINT TAB values:

Or, to take advantage of the TAB control codes, you can do this:

To be sure, the last example is not significantly simpler than the first one, but that is only because the situation calls for skipping just one space between the asterisks. If you want to skip four spaces, `PRINT CHR$(196)` is more elegant than a `PRINT "*"` .

Since the TAB control codes use relative space addressing, it is altogether possible to specify code numbers that will overflow a line, calling for moving the cursor more spaces than remain on a given line. The TAB control function handles that sort of situation by resuming the count on the next line. Anytime you follow a printing operation with `PRINT CHR$(255)`, for example, the next character will be printed directly below the previous one—the cursor will advance 63 spaces to the right of the current cursor position. One might argue that the same position on successive lines are separated by 64 character spaces, and not 63. That is quite true, but bear in mind that the cursor always rests one character space to the right

of the last-printed character; advancing the cursor 63 spaces from that point will position the next-printed character directly under the first one.

### **Working Directly With the Cursor Counter**

With the notable exception of the POKE method of working with the video memory, all the screen formatting techniques described thus far in this chapter take advantage of a cursor counter; a register in the Z-80 microprocessor that keeps track of the one-for-one relationship between video memory addresses and print positions on the crt. These cursor-oriented operations run more slowly than POKEs do, but they are generally easier to use.

It is possible to access that cursor counter, PEEKing at the cursor address or POKEing new cursor addresses into it. The cursor address is accessible in a section of dedicated RAM that Radio Shack calls the "video display control block." The cursor address is carried as a 2-byte number in RAM addresses 16416 and 16417. The least significant byte (lsb) of the cursor's address is saved in location 16416, and the most significant of the two bytes (msb) is in 16417.

So if you want to see the contents of that cursor-position register, you can get to it by doing something such as PRINT PEEK(16417), PEEK(16416). That will give the cursor position as it stands *after* completing that PRINT operation. (The suggested PRINT operation, you see, affects the position of the cursor and hence the contents of the registers PEEKed into.) At any rate, that operation will give you a convincing feeling that you can, indeed, get to the cursor-counting register.

Whenever you PEEK into the cursor register, you will find that the number from the msb position (address 16417) is an integer between 60 and 63. The number from the lsb position (address 16416) is an integer between 0 and 255. Those are the cursor screen coordinates in a decimal format.

The msb divides the screen into four sets of lines, each set having four lines. This accounts for all 16 lines on the video screen. The lsb divides each line into 64 character spaces: 0-63 indicates a character space on the first line of any of the four groups of lines, 64-127 indicates character spaces on the second line in each group of four lines, 128-191 is for the third line in each group, and 192-255 is for the last of the four lines in each group.

So if the cursor happens to be resting at the first character space in the first line on the screen, the msb of the cursor address is 60 and the lsb is 0. But if the cursor is at the second space on the second line, the msb is 60 and the lsb is 65.

Fig. 2-5 summarizes the cursor coordinate addresses for the beginning and end of each line on the screen.

MSB	LSB			MSB	LSB
60	0	-----✓	-----✓	60	63
60	64	-----✓	-----✓	60	127
60	128	-----✓	-----✓	60	191
60	192	-----✓	-----✓	60	255
61	0	-----✓	-----✓	61	63
61	64	-----✓	-----✓	61	127
61	128	-----✓	-----✓	61	191
61	192	-----✓	-----✓	61	255
62	0	-----✓	-----✓	62	63
62	64	-----✓	-----✓	62	127
62	128	-----✓	-----✓	62	191
62	192	-----✓	-----✓	62	255
63	0	-----✓	-----✓	63	63
63	64	-----✓	-----✓	63	127
63	128	-----✓	-----✓	63	191
63	192	-----✓	-----✓	63	255

THE MSB AND LSB FOR CURSOR COORINATES ARE  
STORED AT RAM ADDRESSES 16417 AND 16416, RESPECTIVELY

**Fig. 2-5. Cursor-counter screen address format.**

As mentioned previously, PEEKing into the cursor-position register and PRINTing the results can be a confusing operation because the PRINT operation, itself, advances the cursor position. A better way to make a convincing demonstration is by POKEing coordinates into the cursor-position register and having a PRINT operation print some well-defined character on the screen. Try this:

```

10 CLS
20 INPUT "MSB";MSB
30 IF MSB<60 OR MSB>63 THEN 20
40 INPUT "LSB";LSB
50 IF LSB<0 OR LSB>255 THEN 40
60 CLS
70 POKE 16416,LSB:POKE 16417,MSB
80 PRINT "X";

```

The program requests the msb and lsb coordinates from you. It is goofproofed to prevent you from specifying out-of-range figures. Once you have entered the coordinates, the program clears the screen and prints an X at the point you specified. You will see that the format follows Fig. 2-5 quite nicely.

This is an unusual tool for positioning the cursor on the screen. There are some fine applications for the idea of PEEKing into the cursor-address register, but they will have to await a later discussion.

There is a Level II function that does, in its own fashion, PEEK into the cursor-address register. That is the POS(*x*) function. POS(*x*) looks into the cursor-address register and returns an integer

between 0 and 63, indicating the horizontal position of the cursor at the moment the function is called. The  $x$  argument, incidentally, is a dummy argument and can be any alphanumeric character.

For future thinking it may be helpful to realize that *the msb and lsb numbers representing the cursor-address position are actually 2-byte versions of the actual video memory address*. The first cursor-address position, for example, is 60,0; and in terms of the video memory addresses, that first position is 15360. If you convert that 63,0 combination into hexadecimal, you get 3C00H. And by converting 3C00 hexadecimal into decimal, you get 15360. (These conversion techniques are described in greater detail in a later chapter.)

### SET/RESET Video Memory Addressing

The SET( $x,y$ ) and RESET( $x,y$ ) functions are fairly well documented in the standard TRS-80 literature, and there is no real need to dwell on the general applications here. Bear in mind that  $x$  is an integer between 0 and 127, and  $y$  is an integer between 0 and 47. The SET function turns on a spot of light at its designated  $x,y$  coordinates, and RESET darkens the spot at those character coordinates.

The BASIC interpreter translates those  $x,y$  coordinates into absolute video memory address locations, inserting one of the graphic-code numbers (128–191) as appropriate to the configuration of spots of light in a given character space. See the relationships between the  $x,y$  coordinates and screen positions in Fig. 2-6. Radio Shack's Video/Programming Worksheets spell out this relationship in a clearer and more useful fashion.

It is quite possible to write a BASIC program that will print out the graphic-code number and video address for each graphic element previously generated by the SET/RESET statements. That project is left to the reader as an exercise in applying much of the material presented thus far in this chapter.

### VIDEO CONTROL CODES

Data codes 0 through 31 do jobs other than print characters on the screen. They are video/line-printer function codes that call ROM subroutines for doing certain print-related tasks. See Table 2-2.

The control functions are normally generated by striking the special control keys on the keyboard, but they can be called during the execution of a BASIC program by doing a PRINT CHR\$( $n$ ) statement, where  $n$  is the control-code number from Table 2-2.

The following discussions deal only with the control codes that influence the cursor position on the crt and hence the current address of the video memory.

Y		X	
0	< 0	-----	127 >
1	< 0	-----	127 >
2	< 0	-----	127 >
3	< 0	-----	127 >
4	< 0	-----	127 >
5	< 0	-----	127 >
6	< 0	-----	127 >
7	< 0	-----	127 >
8	< 0	-----	127 >
9	< 0	-----	127 >
10	< 0	-----	127 >
11	< 0	-----	127 >
12	< 0	-----	127 >
13	< 0	-----	127 >
14	< 0	-----	127 >
15	< 0	-----	127 >
16	< 0	-----	127 >
17	< 0	-----	127 >
18	< 0	-----	127 >
19	< 0	-----	127 >
20	< 0	-----	127 >
21	< 0	-----	127 >
22	< 0	-----	127 >
23	< 0	-----	127 >
24	< 0	-----	127 >
25	< 0	-----	127 >
26	< 0	-----	127 >
27	< 0	-----	127 >
28	< 0	-----	127 >
29	< 0	-----	127 >
30	< 0	-----	127 >
31	< 0	-----	127 >
32	< 0	-----	127 >
33	< 0	-----	127 >
34	< 0	-----	127 >
35	< 0	-----	127 >
36	< 0	-----	127 >
37	< 0	-----	127 >
38	< 0	-----	127 >
39	< 0	-----	127 >
40	< 0	-----	127 >
41	< 0	-----	127 >
42	< 0	-----	127 >
43	< 0	-----	127 >
44	< 0	-----	127 >
45	< 0	-----	127 >
46	< 0	-----	127 >
47	< 0	-----	127 >

**Fig. 2-6. SET/RESET screen address format.**

**Table 2-2. Summary of Control Codes and Functions for the TRS-80**

Decimal Code	Control Function
0-7	None
8	Backspace and erase current character
9	None
10	Line feed/carriage return
11, 12	Go to top of next form (line-printer function)
13	Line feed/carriage return
14	Turn on cursor
15	Turn off cursor
16-22	None
23	Convert to 32 characters per line
24	Backspace the cursor
25	Advance the cursor
26	Downward line feed
27	Upward line feed
28	Home the cursor
29	Cursor to beginning of current line
30	Erase to end of current line
31	Clear to end of frame

### **Line Feed/Carriage Return**

Generally speaking, a line feed/carriage return drops the cursor down one line and then moves it to the beginning of that line. An important exception occurs when the cursor is resting on the bottom line of the crt display; doing a line feed/carriage return at that point scrolls the entire display up one line and moves the cursor to the beginning of that bottom line.

Striking the ENTER key generates a line-feed/carriage-return code (code No. 13) anytime the system is responding to characters typed on the keyboard. And when a BASIC program is running, any PRINT statement that is not terminated with a semicolon will call a line feed/carriage return.

You can insert a line feed/carriage return anywhere you wish in a BASIC program by either

```
PRINT CHR$(10);
```

or

```
PRINT CHR$(13);
```

Yes, indeed, you can accomplish the same effect with a simple PRINT primitive, but we are trying to illustrate a principle here that will prove useful later on.

Relating this line-feed/carriage-return operation to the video memory address, the line-feed portion effectively adds 64 to the

current cursor address—this drops the cursor down one line. The carriage-return part of the operation then backs up the address to the point where statement `POS(X)` would return a 0—to the beginning of that line.

When a line feed/carriage return calls for a scrolling of the entire display, the cursor address is kept at the first space on the bottom line, but a 64 is subtracted from the addresses for all characters, except those on the top line. All characters, in other words, undergo an upward line feed, and those on the top line are lost.

### **Home the Cursor, Clear to End of Frame**

Control codes 28 (home the cursor) and 31 (clear to end of frame) can be used separately, but one of their most useful applications use them together.

Homing the cursor sends it to the lowest video memory address: to the extreme upper left-hand corner of the screen. You can implement the function in a BASIC program by doing a

```
PRINT CHR$(28);
```

Any successive printing operations will then be referenced to this home point. The semicolon at the end of the statement is necessary to prevent an automatic line feed/carriage return—an operation that will carry the cursor to the beginning of the second line from the top of the screen.

In terms of video memory addressing, doing a home-cursor operation simply sets the cursor address counter to 15360, or 60,0 if you want to consider the 2-byte address format.

Control code 31 is not really a cursor-moving operation, but it clears the screen from the address of the cursor to the highest video memory address. So if the cursor happens to be situated in the middle of the screen when a BASIC program encounters

```
PRINT CHR$(31);
```

the cursor does not move, but the remainder of the line and all successive lines to the bottom of the screen are cleared.

As far as the internal working of the computer is concerned, doing a clear-to-end-of-frame operation inserts data code 32 (a space) into all address locations from that of the cursor to the highest video memory address (16383).

As mentioned earlier, control codes 28 and 31 are often used together and in that particular sequence. Insert this statement into a BASIC program:

```
PRINT CHR$(28);CHR$(31);
```



That sequence homes the cursor and clears to the end of the frame. In this case, clearing to the end of the frame amounts to clearing the entire screen. The sequence does the same thing as the CLS statement or striking the CLEAR key when the system is in its command mode.

### **Move to Beginning of Line and Erase to End of Line**

Control codes 29 (move the cursor to the beginning of the current line) and 30 (erase to end of current line) work just like codes 28 (home the cursor) and 31 (clear to end of frame). Codes 29 and 30 deal only with the current line of text, however, while 28 and 31 deal with the entire screen.

Whenever you want to return the cursor to the beginning of its line, without erasing any characters along the way, insert this statement into a BASIC program:

```
PRINT CHR$(29);
```

This operation sets the cursor-address counter back to an address representing the beginning of the current line. If POS(X) happens to return a 31 before the operation is executed, it will return a 0 after the operation is done.

Erasing all the characters between the cursor and the end of its line is a matter of writing this statement into a BASIC program:

```
PRINT CHR$(30);
```

This causes the computer to insert a character code 32 (space) into each address location, beginning from the cursor and ending at the extreme right end of the current line. The cursor position is not affected by that clear-to-end-of-line operation.

Erasing an entire line of text, regardless of the cursor's position on that line, is a matter of using these two codes in succession:

```
PRINT CHR$(29);CHR$(30);
```

The first part of the statement moves the cursor to the beginning of the current line, and the second part erases everything to the right of it on that same line. In effect, the statement clears the entire line. The cursor remains at the beginning of the line, perhaps giving you the opportunity to replace it with fresh characters.

### **Turn On/Turn Off the Cursor**

When one is writing a BASIC program or working with the system in its command mode, the TRS-80 automatically turns on the cursor symbol to indicate the printing position of the next character. The cursor figure, however, is normally turned off while running a

BASIC program. Whenever you execute a program that generates a long series of alphanumeric characters, for example, you do not see the cursor figure as it buzzes along, indicating the next character-printing position.

It is possible to turn on the cursor figure and observe its position on the screen during a BASIC program by inserting the statement:

```
PRINT CHR$(14);
```

After that, you will be able to see the cursor, wherever it goes. As you might imagine, being able to see the cursor during the execution of a program can be a valuable aid in keeping track of what is going on and in editing screen graphics.

The turn-on-cursor control code can be executed just one time, and it will remain in effect until some other operation forces it off. So a `PRINT CHR$(14);` is normally inserted at the beginning of a BASIC program.

Turning off the cursor is a simple matter of executing the BASIC statement:

```
PRINT CHR$(15);
```

After seeing this statement, the cursor remains turned off until another turn-on operation occurs.

The next section of this chapter illustrates some important applications of the cursor turn-on control code.

### **One-Step Cursor Operations**

Control codes 24, 25, 26, and 27 increment the cursor position back one space, forward one space, down one line, and up one line. The contents of the video memory are left unchanged in the process.

Here is a summary of these operations as implemented in BASIC:

`PRINT CHR$(24);`—This statement backspaces the cursor position by effectively subtracting a 1 from the cursor-position address.

`PRINT CHR$(25);`—This statement advances the cursor position one space, effectively adding a 1 to the current cursor address.

`PRINT CHR$(26);`—This statement drops the cursor position down one line, but maintains the same horizontal position. This is done by effectively adding 64 to the current cursor address. The downward line-feed operation wraps around to the top of the screen whenever the cursor is on the bottom line.

`PRINT CHR$(26);`—This statement moves the cursor position up one line, but maintains the same horizontal position. The cursor address, in effect, has a 64 subtracted from it. Doing this upward line feed while the cursor is already on the top line causes it to wrap around to the bottom line.

These statements can be written into a BASIC program to achieve some useful formatting effects, doing things such as reserving a portion of the screen for in-line text printing and another portion for video graphics. This is sometimes called *split-screen formatting*, but the standard Radio Shack literature says little, if anything, about it.

It turns out that the four arrow keys on the TRS-80 keyboard will generate the one-step cursor motion codes—if you strike them while depressing the `SHIFT` key. The keys do not generate the right cursor-control codes when the system is in the command mode, however. You must use them in a BASIC execution mode. Here is a little program to demonstrate the point.

```
10 PRINT CHR$(14);CHR$(28);CHR$(31);
20 C$=INKEY$;IF C$="" THEN 20
30 PRINT C$;
40 GOTO 20
```

Line 10 turns on the cursor symbol, sends it home, and clears to the end of the frame. Line 20 then picks up a string character from the keyboard, and line 30 executes the `PRINT` operation on the key character.

As long as you are striking the alphanumeric keys you will see the characters printed on the screen. The cursor symbol will be there to indicate the position of the next-printed character.

The main point of the demonstration is the use of the arrow keys, though. So depress the `SHIFT` key and strike one of the arrow keys. You will see the cursor responding as outlined earlier in this section. The back arrow makes the cursor move to the left, the up arrow makes it move up one line, and so on. You don't get the same sort of response if the `SHIFT` key is *not* depressed.

Also note that you can use this shift/arrow-key feature to move the cursor through printed text without altering it. It deletes the character at its current position, but replaces it when you move the cursor somewhere else on the screen.

This suggests some elementary word-processing operations, giving you an opportunity to edit characters in the text without messing up a whole line of them. You can replace a character with another one by setting the cursor to that character position and striking the key

representing the desired character. Or you can delete the character by striking the space bar.

### **Backspace and Delete Current Character**

Control code 8 works just like the backspace operation just described (control code 24), except that it deletes the character as it goes along. Doing `PRINT CHR$(8);` is the same thing as striking the back-arrow key when the system is the command mode of operation.

## CHAPTER 3

# The Keyboard Environment

Just as the most-used output device for the TRS-80 is its crt screen, the most-used input device is its keyboard. The keyboard environment is neither as sophisticated nor as versatile as the video environment, but understanding its operating features can go a long way toward building some especially useful programs.

In the command mode of operation the keyboard is normally linked to the video display system. Every keystroke produces a well-defined response on the screen—printing characters, spacing, line feeding, and so on. See the SPECIAL NOTE on page 45 for an exception to this general idea.

There are two BASIC statements that let you interact with an on-going program from the keyboard: INPUT and INKEY\$. Both statements are well documented in the standard TRS-80 literature, but the first two sections of this chapter consider them in more detail, illustrating some applications that can be quite useful.

As you might suspect, there are some more subtle ways to deal with the keyboard, PEEKing around in the so-called keyboard memory part of the TRS-80 memory map. As in the case of such schemes in the video environment, some of these PEEK-keyboard routines will appear a bit cumbersome when used in BASIC, but understanding the principles will make it easier for you to grasp the essence of some keyboard operations in machine language. The latter part of the chapter is devoted to some of these unconventional keyboard-input notions.

### THE STANDARD INPUT STATEMENT

The INPUT statement allows in-line interaction with a program, halting the execution of that program until the user strikes the

### *SPECIAL NOTE*

It is possible to link the keyboard directly to a standard Radio Shack line-printer system by doing:

**POKE 16414,141:POKE 16415,5**

Doing this, all PRINT and LIST statements and commands affect the line printer instead of the video display. In effect, it replaces the video display with the line printer.

Since the line printer responds only after seeing the end of a line of text, it prints nothing until you strike the ENTER key or it finds the end of a program line.

Unfortunately, this scheme doesn't always work with nonstandard line-printer systems—those using printer-driving software. In such instances, doing the suggested POKEs sends the system “into outer space,” and the only way to recover is by switching the computer off and on again.

If the scheme works for your line-printer system, you can always return to the normal video link by doing:

**POKE 16414,88:POKE 16415,4**

ENTER key. The main purpose of INPUT is to allow the operator to enter some necessary numerical or string values.

### **Nature of the INPUT Statement**

The *BASIC* INPUT statement has the general form:

**INPUT “*message*”;*variable***

The *message* portion of the statement is optional, so it can be expressed as:

**INPUT *variable***

The *message* option lets you prompt the operator, spelling out the nature of the *variable* to be entered from the keyboard. You can, however, accomplish the same sort of thing with this sequence of statements:

**PRINT “*message*”;INPUT *variable***

On encountering an INPUT statement the system prints a question-mark prompt symbol, followed by the cursor symbol. The question mark indicates that the system is expecting data from the keyboard, and the cursor symbol marks the point on the screen

where that data will appear as the user types it in. The prompt and cursor symbols are generated in ROM, and there is nothing you can do to change their configuration, but you do have some options regarding their position on the screen.

When using the INPUT *"message";variable* format, the prompt and cursor symbols always appear in the two character spaces following the end of the *message*. The semicolon is responsible for this business of inhibiting a line feed/carriage return that would otherwise place the prompt and cursor symbols at the beginning of the next line on the screen. Try as you might, however, you cannot link a *message* with a *variable* in an INPUT statement without using the semicolon—a SYN ERROR message shows up every time.

But if you leave out the *"message";* part of the INPUT statement, you do have some choice in the matter of placing the prompt and cursor symbols. Compare these two sets of operations:

```
PRINT "message";INPUT variable
```

and

```
PRINT "message":INPUT variable
```

In the first instance, the prompt and cursor symbols appear on the same line as the *message*. The semicolon inhibits the line feed/carriage return that normally follows a PRINT statement. In the second case the prompt and cursor appear at the beginning of the line below the *message*.

While the PRINT *"message"/INPUT variable* combination is more cumbersome than a simple INPUT *"message";variable* statement, the former allows for a lot of different kinds of screen formatting. Suppose, for instance, you want the prompting message to appear at the upper left-hand corner of the screen, but, for some reason, you want the keyboard entry for *variable* to appear near the bottom, left-hand corner.

```
10 PRINT @ 0, "INPUT YOUR VALUE";  
20 PRINT @ 832,"";  
30 INPUT V
```

Line 10 PRINTs the prompting message in the upper left-hand corner of the screen and inhibits the line feed/carriage return. Line 20 prints nothing at PRINT @ location 832 (double quotes with no space between them is a nonprinting "null" character) and inhibits the line feed. The prompt and cursor for the INPUT statement in line 30 thus appear at PRINT @ location 832.

The nice feature here is that you can prompt an input and accept it at two different places on the screen without affecting any text or

graphics that might be situated in between. The same sort of thing can be done with many of the cursor-moving video operations described in Chapter 2.

Just bear in mind that an INPUT statement is not fully executed until the user ends the input operation by striking the ENTER key. And striking the ENTER key causes a line feed/carriage return that will scroll the display if the INPUT cursor is on line 14 or 15.

The *variable* portion of the INPUT statement specifies the variable type and assigns the keyboard input data to it. The variable can be either a numeric or string variable, the former expressed without a dollar sign and the latter with a dollar sign. In either case the keyboard entry operation is not complete until the user strikes the ENTER key.

It is possible to use a single INPUT statement to enter values for more than one variable; simply separate the variable names/types by a comma:

```
INPUT NL$,I$,A
```

That INPUT statement will expect three different variables from the keyboard: two strings, followed by a numerical variable. The values for the variables can be entered one at a time or in a sequence.

To illustrate the point consider an example. Suppose that NL\$ represents a person's last name, I\$ is the middle initial, and A is his or her age. If you decide to enter these variables one at a time, the screen format shows a single question-mark prompt symbol when it is time to enter the first variable:

```
? JONES      <ENTER>
??_
```

Typing in the name and striking the ENTER key, you note that the double question-mark prompt appears at the beginning of the next line. The system is expecting the second variable.

```
? JONES      <ENTER>
?? F         <ENTER>
??_
```

On entering the second variable, the prompt and cursor symbols for the third one appear at the beginning of the next line. The number of question marks does not increase beyond two, however. At this point the machine is expecting a numerical variable—the “age” variable in this example.

```
$ JONES      <ENTER>
?? F         <ENTER>
?? 23        <ENTER>
```



Now the INPUT operation is done, and the program goes about its business.

The user can respond to the INPUT statement in a different, and perhaps simpler, fashion. Instead of doing an ENTER after typing each variable, it can be done this way:

```
? JONES,F,23      <ENTER>
```

Separating the variables by a comma does the same thing as ending each variable-entry with an ENTER. Here the user does just one ENTER; and, what's more important in some graphic formats, the whole entry operation takes up just one line.

A complete INPUT statement for this example ought to look like this:

```
INPUT "LAST NAME, MIDDLE INITIAL, AGE";NL$,I$,A
```

This way, everything, including the prompting message, appears on the same line.

### **Some Special Applications of INPUT**

The INPUT statement is a program stopper. Once the system encounters an INPUT statement, everything comes to a halt until the user enters the proper kinds of variables and at least one ENTER keystroke. It is possible to take advantage of the program-stopping feature, using INPUT as a control operation rather than a data-entry operation.

Suppose you have a program that prints out a very long list of data—a list containing more lines than the video display can accommodate at one time. The real problem is that the user probably cannot read and interpret the information as it scrolls up the screen so rapidly.

There are a couple of ways to handle this sort of situation, both using the INPUT statement. To see how they work, suppose you want to inspect the content of addresses 0 through 1023 in your TRS-80 ROM. That represents a lot of lines of information, but here are some programs for doing the job.

```
10 CLS
20 A=0
30 PRINT A,PEEK(A)
40 A=A+1
50 IF A>1023 THEN END
60 GOTO 30
```

This program will print out the addresses and data for ROM locations 0 through 1023, but too quickly to be meaningful. So you might try this:

```

10 CLS
20 A=0
30 INPUT S$
40 PRINT A,PEEK(A)
50 A=A+1
60 IF A>1023 THEN END
70 GOTO 30

```

In this program the system does not print out the next line of information (address and data) until you strike a key on the keyboard. The INPUT statement in line 30 brings everything to a halt until a string variable is entered. It makes no difference what that variable is, because it is not used anywhere else in the program.

The point is that an INPUT statement used in this way lets you display the text one line at a time, giving you plenty of time to inspect the data on the screen. To make the display look a little neater, try putting a semicolon at the end of the PRINT statement in line 40. That way, the information appears on successive lines, and not on alternate ones.

Another way to handle the situation is to let the screen fill with 15 lines of text before the program calls for a keystroke to show the next 15 lines. Try this:

```

10 A=0
20 CLS
30 FOR L=1 TO 15
40 PRINT A,PEEK(A)
50 A=A+1
60 IF A>1023 THEN END
70 NEXT L
80 INPUT S$
90 GOTO 20

```

Now the FOR...NEXT loop between lines 30 and 70 prints 15 lines of information before the INPUT statement in line 80 brings the action to a halt. At that point you can inspect the data at your leisure and then strike any key to view the next 15 addresses.

As nice as the INPUT statement might be in many instances, the fact that it interrupts the program can be a real nuisance in other kinds of situations. That's where the INKEY\$ statement comes into play.

### THE STANDARD INKEY\$ STATEMENT

The INKEY\$ statement causes the computer to scan the keyboard, detecting any keystroke that might occur during the scanning, or strobing, interval. And using a BASIC statement such as C\$=INKEY\$, a string version of the key that is depressed will be assigned to that string variable, C\$.

The nice thing about INKEY\$, compared with INPUT, is that INKEY\$ does not necessarily interrupt the execution of the program. It sidetracks the program during the keyboard-strobing interval, but it doesn't have to bring everything to a halt as the INPUT statement does.

Used as intended, INKEY\$ can be a useful and powerful BASIC statement. Some programmers, however, attempt to use INKEY\$ outside its intended realm, and they get into trouble as a result. Perhaps that accounts for some of the nasty comments that writers sometimes publish about the TRS-80 INKEY\$ function.

One of the first things to bear in mind about INKEY\$ is that it is a string-related statement. An ever-present reminder of that fact is the presence of the dollar-sign symbol at the end of the expression. Getting numerical values into the system via an INKEY\$ function calls for an accompanying VAL(*string*) statement.

Another important feature is that INKEY\$ allows the entry of just one character at a time. Inputting multiple-character values into the system via INKEY\$ calls for executing the statement more than one time and assembling the individual characters as appropriate.

Finally, the system executes the keyboard strobing operation in a very short period, and any program calling for doing a single keystroke during that strobing interval is bound to cause some trouble—making it necessary to tap the key an undetermined number of times before the keystroke and strobe happen to take place simultaneously.

Programmers who neglect any of these three basic notions about INKEY\$ find themselves having some problems with it. Others, who understand these notions, can use the statement quite effectively.

### Dealing With the String Nature of INKEY\$

Here is a continuous, integer-counting program that uses an INKEY\$ function to adjust the time delay between successive counts—without having to interrupt the execution of the program:

```
10 CLS
20 N=0
30 C$=INKEY$:IF C$="" THEN 50
40 C=VAL(C$)
50 FOR T=0 TO 10*C:NEXT T
60 PRINT N
70 N=N+1
80 GOTO 30
```

The statements in line 50 generate the time delay between counts, and you can see that the delay interval depends on the

numerical value of variable C. The larger is the value of C, the longer the delay between successive counts.

The value of C is determined by the statements in lines 30 and 40. Line 30 uses an INKEY\$ statement to pick up a character (a string character) from the keyboard. If that character happens to be a null string (no key depressed), the program jumps down to line 50, allowing no change in the value of C. But if INKEY\$ turns up a numerical string character (0 through 9), the statement in line 40 converts it to a numerical quantity. That value has to be converted to a numerical quantity so that it can be mathematically manipulated in the time-delay loop, line 50.

The main point of this illustration is to show that it is necessary to convert a string-related INKEY\$ variable into a numerical value before it can be treated as a number. Any single-digit number can be inserted into an ongoing program this way, and there is no interruption of the kind that characterizes an INPUT statement.

Beginning programmers sometimes get into trouble with the idea because they forget to account for the null-string condition; they forget the conditional statement in the second part of line 30. Hack off the second statement in line 30, and you will not see the program running as it is supposed to run.

Now, here is another counting program that uses the INKEY\$ statement to adjust the direction of count. Once the program is operational, you can strike the *F* key to count forward, the *R* key to count in reverse, the *S* key to stop and hold the current count, and the *X* key to reset the count to zero and stop it. A rather straightforward decoding sequence translates the INKEY\$ string characters into counting-interval numbers.

```
20 N=0:I=0
30 C$=INKEY$:IF C$="" THEN 80
40 IF C$="F" THEN I=1
50 IF C$="R" THEN I=-1
60 IF C$="S" THEN I=0
70 IF C$="X" THEN 20
80 PRINT @ 0,USING "# # # #";N
90 N=N+I
100 GOTO 30
```

In the next example the INKEY\$-input string characters are converted into their ASCII codes before they are decoded as counting direction/speed parameters. The ASC function in line 35 converts a valid INKEY\$ character into ASCII form, and lines 40 through 75 set up the counting parameters—the I variable takes care of the direction of count, and variable D sets the delay between successive counts.

```

10  C!S
20  N=0:I=0:D=0
30  C$=INKEY$:IF C$="" THEN 80
35  C=ASC(C$)
40  IF C=70 OR C=102 THEN I=1
50  IF C=82 OR C=114 THEN I=-1
60  IF C=83 OR C=115 THEN I=0
70  IF C=88 OR C=120 THEN 20
75  IF C<95 THEN D=0 ELSE D=100
80  PRINT @ 0,USING "###";N
85  FOR T=0 TO D:NEXT T
90  N=N+I
100 GOTO 30

```

The program takes advantage of the fact that the ASCII codes for uppercase and lowercase characters are different. As far as the alphabetical characters are concerned, the lowercase versions are separated by their uppercase counterparts by decimal 32. Consequently, line 40, for example, is sensitive to ASCII 70 (uppercase *F*) or ASCII 102 (lowercase *f*).

So you can control the direction of count by striking the *F* or *R* keys (for forward and reverse, respectively). As in the previous example, striking the *S* key (ASCII 83 or 115) stops and holds the current count, and striking the *X* key (ASCII 88 or 120) stops the count and resets it to zero.

In any event the conditional statement in line 75 is sensitive to lowercase characters—those entered while the SHIFT key is depressed. That being the case, variable *D* is set to 100, and this inserts a noticeable time delay in the counting operation at line 85. Otherwise, the count runs at full speed.

It is also possible to write routines that convert INKEY\$-entered keystrokes into screen-control operations. You are free to assign single-character control codes as you choose. Look at this example:

```

10  C$=INKEY$:IF C$="" THEN RETURN <to calling program>
20  C=ASC(C$)
30  IF C=96 THEN DT=8
40  IF C=97 THEN DT=31
50  IF C=100 THEN 70
60  PRINT DT;:RETURN <to calling program>
70  PRINT CHR$(28);CHR$(31);
80  RETURN <to calling program>

```

This routine, written as a subroutine that is called by some mainline program, allows the user to strike the <SHIFT> *B* key to backspace and erase the current character on the screen (line 30), strike the <SHIFT> *C* key to clear the screen to the end of the frame (line 40), and strike the <SHIFT> *N* key to home the cursor and clear the entire screen (lines 50 and 70). The routine can be expanded

to cover all sorts of control functions, all specified by single key-strokes dreamed up by the programmer.

### **Putting Together Multiple-Character Values From INKEY\$**

One of the not-so-nice features of the INKEY\$ statement is that it permits the entry of just one keyboard character at a time. Single-character INKEY\$ inputs can be quite useful, as demonstrated in the previous discussion, but there are instances where it is more desirable and, perhaps, absolutely necessary to input variable values having more than one character. This is a bit troublesome when using the INKEY\$ function, but it often beats the program-stopping feature of an INPUT statement.

It is somewhat easier to enter multiple-character string values via INKEY\$ than it is to enter multiple-digit numerical values; so let's consider the easier part first.

The following program uses the INKEY\$ function to build up a string value, W\$, from individual keystrokes. The characters are entered one at a time and then assembled (concatenated) to build a single multiple-character string. The string is limited to five characters, and it will truncate the string beyond that point. Limiting the string to five characters is an arbitrary choice, however; the string can be as long as desired.

There are a couple of special features built into this program that are incidental to the main point of the discussion. While these "whistles and bells" add to the complexity of the program, they help make a convincing case for the overall usefulness of the main idea—inputting multiple-character strings from INKEY\$ statements.

```
10 REM ** INKEY$ STRING DEMO **
20 CLS:LN=0:PRINT @ 50,"#"
30 GOSUB 100
40 REM (PUT YOUR REGULAR PROGRAMMING HERE)
50 IF LN>15 THEN 20
60 GOTO 30
70 REM
80 REM
100 REM ** INKEY$ INPUT SUBROUTINE **
110 C$=INKEY$:IF C$="" THEN RETURN
120 C=ASC(C$)
130 IF C=8 THEN 180
140 IF C=13 THEN 200
150 IF LEN(W$)>=5 THEN 170
160 W$=W$+C$
170 PRINT @ 50,"#" + W$;:RETURN
180 FOR CP=50 TO 55:PRINT @ CP,CHR$(32);:NEXT CP
190 W$="";PRINT @ 50,"#";:RETURN
200 PRINT @ 0 + LN*64,W$;
210 LN=LN+1
220 GOTO 180
```

On running this program, you will see a pound sign symbol near the right end of the top line on the screen. This is one of those “whistles-and-bells” features; it is a homebrewed prompt symbol for entering the string characters.

Now, type up to five alphanumerics—combinations of letters and numbers. If you don’t like the string that is being built up to the right of the pound sign, strike the left-arrow key (←) to clear it. This clearing operation is another one of those optional features that are built into the program.

When you are satisfied with the appearance of your multiple-character string entry, strike the ENTER key. That operation clears your entry, resets the pound-sign cursor and, more importantly, prints your entry as a single string “message” at the left-hand side of the screen.

Now you can type in and ENTER another string of characters. As you do that, the list of strings builds downward along the left side of the screen. When the list reaches the bottom of the screen (after entering 15 strings), the list clears and begins again from the top.

The main point of the demonstration is that you can, indeed, enter multiple-character strings with the INKEY\$ function. Now it is time to dig into the program and see how the job is done.

The program is written with the keyboard-entry operations residing as a subroutine, beginning at line 100. The mainline program (the one that calls the INKEY\$-input subroutine) occupies lines 10 through 80.

After initializing the line counter (LN), clearing the screen, and printing the initial # prompt symbol, the mainline program calls the INKEY\$ subroutine at line 30. If you had some additional programming that called for in-line interaction with your input expressions, it would fit between lines 30 and 50.

Basically, the mainline program loops endlessly between lines 30 and 60, calling the INKEY\$ subroutine and doing any program operations you might enter in the area of line 40.

Turning to the subroutine itself, notice that the INKEY\$ characters are carried as variable C\$. The C\$ variable represents the single-character input. In line 160 this character is concatenated with string W\$. This is the point where the individual INKEY\$ characters are assembled into a single multiple-character string. In fact, that is the critical operation—the point of the demonstration. The rest is intended to make things run smoother for the operator.

The statement in line 120, for instance, converts the INKEY\$ input into an ASCII code number. And if C=8 (left-arrow key code), the program jumps to line 180, where the entry is cleared from the screen, W\$ is nulled, and the pound-sign prompt symbol is replaced

on the screen. This simply means that striking the left-arrow key erases the current entry, making it possible to correct any typing errors.

If C=13 in line 140, it means the user has hit the ENTER key. The response in that instance is determined by operations beginning at line 200. Those operations print the assembled string, W\$, on the next line of text along the left side of the screen, increment the line counter (LN), and go back to 180 to erase the entry version of the string.

This notion of concatenating single-keystroke characters from INKEY\$ to make multiple-character strings can lead to some powerful in-line keyboard operations. It allows you to make up multiple-character control expressions that can influence the ongoing activity. Try this program, making just a few additions to the original version (lines 52, 53, and 55).

```
10 REM ** INKEY$ STRING DEMO **
20 CLS:LN=0:PRINT @ 50,"#"
30 GOSUB 100
40 REM (PUT YOUR REGULAR PROGRAMMING HERE)
50 IF LN>15 THEN 20
52 IF W$="DONE" THEN END
53 IF W$<>"CLEAR" THEN 60
55 W$="":GOTO 20
60 GOTO 30
70 REM
80 REM
100 REM ** INKEY$ INPUT SUBROUTINE **
110 C$=INKEY$:IF C$="" THEN RETURN
120 C=ASC(C$)
130 IF C=8 THEN 180
140 IF C=13 THEN 200
150 IF LEN(W$)>=5 THEN 170
160 W$=W$+C$
170 PRINT @ 50,"#" + W$;:RETURN
180 FOR CP=50 TO 55:PRINT @ CP,CHR$(32);:NEXT CP
190 W$="":PRINT @ 50,"#";:RETURN
200 PRINT @ 0 + LN * 64,W$;
210 LN=LN+1
220 GOTO 180
```

Now the program is uniquely sensitive to the strings DONE and CLEAR. On entering DONE, the program immediately ends. Entering CLEAR clears the accumulated lines of text and starts everything from scratch.

This technique ought to suggest some intriguing program-control formats and, indeed, the possibility of making up custom programming languages—interactive ones, at that.

Inputting multiple-digit numerical values via INKEY\$ follows the same general scheme. The point of departure is where the string



version is converted into decimal numbers, and that calls for some math operations. See lines 200 through 230.

```
10 REM ** INKEY$—ENTERED NUMBERS **
20 CLS:LN=0:PRINT @ 50,"#"
30 GOSUB 100
40 REM (PUT YOUR REGULAR PROGRAMMING HERE)
50 IF LN>15 THEN 20
60 GOTO 30
70 REM
80 REM
100 REM ** INKEY$ INPUT SUBROUTINE **
110 C$=INKEY$:IF C$="" THEN RETURN
120 C=ASC(C$)
130 IF C=8 THEN 180
140 IF C=13 THEN 200
145 IF C<48 OR C>57 THEN RETURN
150 IF LEN(W$)>=5 THEN 170
160 W$=W$+C$
170 PRINT @ 50,"#" + W$;:RETURN
180 FOR CP=50 TO 55:PRINT @ CP,CHR$(32);:NEXT CP
190 W$=""':PRINT @ 50,"#"'::RETURN
200 W=0
210 FOR CP=LEN(W$)-1 TO 0 STEP -1
220 W=W+VAL(MID$(W$,LEN(W$)-CP,1))*10[(CP)
225 NEXT CP
230 PRINT @ 0+LN*64,W;
240 LN=LN+1
250 GOTO 180
```

The program develops string W\$ as before, but it is no longer enough to simply concatenate the new keystrokes to get the final result—the finished string has to be converted into a legitimate decimal format.

The decimal version of the number is first set to 0 at line 200. Variable W is the current decimal value of the number. Line 210 picks up the number of characters in the number and, one at a time, assigns them a power-of-10 value that is appropriate to their place in the numeric expression. By the time the program gets to line 230, W is equal to the number entered from the keyboard. It is a true numerical value, as suggested by the fact it is PRINTed as a numerical, rather than a string, variable.

A further modification in line 145 prevents the operator from entering nonnumerical figures.

As you might imagine by now, this technique can be useful for entering the values of numerical variables while some program is running. You can, for instance, change the parameters in a graphing program while the graph, itself, is being generated.

It is very difficult to imagine a more powerful interactive programming technique in BASIC. The INKEY\$ statement is perhaps the

most powerful interactive tool in BASIC, but one of the most maligned. It's too bad so many programmers miss the opportunity to use it.

### **The Loop Requirement for INKEY\$**

An INKEY\$ statement must be placed within a looping operation. If you inspect all the programs suggested thus far for the INKEY\$ statement, you will find it is in a loop of some sort.

INKEY\$ must be placed into a loop to make its operation reliable. The keyboard strobing operation that is called by an INKEY\$ statement runs so quickly that there is very little chance a keystroke will occur at the precise instant INKEY\$ is scanning the keyboard.

Here is an example of a very tight INKEY\$ loop:

```
10 C$=INKEY$:IF C$="" THEN 10
```

In this instance the program “buzzes” on line 10 until a keystroke occurs. It is the same thing as inserting an INPUT statement into a program. The program effectively comes to a halt until the system senses a keystroke.

### **SENSING KEY DEPRESSION WITH PEEK(14463)**

While INPUT and INKEY\$ serve most keyboard input functions quite nicely in BASIC, there are a couple of useful operations that are foreign to those two statements. Some of those operations can be utilized by working with the so-called keyboard memory.

The keyboard memory is a section of the memory map, between 14336 and 15359, inclusively. One of these addresses that is of particular importance is 14463. That address location carries data zero unless one or more keys are depressed. Set up this simple demonstration program, and try things for yourself:

```
10 CLS
20 C=PEEK(14463)
30 PRINT C;
40 GOTO 20
```

When you start this program, you will see strings of zeros being printed across the screen—until you depress a key. The moment you depress a character key (some control keys do not affect 14463), you will see some number other than zero being printed.

What is of special interest here is the fact that the program continues printing those numbers as long as the key is depressed. By contrast, the INPUT and INKEY\$ statements are “stroke” operated. Unless your system is suffering from keybounce, INPUT and

INKEY\$ are one-shot operations; they can tell a key has been depressed, but the duration of a key depression is not relevant.

So a PEEK (14463) gives you a control function that is not described in the standard TRS-80 literature. In effect, the entire keyboard becomes a set of normally open push-button switches. Try this:

```
10 CLS
20 C=PEEK(14463)
30 IF C=0 THEN 20
40 PRINT CHR$(191);
50 GOTO 20
```

This little routine draws a white bar as long as a key is depressed. Release the key, and the drawing stops. It isn't quite so easy to do such a thing with a togglelike INKEY\$ statement. This kind of routine is invaluable whenever you want some action to take place while a key is depressed.

There are some other features of PEEK(14463) that could prove useful as well. Whenever you do a PRINT PEEK(14463) and depress a single key, you see one of the following numbers: 1, 2, 4, 8, 16, 32, 64, or 128. Several different keys produce the same number in this particular family of numbers. PEEK(14463) turns up a 2, for example, whenever you depress A, I, Q, Y, or 1. On the other hand, it turns up a 4 whenever you depress B, J, R, Z, or 3. Table 3-1 summarizes the content of address 14463 for all of the useful key functions.

The code numbers picked up at address 14463 can be decoded and used for specific control applications. Suppose, for example, you want to move a "paddle" figure up and down on the screen, in response to depressions of the up-arrow or down-arrow keys. Depressing the up-arrow key ought to make the figure move up the

**Table 3-1. Decimal Content of Address 14463 While a Specified Key Is Depressed**

Contents of 14463	Key Depressed
0	(No key depressed)
1	@ H P X 0 ( 8 ENTER
2	A I Q Y 1 1 ) 9 CLEAR
4	B J R Z " 2 * : 8BREAK
8	C K S # 3 + ; ↑
16	D L T \$ 4 < , ↓
32	E M U % 5 = - ←
64	F N V & 6 > . →
128	G O W ' 7 ? / SPACE

NOTE: <SHIFT> @ produces a 1 on alternate key depressions. An alphabetical character with SHIFT produces the same value as its non-SHIFT version.

screen, and depressing the down-arrow key should make the figure move downward.

The basic idea in this case is to PEEK(14463) and take appropriate action if that PEEK turns up an 8 (up-arrow depression) or a 16 (down-arrow depression). The following example is becoming a classic demonstration for the PEEK(14463) operation:

```
10  BL=15360:TL=16383:P=15557
20  CLS:POKE P,191
30  M=PEEK(14463):IF NOT(M=8 OR M=16) THEN 30
40  IF M=8 THEN PT=P-64 ELSE PT=P+64
50  IF PT<BL OR PT>TL THEN 30
60  POKE P,32
70  P=PT:POKE P,191
80  GOTO 30
```

Line 30 looks at address 14463 and loops around on that same line until it turns up a value of 8 or 16. At that time line 40 sets the next position for the paddle figure, and line 50 checks for the possibility of running the figure out of video memory range.

Line 60 then erases the old paddle position (replacing it with a space), line 70 draws the figure in its new position, and line 80 loops the whole program back up to line 30, where it looks for the key depression again.

So depressing the up-arrow key develops an 8 at address 14463, and depressing the down-arrow key develops a 16. You will find, however, that some other keys work equally well with the paddle demonstration program. From Table 3-1 you will find a lot of other keys generating 8s and 16s. Those will cause the same action as the up-arrow and down-arrow keys.

The basic problem is that the key-depression sensing activity of address 14463 does not completely decode the keyboard. Rather, it decodes the keys into eight groups. Coupling a PEEK(14463) with an INKEY\$ statement, however, lets you reap the benefits of both: sensing continuous key depression with PEEK(14463) and completely decoding the character code with INKEY\$.

So here is a demonstration program that works as a REPEAT KEY operation. Depress any key listed in Table 3-1, and you will find it being repeatedly written on the screen as long as that key is depressed. The program is fun if you first fill the screen with arbitrary characters, and then "edit" the display by means of the cursor controls (working the arrow keys while depressing SHIFT).

```
10  CLS:PRINT CHR$(14);
20  IF PEEK(14463)=0 THEN 20
30  C$=INKEY$
40  IF C$<>" " THEN R$=C$
50  PRINT R$;
60  GOTO 20
```

In this program, line 10 clears the screen and turns on the cursor. Line 20 looks for a key depression—any key depression. That is half the job. On detecting a key depression, the INKEY\$ statement in line 30 decodes the key, and line 40 sets the value of R\$ so that the program continues drawing that same character long after the INKEY\$ function has reset itself. The program, in other words, continues generating the R\$ character as long as its key is depressed.

The keyboard can thus be completely decoded and yet respond in a push-button fashion that is lacking in either the INKEY\$ or INPUT functions, alone.

The PEEK(14463) operation has one more special feature that can prove useful. Table 3-1 indicates the content of address 14463 whenever *one* key is depressed. But when you depress more than one key at a time, address 14463 tends to show the sum of the basic key codes. Normally, depressing the A key turns up a 2 in address 14463, and depressing the B key generates a value of 4 in that location. Depress both of those keys at the same time, however, and you will find a value of 6 in address 14463. Depress A, B and C, and you will find 2+4+8, or 14, in that address.

So, by depressing more than one key at a time, address 14463 shows the sum of the individual key-depression codes—all the way up to 255 (you cannot get a number larger than 255, no matter how many keys you depress). That leads to some potential applications: setting the speed of an animated figure by depressing more than one control key simultaneously, for instance.

#### SUMMARY

PEEK(14463) will return an integer between 1 and 255 as long as any key, or combination of keys, is depressed.

Table 3-1 shows the content of address 14463 when *one* key is depressed. Depressing more than one key returns the sum of those numbers, up to 255.

### WORKING WITH THE KEYBOARD MATRIX

The TRS-80 keyboard is set up with an 8×8-key format. It at least has the potential for working with eight columns of keys, each having eight keys in them. Not all the rows are filled, however. See Fig. 3-1.

When your TRS-80 is in its command mode of operation some ROM programming routinely scans, or strobcs, the keyboard, looking for a possible key depression. And while a program is running in BASIC, the INPUT and INKEY\$ functions also call for a keyboard-strobe operation.

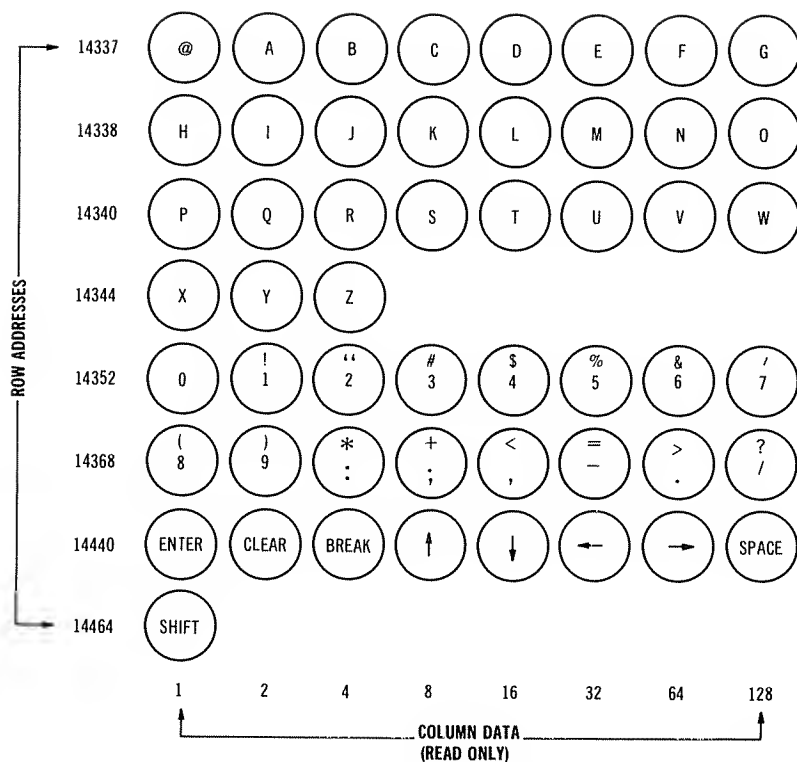


Fig. 3-1. The TRS-80 keyboard matrix (decimal format).

The system, in either case, scans the keyboard one row at a time. If a key in a given row is depressed, the system generates a code number representing that key. This is how the scheme works in a general sense.

The system PEEKs at the row addresses, in sequence and one at a time. Each PEEK returns a zero (if no key is depressed in that row) or a number representing the key that is depressed in that row.

Referring to Fig. 3-1, suppose that you are depressing the A key when the strobing action begins. When the system does a machine-language version of PEEK(14337) it will find a 2 in that location. All subsequent row addresses will turn up zeros. As another example, you might be depressing the P key. This being the case, PEEKs at 14337 and 14338 (addresses of the first two rows) will turn up zeros, but the system will find a 1 when it does a PEEK(14340)—the first bit in the third row.

Now, this strobing scheme can turn up key data that is the same for a number of different keys. The numbers in this case are

identical with those listed in Table 3-1. The system keeps everything straight, however, by keeping track of which row is being addressed at the moment. *A* and *I* keys both turn up a 2, but the system will see that 2 in row address 14337 only when the *A* key is depressed, and it will see the 2 in row address 14338 only when the *B* key is depressed.

To test your understanding of this idea, see if you can answer this question: Which key is being depressed if a 64 turns up in row address 14340? Answer: The *V* key is being depressed.

The machine-language keyboard strobing/decoding routine goes through the same procedure you used to answer that question. The strobing routine runs through all eight row addresses, whether it finds some data along the way or not. This is important because the two SHIFT keys generate a 1 at address 14464 whenever one of them is depressed. So if you are doing a SHIFT while holding down the *A* key, the system finds a 2 at address 14337 and a 1 at address 14464. If the SHIFT key is *not* depressed, address 14337 holds a 0.

The following program lets you interrogate the keyboard matrix. It is essentially a BASIC version of the ROM-based keyboard strobing operation.

```
10 REM ** KEYBOARD MATRIX DEMO **
20 IF PEEK(14463)=0 THEN 20
30 CLS
40 A=14337:PRINT A,PEEK(A)
50 FOR N=0 TO 6
60 P=2↑N
70 A=A+P:PRINT A,PEEK(A)
80 NEXT N
90 IF PEEK(14463)<>0 THEN 90
100 GOTO 20
```

Run the program, then depress a key. The program then scans the keyboard and prints the row addresses and corresponding data as it goes along. Depress the keys in any sequence you choose, including some while the SHIFT key is depressed. Just make sure you hold down the key until the display is completely generated. What you observe on the screen should line up with the information you can glean from Fig. 3-1.

Try depressing more than one key at a time, and see if you can figure out why the display responds the way it does. Compare the results with the discussion of the PEEK(14463) operation in the previous section of this chapter.

Incidentally, you will note from Fig. 3-1 that the row addresses grow successively larger, but the intervals between them are not identical. The row addressing follows a binary/hexadecimal se-

quence; thus there appears to be a lot of wasted “memory space” between row addresses 14400 and 14464.

The row addresses shown in Fig. 3-1 are necessary for addressing just one row at a time. If you PEEK any address between a pair of those row addresses, you are going to address more than one row at a time.

If you do a PEEK(14339), for example, you will address rows 14337 and 14338 at the same time. The result will be the sum of whatever keys might be depressed in those two rows. And what do you suppose happens if you do a PEEK(14463)? This addresses all rows (except the SHIFT row) at the same time, and that is a special case you have already studied in the previous section of this chapter.

Under normal operating conditions the keyboard routine has two phases: it first strobes the keyboard as described here, and then it decodes the information to generate a unique ASCII code for the key being depressed. The decoding is done in ROM, but it is difficult to get at that decoding routine from BASIC. You will be able to access the key-decoding routine rather easily from machine-language programs, thus saving yourself a lot of needless programming—programming aimed at decoding the keyboard.



## CHAPTER 4

# The User's Memory Environment

Every TRS-80 system has a certain memory capacity: 4K, 16K, 32K, or 48K. These figures refer to the amount of RAM space—the number of bytes—available for working BASIC programs or any other kind of memory operations.

In all cases the RAM available to the user begins at address 17129. The top address, however, depends on the system's memory capacity. See Table 4-1.

**Table 4-1. User's Available RAM Space**

System Capacity	Lowest Address	Highest Address
4K	17129	20479
16K	17129	32767
32K	17129	49151
48K	17129	65535

If you figure the difference between the lowest address and the highest address in Table 4-1, you won't come up with the number representing system capacity. These extremes indicate the range of RAM addresses that are available for user's programs.

The RAM space actually begins at 16384 in all Level II machines. The memory space between 16384 and 17129 is devoted to internal operations: restart vectors 1-7, device control blocks, the i/o buffer, and so on. See any of the memory maps supplied by Radio Shack for details.

Some of the demonstrations in Chapters 2 and 3 involved working in that lower part of the RAM space. The real working memory, however, begins at 17129, and that's the starting place for the discussions in this chapter.

The first part of the chapter looks into the i/o buffer and shows how BASIC statements are assembled in the user's memory space. The second part deals with the matter of the so-called protected memory—a section of the user's memory that is set aside for special applications through the MEMORY SIZE? entry routine.

### ORGANIZATION OF THE USER'S MEMORY SPACE

Fig. 4-1 is a general memory map of the memory space available for writing and executing programs. The range of this memory space is clearly defined for every machine, but the space devoted to the various categories of information stored there is not.

To be sure, BASIC program text always begins at address 17129 and builds upward from there. Program text is an almost-literal listing of a BASIC program; so the longer the program is, the farther it builds up into the user's memory space.

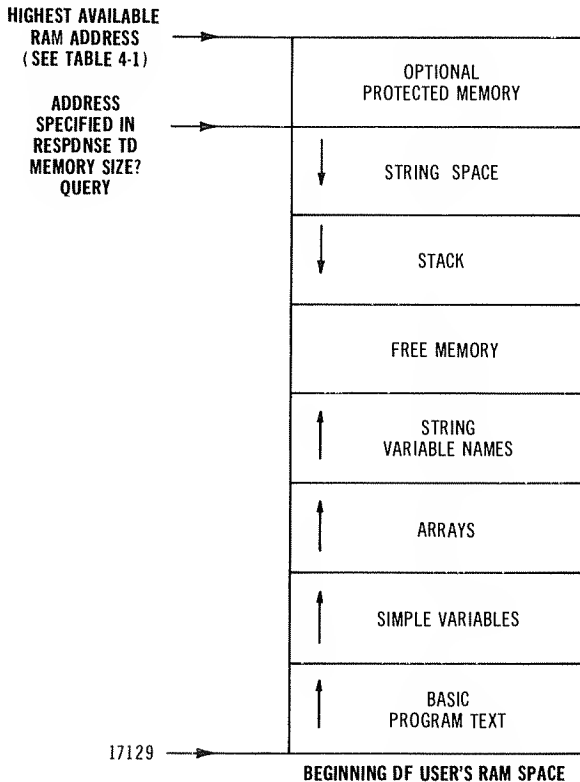


Fig. 4-1. Memory map for BASIC programs.

But there's more in the user's memory space than the program text. As the BASIC program is entered and executed, simple variables, arrays, and string variable names are appended to the top of the BASIC program text. A memory map cannot specify the starting address of simple variables, for example, because that address varies with the size of the program text. The same is true for the array and variable-name addresses.

The best we can do is illustrate their relative positions on the memory map.

There is also some of this dynamic sort of memory space allocated for strings and a stack. (The stack is responsible for keeping track of things while a BASIC program is running nested arithmetic, logic, and GOSUB routines.) These two memory spaces are not appended to the others, however. As indicated in Fig. 4-1, string space and stack begin at a high address and build downward. The more complex the string and stack operations become, the farther they grow down toward the lower end of the user's RAM space.

Incidentally, the starting address for the string-space allocation depends on the computer's memory capacity and any reply the user might make to the MEMORY SIZE? query. There's much more to say about that later in this chapter.

The main point of the present discussion is that during the entry and execution of a BASIC program, some of the program elements load upward into the user's memory space, and some begin at the top and grow downward.

It figures, then, that there is a no-man's land in between. This is the user's memory space that is commonly called *free memory*. There's no telling exactly where that free memory is located; all we know for sure is that it stands between the upward-growing variable-name space and downward-growing stack.

Even though it is pointless to figure the exact location of the free memory, you can always determine how large it is—how many bytes of memory are left in that space. To do this, simply ENTER a PRINT MEM. That's the purpose of the MEM statement.

The term "free memory" isn't really a misnomer, but it is often misunderstood. It must not be interpreted as a section of memory you are free to use as you wish; it is a bad idea to tinker around in the free memory. It isn't *free* in the sense that you are free to use it for your own purposes (the protected memory plays that role for you). Rather, its size indicates how much memory remains for entering and executing BASIC programs.

As the program text and stack grow, they decrease the free memory space. And that amount of free memory doesn't necessarily remain constant after a program is fully entered. It can vary during the execution of a program, especially if that program has a lot of

stack-related operations—nested math and logic operations or nested GOSUBs. Try this little demonstration program:

```
10 PRINT MEM
20 GOSUB 40
30 PRINT MEM:END
40 GOSUB 60
50 PRINT MEM:RETURN
60 GOSUB 80
70 PRINT MEM:RETURN
80 GOSUB 100
90 PRINT MEM:RETURN
100 PRINT MEM:RETURN
```

This program contains a number of nested subroutines, and each one of them prints out the amount of available free memory. You will see very clearly that the size of the free memory space changes during the execution of a stack-related program.

The practical significance of this demonstration becomes apparent when writing very long and complicated BASIC programs—programs that occupy nearly all available memory space. When you are through entering the program and do a PRINT MEM to check on the amount of free memory still available, you might get a comfortable figure. That can lull you into thinking that you did, indeed, manage to get that nasty program into your available memory. But then you might get a rude awakening to stack operations when you try running the program, and the stack space collides with the variable-names space—the program crashes for a lack of available memory.

Understanding the significance of the free memory space, and making certain that a long program leaves a generous amount of it can save some real disappointments.

The “protected memory” is an optional part of the user’s memory space. It is always situated at the very top of the available RAM space, beginning from an address specified during the MEMORY SIZE? power-up routine. If you respond to that power-up query by simply striking the ENTER key, there will be no reserved memory, and the string space and stack will grow downward from the highest available RAM address—addresses specified for you in Table 4-1.

But if you respond to MEMORY SIZE? with a valid RAM address, the string and stack spaces will begin building downward from that address; all available memory above that point will be reserved for things such as machine-language programs and any other memory tinkering you may want to do.

After running through the previous discussion of the free memory space, you ought to see the value in reserving as little protected

memory as possible. If you are overly generous and protect too much memory, you will leave little space for entering and executing programs in BASIC. In effect, every byte you reserve for protected memory steals a byte from the free memory; and we have already discussed what can happen when the system runs out of free memory.

In actual practice, however, a programmer rarely runs out of free memory because he or she has reserved some protected memory space. Usually, the purpose of reserving protected memory is to provide space for machine-language programs. When using machine-language programs any BASIC programming used in conjunction with them is generally rather short and unsophisticated. In fact, a purely machine-language program can occupy the entire user's memory space, because there will be none of the memory-gobbling BASIC operations involved.

A BASIC program will never run into any protected memory for the simple reason that the built-in BASIC entry and execution routines never specify addresses larger than the upper string-space address entered in response to MEMORY SIZE?. Of course you can write BASIC programs that POKE into the protected memory space, and we will be doing that in a later section of this chapter. But the BASIC routines, themselves, never get into that upper-address area—if it has been specified.

## THE I/O BUFFER

In all Level II systems the i/o buffer occupies RAM addresses 16870 through 17127. All keyboard-entry operations take place through this memory space. Whenever you are typing in a line of BASIC programming, for example, it first goes into that buffer, and it is transferred to the program text space only after you strike the ENTER key. During the execution of a BASIC program the keyboard response to INPUT, INKEY\$, and any other kinds of in-line entry operations go to the buffer first.

Strings and lines of BASIC text are limited to 255 characters because they must pass through the i/o buffer first; the buffer isn't much larger than that.

If you want to take a look at the sort of data that goes into the i/o buffer, try this program from the command mode (do not use a line number):

```
A=16870:FOR I=0 TO 63:PRINT PEEK(A+I);NEXT I
```

As you type in this line of text the information goes into the i/o buffer; and when you end the operation by striking the ENTER key

**Chart 4-1. I/O Buffer Contents**

65	213	49	54	56	55	48	58	129	32	73	213	48	32	189
32	54	51	58	178	32	229	40	65	205	73	41	59	58	135
32	73	0	0											

the program displays itself as it appears in the i/o buffer. The resulting display resembles the sequence of numbers in Chart 4-1.

There will be more to your display than shown here, but, as you will see shortly, anything following a pair of zeros in succession is irrelevant—it's just bits and pieces of data left over from previous i/o operations.

If you compare that string of numbers in Chart 4-1 with an ASCII character table and the command-level program you entered, you can see something of a pattern in the whole thing.

The first number in the i/o listing, for example, is 65. It is no mere coincidence that 65 is the ASCII code for the letter *A* and that the first character in the program is an *A*. Skip the second number, 213, for a moment, and look at the third number in the i/o printout: it's a 49. This is the ASCII code for numeral *1*—the next alphanumeric character in the program.

Running through this process on your own, you will find the i/o buffer containing the ASCII code numbers for all variable names, constants, and punctuation in the program.

But command words and operators are not represented in an ASCII form in the i/o buffer. Rather, they are shown as special *compression code* numbers. The equal sign, for example, appears as a 213 everywhere in the i/o buffer. The FOR part of the FOR...NEXT command shows up as a 129.

Variable names and values appear in the i/o buffer in a literal, ASCII code format.

Command words and arithmetic/logic operators appear in the i/o buffer as single-number *compression codes*.

Of course, the point of loading command words and operators as single-number compression codes is to make more room in the i/o buffer (and, ultimately, the program memory space) for more BASIC operations.

Table 4-2 is a complete listing of command words and operators that have special compression codes. This table will help you "disassemble" the i/o listing in Chart 4-1. Now you should be able to see how that one-line program appears in the buffer.

**Table 4-2. Compression Codes for BASIC Commands and Operators**

Command	Compression Code	Command	Compression Code	Command	Compression Code
A\$S	217	INKEY\$	201	RND	222
AND	210	INP	219	RSET	172
ASC	246	INPUT	137	RUN	142
ATN	228	INSTR	197	SAVE	173
AUTO	183	INT	216	SET	131
CD\$	241	KILL	170	SGN	7
CHR\$	247	LEFT\$	248	SIN	226
CINT	239	LEN	243	SQR	205
CLEAR	184	LET	140	STEP	204
CLOAD	185	LINE	156	STOP	148
CLOSE	166	LIST	181	STR\$	244
CLS	132	LOAD	167	STRING\$	196
CMD	133	LOC	234	SYSTEM	174
CONT	179	LOF	235	TAB(	188
COS	225	LOG	223	TAN	227
CSAVE	186	LPRINT	175	THEN	202
CSNG	240	LSET	171	TIME\$	199
CVD	232	MEM	200	TO	189
CVI	230	MERGE	168	TROFF	151
CVS	231	MID\$	250	TRON	150
DATA	136	MKD\$	238	USING	191
DEF	189	MKI\$	236	USR	193
DEFD\$	155	MKS\$	237	VAL	255
DEFINT	153	NAME	169	VARPTR	192
DEFSNG	154	NEW	187	+	205
DEFSTR	152	NEXT	135	-	206
DELETE	182	NOT	203	*	207
DIM	138	ON	161	/	208
EDIT	157	OPEN	162	[	209
ELSE	149	OR	211	>	212
END	128	OUT	160	=	213
EOF	233	PEEK	229	<	214
ERL	194	POINT	198	&	38
ERR	195	POKE	177	'	251
ERROR	158	POS	220		
EXP	224	PRINT	178		
FIELD	163	PUT	165		
FIX	242	RANDOM	134		
FN	190	READ	139		
FOR	129	REM	147		
FRE	218	RESET	130		
GET	164	RESTORE	144		
GOSUB	145	RESUME	159		
GOTO	141	RETURN	146		
IF	143	RIGHT\$	249		

## "DISASSEMBLING" A BASIC PROGRAM

When you are writing in a line of BASIC program text the information first goes into the i/o buffer, and it accumulates there until you strike the ENTER key. As explained in the previous discussion the text in the buffer shows ASCII character codes for all numbers and variable names, but special compression codes for the commands and operators.

As shown in the i/o "disassembling" program information in Chart 4-1, the line ends with a double-zero marker. Those zeros are added automatically to the end of a line of program text in the i/o buffer, and they carry over to the program memory as the line is transferred there.

Here is a program that makes the point:

```
10 A=17129:FOR I=0 TO 63:PRINT PEEK(A+I):NEXT I
```

This program is very similar to the one used for looking at the contents of the i/o buffer. There are just two differences. First, this program is entered with a preceding line number. That line number has to be present in order to make the transfer from the i/o buffer to the program text space in the user's memory. Second, the starting location for the searching operation is different. Instead of starting at the beginning of the i/o buffer (A=16870), this one starts at the beginning of the BASIC program text space (A=17129).

Do a RUN on this program, and you will see the contents of the first 64 locations in the BASIC program text memory. It should look like the data in Chart 4-2.

**Chart 4-2. Contents of BASIC Program Text Memory Space**

14	67	10	0	65	213	49	54	56	55	48	58	129	32	73
213	48	32	189	32	54	51	58	178	32	229	40	65	205	
73	41	59	58	135	32	73	0	0						

Again, there will be more text on the screen than shown in this table, but anything past the double-zero point is irrelevant. Now compare Charts 4-1 and 4-2. They represent the same program, but the former is pulled from the i/o buffer and the latter as it appears after being transferred to the BASIC program text space.

The only significant difference between those two printouts is the presence of 4 extra bytes of information at the beginning of the program-text version in Chart 4-2. Skip those first four numbers, and the printouts are identical.

Indeed, the data set up in the i/o buffer is transferred to the BASIC text memory as ASCII characters and BASIC compression codes.



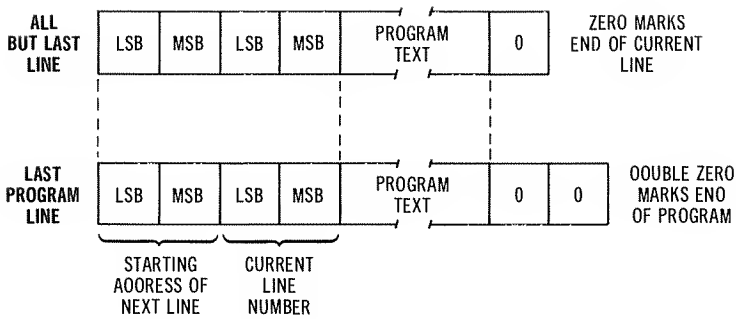


Fig. 4-2. BASIC text format.

What are those first 4 bytes in the program text listing? Well, they represent the program line number and the memory address of the next program line. Every line of BASIC text stored in the program memory starts out this way: 2 bytes for the address of the beginning of the next program line, and 2 bytes indicating the line number of the current line. See Fig. 4-2.

Interpreting the information in these first 4 bytes of each program line is a bit messy, because they are formatted as decimal versions of 2-byte numbers. So if you are confused by the discussion that follows, you should consult Appendix A to see the rationale behind it.

The first 2 bytes in Chart 4-2 indicate the address of the next line of text in the program. Those numbers are 14, 67. Converting them to a hexadecimal format, they are 0E, 43. Turning them around so that the msb is in its loading position, the hexadecimal address of the next program line is 430E. If you don't like working in hex, this number converts to decimal 17166. So whenever a second line is added to this program, it will begin at user's memory address 17166.

Any BASIC program begins at address 17129, and the line under discussion happens to end at 17165 (remembering that 17166 is the starting address of the next line). If you take the difference between the two numbers and add 1, you end up with the number of bytes that program line occupies in program memory: 37 bytes. Count the number of bytes in the listing of Chart 4-2, remembering to exclude the second of the two zeros marking its end. Sure enough, there are 37 bytes in that line.

Of course the number of bytes in a BASIC program line varies with the number of characters and commands you put into it. This one just happens to use 37 bytes of BASIC program memory.

Perhaps the discussion is getting a bit off the track here. The main point is that the first 2 bytes of any line of BASIC text that

is stored in program memory represent the starting address of the next line—the actual memory address.

The second 2 bytes reflect the current line number, as indicated in Fig. 4-2. In this particular case the line number is 10, and that shows up clearly in the third and fourth bytes of the listing in Chart 4-2.

A program made up of more than one line of text, as most BASIC programs are, will have a zero at the end of each one. The zero marks the end of a program line, and that's what the BASIC compiler looks for when executing a BASIC program. The last line of program text also ends with a zero, but there is a second zero to mark the end of the program. And that's how the BASIC interpreter knows it has reached the end of a program, whether the END statement is used or not.

As noted in an earlier demonstration, every line of text from the i/o buffer concludes with that double-zero, end-of-program marker. In a sense the i/o buffer assumes that every line you ENTER is the last one. But as you ENTER additional lines of program text, the line just ENTERED begins at the address of the second zero attached to the end of the previously entered line. So no matter how long or short your BASIC program might be, single zeros separate the lines of text in the program memory, and the last-entered line of text ends with two zeros.

Here is a program that searches the user's memory and displays the data for each program line. It ignores any zeros appearing in the first 4 bytes of each line but begins a new line of printout after seeing a single zero at the end of a line. The program stops under two conditions: whenever it sees a double-zero combination that marks the end of all the programming in the BASIC program text memory, and when it sees an END command—compression code 128.

If this is the only programming in the BASIC program text memory, it "disassembles" itself for you, showing you all the BASIC

```
1000 REM ** BASIC 'DISASSEMBLER' **
1010 CLS
1020 A=17128
1030 FOR B=1 TO 4
1040 A=A+1
1050 PRINT PEEK(A);
1060 NEXT B
1070 A=A+1
1080 PRINT PEEK(A);
1090 IF PEEK(A)=0 AND PEEK(A+1)=0 THEN 1150
1100 IF PEEK(A)=0 THEN 1140
1110 IF PEEK(A)=128 THEN 1130
1120 GOTO 1070
1130 A=A+1:PRINT PEEK(A):INPUT S$:GOTO 1030
1140 PRINT:GOTO 1030
1150 PRINT PEEK(A+1)
```

text information for each line. But if you enter a short BASIC program, using line numbers smaller than 1000, you can "disassemble" that program by doing a RUN 1000. Just be sure to end your custom program with an END statement. Otherwise the program will run through itself.

This discussion of how the BASIC program is formatted in the user's memory space ought to suggest some utility programs. You can, for instance, write a program that will search another program, looking for the occurrence of certain functions or variable names. The program can list the line numbers for you or count the number of occurrences.

Probably the most sophisticated and useful utility program is one that renumbers the BASIC lines for you. While such a program is beyond the scope of this book, you might be able to see how it would work. You can write a utility that will search out the line numbers (the second 2 bytes from the beginning of each line) and modify them in any way you choose. The tricky part of a renumbering utility is keeping track of line-number changes for GOTO, GOSUB, ON...GOTO, and ON...GOSUB statements. Finding such statements is easy; just look for their compression codes (Table 4-2), but keeping track of how the line numbers should change is another matter.

In any event, bear in mind that a BASIC program always begins loading at address 17129 and builds upward from there.

## PROTECTING MEMORY SPACE

The only section of the user's memory space that is available for any sort of non-BASIC tinkering is the optional protected memory. The amount of protected memory depends on your response to the system's MEMORY SIZE? query at power-up. The number ENTERed at that time sets the lowest available address for the protected memory, and it runs from that point to the highest memory address available on your system. See Table 4-1 for those top addresses.

As an example, suppose you want to save 1024 bytes (1K) for protected memory. If you have a 16K system, subtracting 1024 from 32767 yields the lowest address of the protected memory: 31743. Respond to the MEMORY SIZE? query by entering 31743, and you will have 1024 bytes reserved for your own memory operations; a space running from address 31743 to 32767.

There are occasions when it is necessary to save more than one block of protected memory. A user, for instance, might want 512 bytes for one kind of job, 64 for another, and maybe 1024 for yet another application. The MEMORY SIZE? query allows only one

$$SA = MA - BS$$

where

SA = the starting address of the protected memory space,

MA = the highest memory address available (Table 4-1),

BS = the number of bytes to be protected.

response. What will it be? Simply add up the byte sizes for the individual blocks, and subtract the results from the highest available address. In this case the total amount of protected memory required is  $512 + 64 + 1024 = 1600$  bytes. If the system has a 16K memory,  $32767 - 1600 = 32167$  gives the starting address of the protected memory space. That's the number to ENTER in response to MEMORY SIZE?

The same situation sometimes arises in a different light. The writer's line-printer system, for example, operates from a machine-language driver program. That program must be entered into the system before it will respond to LPRINT and LLIST commands. Now, the literature supplied with that software driver calls for responding to MEMORY SIZE? by entering 32512. Apparently that is the starting address of the driver routine in the protected memory space.

But for the sake of doing some special experiments with protected memory, I have to set aside some additional space, say 1024 bytes. Remembering that I cannot use any addresses from 32512 and up, the additional protected memory cannot have addresses exceeding 32511—otherwise my memory tinkering might mess up the line-printer program. So with 32511 being the highest address to be used, I subtract the additional bytes of reserved memory (1024) from 32511. That turns up the number 31487. Thus I should respond to MEMORY SIZE? with 31487. That being the case, I can begin my special memory projects at address 31487, and as long as I do not tinker around in memory addresses greater than 32511, the line-printer driver program will remain intact.

Unfortunately, there is no way to protect blocks of memory within the protected memory area. The programmer simply has to be very careful about the matter, perhaps writing BASIC programs that do not allow things to be POKEd into certain blocks of protected memory.

### **SOME SPECIAL MEMORY OPERATIONS**

Most TRS-80 literature refers to protected memory in the context of machine-language programs: specifically, machine-language pro-

grams that are used in conjunction with BASIC programs. This is a common situation and this book deals with it later on, but it isn't the only reason for reserving some protected memory. Protected memory can be used as a repository for information that has nothing at all to do with machine-language programming. And that is the point of the following discussions.

### **Saving One-Line Text in Memory**

The following BASIC program, PUTTING STUFF IN, lets you type some text on the screen and save it in memory. Every character is loaded into memory as you enter it, and the entering operation ends only as you strike the pound-sign key (#).

The program assumes that the characters will be entered into successive memory locations, beginning at 31487. This means you have to save some protected memory by answering MEMORY SIZE? with that number. Of course, you are free to change that starting point of the protected memory; but if you do, be sure to change the initial value of variable A in line 1010 accordingly.

The program further assumes you will not wish to enter more than 1000 characters. That limit is set by IM=1000 in line 1010. That, too, can be changed if you wish.

```
1000 REM ** PUTTING STUFF IN **
1010 A=31487:I=0:IM=1000
1020 CLS:PRINT CHR$(14);
1030 C$=INKEY$:IF C$="" THEN 1030
1040 IF I>IM THEN 1110
1050 PRINT C$;POKE A+I,ASC(C$)
1060 IF C$="#" THEN 1090
1070 I=I+1
1080 GOTO 1030
1090 PRINT:PRINT "TEXT ENTERED"
1100 END
1110 PRINT:PRINT "MEMORY FULL"
1120 END
```

So, get to the MEMORY SIZE? query by turning the TRS-80 off and on, and respond by entering the beginning address of some protected memory—31487 in this example. Then do a RUN.

You will see the cursor on the screen, because the PRINT statement in line 1020 turns it on for you. Then type in some text, anything you want to enter into the protected memory space. When you are done, strike the # key.

Striking the # key signals the end of the test-writing operation, and the program will confirm that condition by printing TEXT ENTERED.

Looking through some of the other lines, note how the current character, carried as string C\$, is both PRINTed on the screen and

POKEd into protected memory at line 1050. Variable I, initialized at zero in line 1010, is the character counter. It is incremented after every character operation at line 1070.

If the value of I should ever exceed IM, the program brings up a MEMORY FULL message and stops everything. The value of IM (the maximum number of characters to be written into the protected memory space) is set to 1000 in line 1010.

So this program, PUTTING STUFF IN, gets up to 1000 characters loaded into successive memory locations in the protected memory space. Just be sure to end the operation by striking the # key. Now the problem is to get to that saved text material and print it out on the screen again. This will confirm that the program works.

The next program, GETTING STUFF OUT, runs through the protected memory, printing out saved characters until it sees the pound-sign character.

```
2000 REM ** GETTING STUFF OUT **
2010 A=31487:I=0:IM=1000
2020 C$=CHR$(PEEK(A+I))
2030 IF C$="#" THEN END
2040 PRINT C$;
2050 I=I+1
2060 IF I>IM THEN 2080
2070 GOTO 2020
2080 PRINT:PRINT "OUT OF MEMORY"
2090 END
```

On doing a RUN 2000, *GETTING STUFF OUT* PEEKs at successive addresses in the protected memory (line 2020) and prints the characters contained therein (line 2040). The reading and printing operation continues until it turns up the pound-sign character (line 2030) at the end of the text or it runs out of memory (lines 2060 and 2080).

With those two programs entered into your system, check out the operation of the scheme by writing in some text, making sure you don't strike the # key until you want to end the text-entering phase. Then do a RUN 2000 to see the message you saved in protected memory.

It is possible to slick up the operations a bit, using PUTTING STUFF IN and GETTING STUFF OUT as subroutines for a nice mainline program. Here, then, is 1-LINE TEXT MEMORY DEMO:

```
100 REM ** 1-LINE TEXT MEMORY DEMO **
110 CLS
120 INPUT "READ OR WRITE";F$
130 IF LEFT$(F$,1)="W" THEN 150
140 IF LEFT$(F$,1)<>"R" THEN 120 ELSE 170
150 GOSUB 1000
```

```

160 GOTO 120
170 CLS:GOSUB 2000
180 END
190 REM
200 REM
1000 REM ** PUTTING STUFF IN **
1010 A=31487:I=0:IM=1000
1020 CLS:PRINT CHR$(14);
1030 C$=INKEY$:IF C$="" THEN 1030
1040 IF I>IM THEN 1110
1050 PRINT C$;POKE A+I,ASC(C$)
1060 IF C$="#" THEN 1090
1070 I=I+1
1080 GOTO 1030
1090 PRINT:PRINT "TEXT ENTERED"
1100 RETURN
1110 PRINT:PRINT "MEMORY FULL"
1120 END
1130 REM
1140 REM
2000 REM ** GETTING STUFF OUT **
2010 A=31487:I=0:IM=1000
2020 C$=CHR$(PEEK(A+I))
2030 IF C$="#" THEN END
2040 PRINT C$;
2050 I=I+1
2060 IF I>IM THEN 2080
2070 GOTO 2020
2080 PRINT:PRINT "OUT OF MEMORY"
2090 END

```

If PUTTING STUFF IN and GETTING STUFF OUT are already resident in your BASIC memory space, all you have to do is enter the mainline program (lines 100 through 200) and change line 1100 to RETURN instead of END. The program gives the option of reading or writing into the protected memory.

If you choose to write, respond to the READ OR WRITE message with a W and enter your text, ending with a pound sign. The program immediately returns to the READ OR WRITE message and, at that point, if you want to see a printout of the text, respond with a R. After printing out the message saved in the protected memory space, the program ends.

Incidentally, the text remains in the protected memory space, even after the BASIC program ends. You can confirm that fact by doing another RUN and specifying a READ operation right away.

### **Saving a Whole Screen of Text/Graphics in Memory**

It isn't always easy to compose elaborate graphics on the TRS-80 screen. The SET/RESET feature eases the task a bit, however, by saving you the trouble of building graphics out of the graphic

character set and doing the appropriate sequence of POKE statements.

But no matter how you choose to go about composing a graphic, it is possible to dump the finished product into some user's memory space and call it back to the screen, rather quickly, whenever you want.

The following demonstration uses some SET statements to draw the face with a "have-a-nice-day" smile on the screen. As soon as the drawing operation is done, the entire screen—the video memory—is transferred to user's memory space, beginning at address 19454. After that, this message appears on the screen: GRAPHIC IS DRAWN AND SAVED/ ENTER X TO GET IT FROM MEMORY. When you respond to the message by entering an X, the graphic is duplicated from user's memory to the video memory. Note that this transfer takes place quite a bit faster than the original drawing operation.

Once the graphic is drawn and saved (program lines 20 through 180) you can recall it to the screen at any later time by doing a RUN 200. This assumes, of course, you have done nothing that will disturb the contents of the memory space where it is saved.

```
10 REM *** SCREEN TRANSFER DEMO ***
20 CLS
30 FOR A=0.02 TO 6.28 STEP 0.02
40 X=COS(A):Y=SIN(A)
50 SET T(34*X+64,16*Y+23);SET(4*X+48,2*Y+18);SET(4*X+80,2*Y+18)
60 NEXT A
70 FOR A=0.02 TO 3.14 STEP 0.02
80 SET T(16*COS(A)+64,8*SIN(A)+25)
90 NEXT A
100 V=15360:M=19454
110 FOR I=0 TO 1024
120 POKE M+I,PEEK(V+I)
130 NEXT I
140 CLS
150 PRINT "GRAPHIC IS DRAWN AND SAVED"
160 PRINT "ENTER X TO GET IT FROM MEMORY"
170 INPUT S$
180 IF S$="X" THEN 200 ELSE 170
200 V=15360:M=19454
210 CLS
220 FOR I=0 TO 1024
230 POKE V+I,PEEK(M+I)
240 NEXT I
```

You can replace program lines 10 through 90 with any text or graphic-drawing routine you choose. Just work things out so that your program goes to line 100 when it is time to transfer the screen material to user's memory.



Of course, you can change the location of the graphic storage space by changing the value of the M variable in lines 100 and 200. This variable represents the starting address for the storage area, and the routine occupies 1024 bytes from there.

Then, too, you can change the line numbers if you find you do not have enough space for your drawing routine. Just bear in mind that lines 100 through 180 transfer the screen information to user's memory, and lines 200 through 240 bring it back.

This is not really the most efficient sort of graphic-storage technique; a better one would use some space compression techniques to eliminate the storage of long strings of spaces. But this technique does illustrate one powerful use for user's RAM space.

## CHAPTER 5

# Linking BASIC and Machine Language With USR

Any protected memory can be used as a repository for data—data being used in conjunction with some BASIC programming. This notion was described and demonstrated in the latter part of Chapter 4.

But that same sort of protected memory can also be used as a place for storing machine-language programs. And those programs, like the data that could be stored there as well, can be used in conjunction with some BASIC programming. This is the main topic of this chapter.

Using BASIC and machine-language programs together can be useful at times and, certainly, instructive at any time. The idea is especially useful when you are writing special utility programs that manipulate the character of a resident BASIC program; and when it comes to writing BASIC programs that call for some high-speed graphics, a machine-language subroutine can be invaluable.

As far as being instructive, the idea of using BASIC and machine-language at the same time can be a useful learning tool, especially in cases where the user is fairly well acquainted with BASIC but knows little about machine-language programming.

When using BASIC and machine language at the same time, the BASIC programming always resides at the lower end of the user's memory space. Not much can be done about changing that position for the BASIC, because it is automatically established by some TRS-80 ROM programming. (See the details of this particular matter in Chapter 4.)

The machine-language part of the scheme always resides in a portion of user's memory that is above all the BASIC—in some protected memory space. You, the user, must establish the size, or

starting address, of that protected memory by making an appropriate response to the MEMORY SIZE? query. Again, the details for this operation are outlined for you in Chapter 4.

In a practical sense, then, the first thing you must always do when combining BASIC and machine-language programming is to set aside some protected memory space for the machine-language part. And you must keep track of the number you specify in response to the MEMORY SIZE? query, because it marks the lowest possible address you can use for storing the machine-language program.

Suppose the BASIC is automatically set to the lower part of the user's memory, and you establish the starting point for the machine-language programming in the upper part of the user's memory. Now, how do you get a BASIC program to call the machine-language program, and vice versa? Answer: Use the BASIC USR statement.

The USR statement is your link between BASIC and any machine-language programming. Actually, the USR statement works much like a GOSUB; whenever the computer is running a BASIC program and encounters a USR, operations are sent to a specified starting address for a machine-language program. Just how that starting address is specified will be discussed in a moment.

So a USR statement gets operations from BASIC to machine language. But how, if you want, do you get back from a machine-language program to BASIC? How, if USR works like a GOSUB, do you carry out the RETURN part of the call?

Returning from a machine-language program to BASIC is a matter of ending the machine-language subroutine with a machine-language version of a RETURN statement. There are several different kinds of machine-language RETURN instructions, and it is up to you, the machine-language programmer, to select the right one for the job at hand. But the point is that a machine-language version of a RETURN sends control back to the BASIC line that follows the USR statement.

### **GETTING SET UP FOR A USR OPERATION**

While the TRS-80 system automatically initializes a great many different operations, it does not initialize the USR function. You have to prepare the way yourself.

First, you must set aside some useful amount of protected memory for the machine-language programming. As described in Chapter 4, that is a matter of responding to the MEMORY SIZE? query with an address that marks the beginning of the protected memory space.

Keep track of that number, because it represents the lowest address you can call from BASIC.

The next step is to prepare some critical parts of the BASIC

routine. Schemes that combine BASIC and machine language always begin operation in BASIC; and one of the critical BASIC operations is to establish the starting address of the machine-language subroutine. USR works like a GOSUB, but, just as you must specify a line number for a GOSUB statement, you must also specify a starting address for a USR-called machine-language subroutine.

Now, that starting address can be the lowest-available address in protected memory—the same address you entered in response to the MEMORY SIZE? query. But it can be any address within the protected memory space.

That starting address is specified as a 2-byte decimal number that is POKEd into addresses 16526 and 16527. The lsb of the 2-byte version of the starting address goes into 16526 and the msb goes into 16527. Suppose, for example, you want to begin a BASIC-called machine-language subroutine at 31404. Converting that decimal number to a 2-byte decimal format (as described in Appendix A), it comes out to be 122 172—msb and lsb, respectively.

To call a machine-language subroutine that begins at address 31404, the BASIC program must include the following:

```
POKE 16526,172;POKE 16527,122
```

Whenever the computer encounters a USR statement, it first looks to addresses 16526 and 16527 to find out where the machine-language routine is to begin. Those two “magic numbers” never change. It is up to you, however, to specify the 2-byte version of the starting address in those memory locations.

To initialize a machine-language subroutine from BASIC, you must:

```
POKE 16526,lsb;POKE 16527,msb
```

where *lsb* and *msb* are the least significant and most significant bytes of a 2-byte version of the starting address.

Anytime after specifying the starting address of the machine-language subroutine, you are free to write in a USR statement to call it. POKing the starting address does not actually call the subroutine; the USR statement does that.

Most programmers who use BASIC and machine language together specify the starting address—with the two POKE statements—rather early in the BASIC program. That way they can be sure the subroutine is properly initialized when they are ready to write in the USR statement.

So far in this discussion you have seen how to set aside some protected memory space for a machine-language program and then specify the starting address of a machine-language subroutine. The final step in the operation is to specify the USR statement itself.

The expression USR is not a complete statement, but it must be presented as a complete statement if the BASIC program is to execute it—to call the machine-language subroutine you have specified.

A complete USR statement can take on several forms, depending on what you want it to do:

1. Call a machine-language subroutine and return to BASIC without passing any values in either direction.
2. Carry a value from BASIC to the machine-language subroutine but return to BASIC without passing a value.
3. Call a machine-language subroutine without passing a value but return to BASIC, carrying a value.

There are several variations that combine a couple of these basic kinds of USR-oriented operations, but these are the fundamental ones.

In the first case, passing no values between the two kinds of programs, the USR in BASIC simply calls a machine-language subroutine that does a job that does not generate any numbers that are useful to the BASIC program. Perhaps the machine-language program generates some high-speed animation graphics and does nothing to generate any numbers that are useful to later BASIC operations. The USR statement that calls the subroutine is likewise lacking in any numerical information that is useful to the animation sequence.

Calling a machine-language subroutine, without passing any values, can be done with a statement such as

```
X=USR(X)
```

where the Xs are dummy variables—variables arbitrarily chosen simply to fit the USR function into a complete BASIC statement. A statement such as A=USR(B) would accomplish exactly the same thing.

Assuming that a machine-language program is already loaded into protected memory, beginning at address 31488, a typical BASIC sequence could look like this:

```
10 POKE 16526,0:POKE 16527,123
20 CLS
30 X=USR(X)
40 PRINT "IT'S DONE"
50 END
```

Line 10 in this example establishes the calling address at 31488, line 20 simply clears the screen, and the USR statement in line 30 calls the machine-language program. And when the system is through with the machine-language subroutine (and assuming that subroutine ends with a RETURN-like instruction), operations return to BASIC at line 40, printing the IT'S DONE message and coming to an END at line 50.

No numerical values are passed from BASIC to the machine-language subroutine, and none return to BASIC when the subroutine is done.

The second major class of USR calls passes some value from BASIC to the machine-language subroutine. Maybe you would like the machine-language subroutine to print some characters on the screen, with those characters being specified from the keyboard while the BASIC part of the program is still running. That is just one sort of operation that calls for passing a variable value from BASIC to machine language.

Calling a machine-language subroutine, and passing a value to it from BASIC, can be done with a statement such as:

```
X=USR(Y)
```

where X is still an arbitrarily chosen dummy variable, but Y is some integer between -32,768 and 32,767. The dummy variable, X, is necessary only to fit the USR function into a complete BASIC statement. A statement such as X=USR(65), for example, will carry decimal number 65 to the machine-language program, where it will be used for some particular purpose.

Here is a program that passes an ASCII character code, in decimal, to a machine-language program:

```
10 POKE 16526,0:POKE 16527,123
20 CLS
30 INPUT S$
40 S=ASC(S$)
50 X=USR(S)
60 PRINT "DONE"
70 GOTO 30
```

The program assumes that some sort of machine-language program is in protected memory and begins at address 31488. That starting address is specified by the POKE statements in line 10. Line 20 clears the screen, and line 30 allows you to input some value for S\$ from the keyboard. Line 40 assigns the ASCII version of that variable to S, and line 50 calls the machine-language subroutine, passing the value of S to it. When the system is through executing the machine-language subroutine, operations return to line 50 of the

BASIC program, PRINTing a DONE message and looping back up to line 30, giving you a chance to pass another ASCII character code to the subroutine.

The argument in the USR statement—the term enclosed in parentheses—can be a function as well as an integer value. Taking advantage of this feature, the previous example can be shortened somewhat: omit line 40, and replace line 50 with `X=USR(ASC(S$))`.

Recall that any value passed from the USR to the machine-language program must be an integer value. If you are doing some mathematical operations in BASIC that might pass a value that isn't an integer, you can goofproof the system by doing something such as `X=USR(INT(Y))`, where Y can be any value—integer or otherwise—between -32,768 and 32,767.

Finally, you can always pick up a numerical value from a machine-language program and pass it back to the BASIC programming. Perhaps your machine-language programming does some math or logic operations for you, and you want to return the results to BASIC.

Calling a machine-language program and bringing back a value when it returns to BASIC can be done with a statement such as

```
Y=USR(X)
```

where X is a dummy variable used to make a complete BASIC statement, and Y is the value returned from the machine-language program. That value returned from machine language will be some integer between -32,768 and 32,767.

Here is a program that picks up a value from a machine-language program that it calls:

```
10 POKE 16526,0:POKE 16527,123
20 CLS
30 Y=USR(X)
40 PRINT Y
50 END
```

This one assumes a machine-language program begins at address 31488, a point specified by the POKes in line 10. Line 20 clears the screen, and line 30 calls the machine-language subroutine. When the system returns from that subroutine, line 30 has called for assigning the numerical result to variable Y. Line 40 then prints that value for you.

It is possible to shorten the BASIC program in this case by omitting line 40 and rewriting line 30 this way:

```
30 PRINT USR(X)
```

This modification bypasses the need for variable Y, and it illustrates the fact that a USR function can be used in more than one sort of complete BASIC statement. The USR function *is* a function, and it can be used as such—just as ABS, INT, SQR, and a whole host of BASIC functions are used.

This completes the discussion of the three fundamental USR operations, at least from the BASIC point of view. Much of the material remaining in this chapter considers the same sorts of things from the machine-language viewpoint.

Before getting down to the machine-language view of things, consider an important variation of two kinds of USR operations: combining those passing a value from BASIC with those passing a value back from the machine-language program. Such a scheme can look like this in BASIC:

```
10 POKE 16526,0:POKE 16527,123
20 INPUT S$
30 PRINT USR(ASC(S$))
40 GOTO 20
```

Assuming once more that a machine-language program is resident and beginning at address 31488 (as specified by line 10), the BASIC program picks up a value for S\$ from the keyboard at line 20. Line 30 then passes the ASCII value of that string to the machine-language program, and after the machine-language program has presumably done some manipulations with that value, line 30 also brings back the result to BASIC and prints it on the screen.

So line 30 actually does three things: it calls a machine-language subroutine, passes an ASCII version of S\$ to it, and then brings back the result from machine language, printing it on the video display.

### **A PRELIMINARY NOTE ABOUT MACHINE-LANGUAGE PROGRAMMING**

Machine-language instructions are stored in memory as 1-byte binary numbers. Every instruction that the Z-80 microprocessor (the one used in your TRS-80) can perform is represented by one of those 1-byte *op codes*.

Appendix B is a complete listing of the Z-80 instruction set, including the op codes and almost-plain-English source-code mnemonics. Every machine-language instruction used in this book can be found somewhere in that listing.

So you can find the complete instruction set for Z-80 machine-language programming in this book, but you will find no details concerning the actions of these instructions nor any special dis-



cussions of fundamental machine-language programming. All of that sort of thing is beyond the scope of this book. Indeed, it is unfortunate that it is impractical to write a single volume that both spells out the special features of the TRS-80 system and teaches beginners how to handle machine-language programming. It is simply too much for one book.

The emphasis of this book is on *what* can be done with your Level II TRS-80. The *how* of the matter is described in general terms and does not (cannot) get down to the finer details of machine-language programming.

You can certainly work the examples and demonstration programs without having any knowledge of machine-language programming, but you have to learn that sort of programming on your own if you want to extend the examples and demonstrations to suit your needs.

Any good book on the Z-80 microprocessor and its instruction set will do. Learn that sort of thing from another source, and this book will help you fit it into the context of your TRS-80 machine.

### **POKEING IN MACHINE-LANGUAGE PROGRAMS FROM BASIC**

When linking BASIC and machine-language programs with the USR function, the BASIC program generally performs two functions. It calls the machine-language program as outlined in several different ways earlier in this chapter. But it also loads the machine-language program into protected memory for you.

Thus there are at least two phases to the BASIC part of the scheme: loading the machine-language program and calling it when the loading is done.

After you have written a machine-language program on paper it must be POKEd into its memory space, one byte at a time and in the proper sequence.

The simplest possible machine-language program in a USR-linked scheme is one that does nothing more than return operation to the BASIC program. The instruction code for an unconditional RETurn is decimal 201. So if a 201 is POKEd into the starting address of the machine-language program, a USR function will call it from BASIC, and it will return control immediately back to the BASIC program. Here is a complete example that you can run on your machine:

```
10 POKE 16526,0:POKE 16527,123
20 POKE 31488,201
30 X=USR(X)
40 CLS:PRINT "IT WORKED!!"
50 END
```

It isn't a very exciting program, but it serves as an example of the discussion at hand.

Line 10 establishes the starting point of the machine-language program at 31488. The 0 and 123, you recall, are 2-byte versions of the starting address. The purpose of that line is to tell the USR function where it is supposed to begin execution of the machine-language program.

Line 20 represents the main point of this discussion. That statement POKes the one-instruction machine-language program into memory address 31488. The 201 is a decimal version of the Z-80's unconditional RETurn instruction.

Now, it is no coincidence that line 10 sets the starting point at address 31488 and line 20 pokes the first (and only) machine-language instruction into that same address. The first instruction and the starting point of the USR-called subroutine must be the same.

By the time the system completes its execution of line 20 the BASIC program has loaded the machine-language program. The next step is to execute that machine-language program, and that is done by the USR statement in line 30.

It should be apparent that the machine-language program must be entered before it is called by the USR function. The statement in line 20 must occur sometime before the statement in line 30.

The statement in line 10—establishing the calling address for the USR function—must occur before USR does but not necessarily before the machine-language-loading operation. Line 10, for example, could be moved to, say, line 25, and things would work just as well. It is a fairly well established convention, however, to specify the starting address of the machine-language programming very early in the BASIC routine.

So line 20 calls the one-instruction machine-language program. And when control returns to BASIC, things pick up at line 40, clearing the screen and printing IT WORKED!!.

Incidentally, if you failed to POKE the right instruction into memory at line 20, the system might not return at all. It would most likely go "into outer space," making it necessary to do a RESTART (depressing the RESTART push button on the back of your TRS-80 unit) or a total power-up.

Seeing the IT WORKED!! message in this case confirms the proper function of that one-instruction subroutine.

When loading a machine-language program from BASIC POKE statements, the op codes must be entered in a decimal format. TRS-80 BASIC, you see, is a decimal-oriented language. This poses a minor difficulty, because the Z-80 instruction set specifies the op codes in a hexadecimal format. So an unconditional RET in Ap-

pendix B is shown as C9. This is a hexadecimal number. But the same instruction is represented as a 201 in that last example. Where does the 201 come from? It is the decimal version of hexadecimal C9.

The moral of the story is this: *hexadecimal op codes from the Z-80 instruction set must be converted to decimal before they can be entered into memory by way of POKE statements.* See Appendix A, Table A-1, to make the necessary conversion.

Most machine-language programs have more than one instruction in them—a lot more. Each one of them has to be loaded into memory by a POKE statement; if you have a lot of instructions, you can end up with a whole lot of POKE statements.

Things are generally simplified, however, by enclosing the instruction-POKEing statement within a FOR . . . NEXT loop. And if that FOR . . . NEXT loop also contains a READ statement, the instructions can be written in their proper sequence in some DATA lists.

Here is a working example:

```
10 POKE 16526,0:POKE 16527,123
30 FOR I=0 TO 5:READ D:POKE 31488+I,D:NEXT I
40 CLS
50 X=USR(X)
60 PRINT:PRINT
70 END
100 DATA 62,65,50,0,60,201
```

Line 10 sets the starting point of the machine-language program at 31488. This assumes, of course, you have previously set aside some protected memory that includes that address. The program is loaded into protected memory by the statements in line 30.

Line 30 includes a READ statement that references the DATA list in line 100. Line 100, incidentally, is the machine codes to be entered. In this case it is a 6-byte program. So as variable I is incremented from 0 to 5, the READ statement picks up the DATA items in sequence, and the POKE statement inserts them into successive memory locations, beginning at 31488—the place already specified as the starting point in line 10.

Once the system completes the execution of line 30 the machine-language program is completely loaded, and the next series of BASIC operations clears the screen, executes the machine-language program (line 50), does a couple of PRINTs, and then ENDS.

We'll discuss the machine-language program, itself, shortly. What is more important now is the way the BASIC program loads the machine-language program.

The statements shown here in line 30 have a form that is almost universally accepted as the general procedure for loading a machine-language program from BASIC. The features thus call for a detailed analysis.

The first part of the line reads, in general:

```
FOR i=0 TO n
```

where

*i* is any numerical variable name,

*n* is the number of machine-language bytes to be loaded, minus 1.

There are six items in the DATA list in the previous example, so the first part of the loading sequence should read: FOR I=0 TO 5. If there happened to be 25 items in the DATA list, you would write FOR I=0 to 24.

The second part of the loading sequence reads:

```
READ d
```

where *d* is any valid numerical variable name.

The third part looks like this, when presented in a general form:

```
POKE a+i,d
```

where

*a* is the starting address of the machine-language program,

*i* is the same variable used in the first phase of the loading operation,

*d* is the same variable READ in the previous statement.

The value of address *a* in this statement must be a conventional decimal version of the starting address specified in a 2-byte decimal format at the beginning of the program. In this case, address 31488 is the conventional decimal form of 123 0.

Putting it all together, the line for entering a machine-language program from BASIC looks like this:

```
FOR i=0 TO n:READ d:POKE a+i,d:NEXT i
```

where

*i* and *d* are any convenient numerical variable names,

*a* is the starting address of the machine-language program,

*n* is the number of bytes in the machine-language program, minus 1.

This entire scheme for loading a machine-language program from BASIC assumes you have listed the instructions and data as decimal numbers in a DATA list. In the previous example the DATA list read: 62,65,50,0,60,201. A disassembled version is shown in Table 5-1.

**Table 5-1. A Typical Decimal-Oriented Machine-Language Program**

Address	Op Code	Mnemonic	Comment
123 0	62 65	LD A,65D	;LOAD ASCII "A" INTO REGISTER A
123 2	50 0 60	LD (15360D),A	;LOAD A TO THE VIDEO MEMORY
123 5	201	RET	;RETURN TO BASIC

Table 5-1 is a formal, assembly-language version of the program executed by a USR function in line 50 of the previous BASIC program. It prints the letter A at the first video memory location.

The assembly-language program in that table has some features that are common to any sort of Z-80 source and op-code listing. One difference is that the op codes are presented in a decimal, rather than hexadecimal, format. The addresses are also in a decimal format—a 2-byte decimal format, to be exact. The justification for using addresses in this peculiar 2-byte decimal format will become clear in a later section of this chapter. The mnemonics are standard Z-80 mnemonics, and the comments are typical of those used with any machine-language programming.

If you are totally confused by the listing in Table 5-1, you will have to consult a Z-80 assembly-language book or manual (e.g. look at Radio Shack's *TRS-80 Assembly-Language Programming*).

When you get tired of seeing this program do nothing but print a single A in the upper left-hand corner of the screen, change the number in the second item of the DATA list to some other number between 33 and 191. Then run the program again.

That little exercise, incidentally, suggests how easy it is to modify a machine-language program that is loaded from BASIC; just change the item or items in the DATA list and, if necessary, adjust the value of *n* in the FOR *i*=0 TO *n* part of the loading routine.

### **SOME PROGRAMMING EXAMPLES**

The following examples of USR-called programs are intended to illustrate the range of interaction that is possible between a BASIC program and a machine-language program that it can call. In every instance the BASIC program not only controls the calling of a machine-language program but loads it into a specified section of protected memory as well. All examples assume you have responded to MEMORY SIZE? with a number no greater than 31488.

#### **A Simple Calling Routine From BASIC**

Here is a decimal-oriented, machine-language program that simply draws a wide white bar across the top line of the screen:

```

;WIDE BAR ROUTINE
123 0    33 0 60    START:    LD HL,15360D    ;POINT TO START OF VIDEO
123 3    62 63      LD A,63D      ;END OF COUNT TO A
123 5    54 191     AGAIN:    LD (HL),191D    ;SET GRAPHIC CHARACTER
123 7    1B9        CP L          ;LOOK FOR END
123 8    200        RET Z         ;RETURN IF AT END
123 9    35         INC HL        ;NEXT VIDEO POINT
123 10   195 5 123  JP AGAIN     ;DRAW AGAIN
;END

```

Basically, the program loads character code 191 (full rectangle) into the 64 character locations across the top of the screen. Here is a BASIC program for loading and executing it:

```

0 REM ** CALL A WHITE BAR **
10 POKE 16526,0:POKE 16527,123
20 FOR I=0 TO 12:READ D:POKE 3148B+I,D:NEXT I
30 CLS
40 X=USR(X)
50 PRINT:PRINT
60 END
100 DATA 33,0,60,62,63,54,191,1B9,200,35,195,5,123

```

The BASIC program loads the machine-language program from the DATA list in line 100, using the procedure in line 20. After that, it clears the screen (line 30), calls the machine-language program, and does a couple of PRINTs before coming to an END. Of course the white bar is drawn the moment the BASIC program calls the machine-language program via a USR statement. The RET Z instruction in the machine language ensures that the operations return to BASIC when the line-drawing operation is done.

### Conditioning a Call at the BASIC Level

In some actual applications you might not want to execute the machine-language portion of a program unless certain conditions are met at the BASIC level. Using the same machine-language program--the one that draws a white bar across the top of the screen--you can make the drawing of the line contingent on striking the B key. Here's the revised BASIC program:

```

0 REM ** CONDITIONAL CALL FROM BASIC **
10 POKE 16526,0:POKE 16527,123
20 FOR I=0 TO 12:READ D:POKE 3148B+I,D:NEXT I
25 C$=INKEY$:IF C$="" THEN 25
30 IF C$="B" THEN 40
35 PRINT "NOT YET":GOTO 25
40 CLS:X=USR(X)
50 PRINT:PRINT
60 END
100 DATA 33,0,60,62,63,54,191,1B9,200,35,195,5,123

```

This BASIC program still loads the machine-language subroutine. See lines 20 and 100. Doing the USR function, however, is contingent on striking the *B* key. Until that happens, striking any other key yields a NOT YET message.

Conditioning the call of a machine-language subroutine can be a rather sophisticated operation. For example, you might work out a BASIC program of this sort, allowing it to execute the USR function (and hence the machine-language program) only on typing the string CALL BAR, or something like that.

### Repetitive Calling

The following program draws a solid bar of light across the top of the screen. Everything remains unchanged until you depress the *F* key. At that time the bar begins flashing on and off, and it remains flashing until you release the *F* key.

The BASIC program is responsible for sensing the depression of the *F* key and calling a machine-language program that first turns off the light, and then turns it on again. After completing that one cycle, control returns to BASIC, where the *F* key is checked again. If that key is still depressed, the control returns to the machine-language program for another off/on cycle. But when control returns to BASIC and at that time the *F* key is no longer depressed, the program continues looping in BASIC until the *F* key is depressed again.

This is an example of looping around and around between a BASIC and machine-language program. Such loops have some powerful applications, especially when it comes to doing some real-time, animated graphics.

The flashing effect, by the way, is far less satisfactory if it is generated—if the bar is alternately erased and drawn—in purely BASIC terms. Even POKE graphics do not run fast enough to create a good flashing effect of such a long bar on the screen. Machine-language graphics certainly does the job, as you will see when running this program.

Here is the BASIC part of the program:

```
0 REM ** FLASHING BAR WITH BASIC CONTROL **
10 POKE 16526,0:POKE 16527,123
20 FOR I=0 TO 37:READ D:POKE 31488+I,D:NEXT I
30 CLS:PRINT @ 64,"DEPRESS THE F KEY TO GET A FLASHING BAR OF
   LIGHT..."
40 IF PEEK(14463)=0 THEN 40
50 C$=INKEY$
60 IF C$<>" " THEN R$=C$
70 IF R$<>"F" THEN 40
80 X=USR(X)
90 GOTO 40
```

```

100 DATA 6,128,205,11,123,6,191,205,11,123,201
100 DATA 33,0,60,62,63,112,189,202,25,123,35,195,16,123
120 DATA 22,225,14,225,13,194,29,123,21,194,27,123,201

```

Line 10 sets the starting address of the machine-language program at 31488, and line 20 loads that program from DATA lines 100, 110, and 120. That is a 38-byte program there.

Line 30 simply clears the screen and prints a prompting message at the second line on the screen.

The real looping action, the main point of the demonstration, begins at line 40. This line, working together with lines 50 and 60, views the keyboard as a big set of normally open push-button switches. Line 70, however, maintains the key-searching loop until you depress the *F* key.

Line 80 uses the USR function to call the machine-language subroutine—a routine that flashes the white bar off and on one time. After completing just one flashing cycle in machine language, control returns to BASIC at line 90.

Now, line 90 restarts the key-sensing loop; and if the *F* key is still depressed, line 80 is executed again, and the bar of light flashes. As long as that *F* key is depressed, the USR function and the machine-language subroutine become part of a loop. Line 80 is removed from the loop only when the *F* key is not depressed.

So, indeed, machine-language programs can be written into a BASIC program loop.

In passing, notice the short time delay between the time you do a RUN for the program and the appearance of the prompting message. That delay will occur only when you start this program from scratch, and it is caused by the time required for POKEing the machine-language op codes into your protected memory space.

The machine-language program, itself, is justified in its hand-assembled form in Table 5-2. The number of jumps and calls in the routine clearly demonstrates the value of specifying the addresses as 2-byte numbers. Note that only the lsb of those addresses changes. The same msb is good for up to 256 memory locations.

### PASSING A VALUE TO THE MACHINE-LANGUAGE SUBROUTINE

As stated earlier in this chapter the statement `X=USR(Y)` will pass the value of the argument, the *Y*, to the subroutine it calls, provided *Y* is an integer between -32,768 and 32,767. This is stating the case in a very general way, however. It is more precise to say that the `X=USR(Y)` statement makes the value of the argument *available* to the machine-language program it calls.



**Table 5-2. Machine-Language Program Called by the BASIC Program  
FLASHING BAR WITH BASIC CONTROL**

Address	Op Code	Mnemonic	Comment
123 0	6 128	START: LD B,128D	;SET B FOR OFF
123 2	205 11 123	CALL DRAW	;DRAW "OFF"
123 5	6 191	LD B,191D	;SET BE FOR ON
123 7	205 11 123	CALL DRAW	;DRAW "ON"
123 10	201	RET	;RETURN TO BASIC
123 11	33 0 60	DRAW: LD HL,15360D	;START OF VIDEO MEMORY
123 14	62 63	LD A,63D	;SET END IN A
123 16	112	AGAIN: LD (HL),B	;DRAW CHARACTER OF 8
123 17	189	CP L	;LOOK FOR END
123 18	202 25 123	JP Z,TIME	;IF END, DO TIME DELAY
123 21	25	INC HL	;ELSE SET FOR DRAW AGAIN
123 22	195 16 123	JP AGAIN	;JUMP TO DO AGAIN
123 25	22 255	TIME: LD D,255D	;SET MSB OF TIME DELAY
123 27	14 255	SETC: LD C,255D	;SET LSB OF TIME DELAY
123 29	13	DECC: DEC C	;COUNT DOWN LSB
123 30	194 29 123	JP NZ,DECC	;DO AGAIN IF NOT DONE
123 33	21	DEC D	;COUNT DOWN MSB
123 34	194 27 123	JP NZ,SETC	;DO AGAIN IF NOT DONE
123 37	201	RET	;RETURN AT END OF TIME

Whenever a value is to be passed from BASIC to a machine-language program, the first step in the machine-language program must be a CALL 2687D. The value passed from BASIC, in other words, has to be pulled from somewhere else by another subroutine—one executed by doing the CALL 2687.

And after that CALL is executed, the value of the argument will appear in the Z-80's HL register pair. It is then up to the programmer to do something with it.

A machine-language CALL 2687D places the value of a passed term into the HL register pair.

Here is a machine-language program and its BASIC calling program that illustrate the point:

```

123 0    205 127 10    CALL 2687D    ;GET THE PASSED VALUE
123 3     125         LD A,L        ;LOAD LSB TO A
123 4     33 0 60     LD HL,15360D  ;POINT TO START OF VIDEO
123 7     119         LD (HL),A     ;PUT VALUE ONTO SCREEN
123 8     201         RET           ;RETURN TO BASIC

```

```

0 REM ** CARRY A CHARACTER TO SUBROUTINE **
10 POKE 16526,0:POKE 16527,123

```

```

20 FOR I=0 TO 8:READ D:POKE 31488+I,D:NEXT I
30 CLS
40 C$=INKEY$:IF C$="" THEN 40
50 X=USR(ASC(C$))
60 GOTO 40
100 DATA 205,127,10,125,33,0,60,119,201

```

The BASIC portion of the scheme loads the machine-language program, then looks for a character string from the keyboard. On finding one, line 50 both converts the character string to its ASCII value and calls the machine-language subroutine. And after the machine-language routine does its job—to be described shortly—the BASIC program simply loops back to line 40, giving you a chance to enter another character.

Now look at the machine-language part of the project. The BASIC program passes `ASC(C$)` to the machine-language subroutine, and the first instruction in that subroutine calls the value from its mysterious place in the TRS-80 memory. On completing execution of that instruction the value passed from BASIC resides in the Z-80's HL register pair.

The second line in the machine-language program loads the content of register L into register A. The lsb of the value passed to the subroutine is thus saved in that register. There is no need to save the content of the H register in this particular case, because ASCII values cannot exceed 255 and are thus completely specified in a single, 1-byte register—the L register in this case.

So the second line saves the ASCII character code in register A. That is an important step, because the next instruction uses the HL pair for another purpose—to point to the beginning of the video memory. Then the ASCII code is moved from register A to the video memory, where the character appears on the screen.

The last line in the machine-language program simply returns operations to BASIC.

Not a very exciting program, perhaps, it simply prints whatever character is entered from the keyboard. But the entry operation is done in BASIC and the printing is done in machine language; and that illustrates the main point of this discussion: passing a value from BASIC to machine language.

The next BASIC/machine-language routine is another example of passing a value from BASIC. Unlike the previous example, this one has the potential for working with numbers that occupy a register pair—passing numbers that may be larger than 255.

The BASIC program loads the machine-language subroutine, then enters a loop that allows you to enter decimal numbers between 0 and 1023. For the purposes of this demonstration, those numbers represent the range of plotting positions on the crt. The numbers are

123	0	205 127 10	CALL 2687D	;GET THE PASSED VALUE TO HL
123	3	235	EX HL,DE	;GET PASSED VALUE TO DE
123	4	33 0 60	LD HL,15360D	;POINT TO VIDEO
123	7	175	XOR A	;CLEAR A REGISTER
123	8	54 131	DRAW:LD(HL),131D	;DRAW LINE SEGMENT
123	10	187	CP E	;LOOK FOR END OF LSB VALUE
123	11	194 16 123	JP NZ NEXT	;JUMP IF NOT DONE
123	14	186	CP D	;LOOK FOR END OF MSB VALUE
123	15	200	RET Z	;IF DONE, RETURN TO BASIC
123	16	27	NEXT:DEC DE	;COUNTDOWN VALUE
123	17	35	INC HL	;SET NEW VIDEO POINT
123	18	195 8 123	JP DRAW	;DRAW AGAIN

```

0 REM ** PRINT @ LINE DRAW **
10 POKE 16526,0:POKE 16527,123
20 FOR I=0 TO 20:READ D:POKE 31488+I,D:NEXT I
30 CLS
40 INPUT Y
50 IF Y<0 OR Y>1023 THEN 40
60 CLS
70 X=USR(INT(Y))
80 PRINT:GOTO 40
100 DATA 205,127,10,235,33,0,60,175
110 DATA 54,131,187,194,16,123,186,200
120 DATA 27,35,195,8,123

```

comparable to those used for doing PRINT @ operations in BASIC.

Instead of printing a character at a single point, however, the machine-language part of the program draws a continuous white line, beginning at point 0 on the screen and running to the value passed from the BASIC program.

The endpoint-coordinate of the line is INPUT at line 40. The conditional in line 50 goofproofs the system from numbers outside the video-memory range. Line 70 calls the machine-language subroutine, but also goofproofs the system against any fractional numbers the user might try to enter. Recall that numbers passed to a machine-language subroutine must be integer values. (Things will happen if you attempt to pass a noninteger value, but there will seem to be no rational correspondence between the specified number and the response of the machine-language program.)

Notice that the machine-language program begins with the mandatory CALL 2687. The step is mandatory, at least, if you want to work with the value passed from the BASIC part of the scheme. That operation always puts the passed value into the HL pair, but this writer wanted to use the HL pair as a video-memory pointer; so the next step is to exchange the HL and DE register contents. That effectively moves the passed value to the DE pair.

The remainder of the machine-language program simply plots character code 131 onto the screen, counts down the value passed

from BASIC, and moves to the next character space on the screen.

The machine-language routine continues looping in this fashion until the value passed to it is counted down to zero. Then the control is returned to BASIC, where you have a chance to enter another endpoint number between 0 and 1023.

The main points demonstrated by these two programs are:

1. A value passed from BASIC to machine language must be picked up at the machine-language level by executing a CALL 2687 (decimal).
2. A value picked up by a CALL 2687 will appear in the Z-80's HL register pair.
3. You are free to manipulate that value as it is passed to the HL register pair.

### **PASSING VALUES FROM MACHINE LANGUAGE TO BASIC**

A BASIC statement such as `Y=USR(X)` will call a machine-language subroutine and return to BASIC with a value assigned to the Y variable. An alternate form is `PRINT USR(X)`: a statement that will immediately print the value returned from the machine-language routine.

But in order to pass a value from a machine-language routine to BASIC, that value must be residing in the Z-80's HL register pair, and there is a very special, mandatory technique for returning to BASIC.

Whenever you want to pass a value back to BASIC, the machine-language program must end with this instruction:

```
195 154 10    JP 2714D    ;RETURN TO BASIC WITH VALUE
```

Such a routine does not return to BASIC with one of the usual sorts of RETurn instructions. Technically speaking, that mandatory JUMP instruction sends operations to a place in TRS-80 ROM that does the RETurn job for you, and in the meantime it takes care of the operations involved in assigning the value in the HL register pair to the USR function in the BASIC listing.

So when you are ready to get out of a machine-language subroutine and return a value to BASIC, you must make sure the value to be passed is in the HL pair, and then you must execute the special JUMP instruction, JP 2714D.

The following BASIC/machine-language program illustrates the point. This one counts and prints integer values. The counting operation, however, is done at the machine-language level, while the PRINTing is done in BASIC. That sort of scheme calls for passing the current count from machine language to BASIC.

```

123 0    42 10 123  LD HL,(COUNT)    ;FETCH COUNT TO HL
123 3    35          INC HL           ;COUNT BY ONE
123 4    34 10 123  LD (COUNT),HL    ;SAVE COUNT AT END OF PROGRAM
123 7    195 154 10 JP 2714D          ;RETURN VALUE TO BASIC

      0  REM ** COUNT AT MACHINE LEVEL **
      10 POKE 16526,0:POKE 16527,123
      20 FOR I=0 TO 11:READ D:POKE 31488+I,D:NEXT I
      40 PRINT USR(X);
      50 FOR T=0 TO 10:NEXT
      60 GOTO 40
     100 DATA 42,10,123,35,34,10,123,195,154,10,0,0

```

The BASIC portion of the program loads the machine-language program and then enters an endless loop between lines 40 and 60. That loop calls and prints a value from the machine-language subroutine and then does a short time delay at line 50. The time delay simply slows down the operation so you can observe it better. The only way to end the program is by striking the *BREAK* key.

The most important part of the machine-language routine is the last line. This is the mandatory instruction—the one that makes it possible to assign the number in the HL pair to the USR function and return to BASIC.

The count, itself, is saved at two bytes of memory just above the main program listing. Saving the count in that fashion is absolutely necessary, because most of the Z-80 registers are affected by the return to BASIC, and a count residing in them will be lost. The two zeros at the end of the DATA listing initialize that “counter” to zero but only when the program is first RUN. After that, the “counter” carries the current count.

In summary:

1. A value to be passed from machine language to BASIC must reside in the Z-80's HL pair at the moment the return occurs.
2. When carrying a value back to BASIC, the return must be carried out by executing a JP 2714 (decimal).
3. The value assigned to the USR function may be treated as any BASIC value.

## PASSING VALUES BACK AND FORTH

Some of the most useful kinds of BASIC/machine-language programs are those that pass one value from BASIC, perform some mathematical or logic operations on the value, and return the result to BASIC. The back-and-forth process is really just a matter of combining the principles described in the two previous sections of this chapter: passing values from BASIC to machine language, and passing values from machine language to BASIC.

The following program illustrates the point. This program allows you to enter any integer between -32,768 and 32,767. The integer is passed to the machine-language part of the program and a 2-byte version of the number is returned. In other words, the program carries out the tedious task of converting a standard decimal number into its 2-byte decimal form—something that frequently troubles programmers who are working with BASIC and machine language together.

123 0	205 127 10	START:	CALL 2687D	;VALUE TO HL PAIR
123 3	17 25 123		LD DE,PHASE	;SET PHASE ADDRESS
123 6	26		LD A,(DE)	;FETCH CURRENT PHASE
123 7	254 0		CP 0	;ZERO PHASE?
123 9	194 17 123		JP NZ,MSB	;IF NOT, THEN MSB
123 12	60		INC A	;SET FOR PHASE 1
123 13	18		LD (DE),A	;SAVE PHASE
123 14	195 20 123		JP DONE	;AND JUMP TO DONE
123 17	61	MSB:	DEC A	;SET FOR PHASE 0
123 18	18		LD (DE),A	;SAVE PHASE
123 19	108		LD L,H	;MSB TO L REGISTER
123 20	38 0	DONE:	LD H,0	;CLEAR H REGISTER
123 22	195 154 10		JP 2714D	;RETURN TO BASIC ;WITH VALUE

```

0 REM ** CONVERT TO 2-BYTE DECIMAL **
10 POKE 16526,0:POKE 16527,123
20 FOR I=0 TO 25:READ D:POKE 31488+I,D:NEXT I
40 CLS:INPUT "ENTER STANDARD DECIMAL (-32768 TO 32767)";N
50 IF N<-32768 OR N>32767 THEN 40
60 PRINT N="";
70 FOR PH=0 TO 1:PRINT USR(N);:NEXT PH
80 PRINT:PRINT
90 INPUT"STRIKE ENTER KEY TO DO AGAIN";$:GOTO 40
100 DATA 205,127,10,17,25,123,26,254,0,194,17,123
110 DATA 60,18,195,20,123
120 DATA 61,18,108,38,0,195,154,10,1

```

As in all previous examples the BASIC portion of this program both loads the machine-language portion of the program and performs some important i/o operations in its own right. The machine-language program is loaded by line 20, calling on the DATA items in lines 100, 110, and 120. The actual working part of the program begins at line 40 and loops endlessly between that line and line 90. The machine-language subroutine is called by the PRINT USR(N) statement in line 70.

The general idea of the program is to enter a decimal integer in response to the INPUT statement ENTER STANDARD DECIMAL (-32768 TO 32767). Line 50 goofproofs that entry, making sure the value passed to the machine-language subroutine is within its acceptable range (a matter described earlier in this chapter).

For the sake of doing some tidy formatting, line 60 reprints your entry and inserts an equal sign. The semicolon preceding the equal sign suppresses the usual line-feed/carriage-return operation, allowing the results of the machine-language program to be printed on that same line.

The FOR...NEXT statements in line 70 imply a two-phase operation, as far as the execution of the machine-language program is concerned. This two-phase operation allows PRINTing of the msb of the result on the first pass and PRINTing of the lsb on the second pass.

The same value is passed to the machine-language program in both phases—the value of N entered in response to the INPUT message. The machine-language program passes back a value two different times. The first value passed back to BASIC is the msb of the 2-byte decimal number, and the second value in the sequence is the lsb of that number.

So if you respond to the INPUT message with a 15360, the PRINTed results on the screen look like this: 15360=60 0. And, sure enough, the 2-byte decimal version of 15360 (the beginning address of the TRS-80 video memory) is 60 0, a figure used in several previous demonstrations.

The main point of the demonstration lies in the first and last lines of the machine-language portion of the program. The first line does the CALL 2687 (decimal) to insert a value passed from BASIC into the Z-80's HL pair. And the last line does the JP 2714 (decimal) to carry a value from machine language to BASIC.

### **SAVING BASIC/MACHINE-LANGUAGE PROGRAMS ON CASSETTE TAPE**

One of the advantages of using BASIC and machine-language programs together is that they are so easy to save on cassette tape. Since the machine-language part of the scheme is loaded from BASIC, the entire program can be saved on tape via the usual CSAVE command. Likewise, the programs can be loaded into the TRS-80 by means of the usual CLOAD command.

In short, there is no need to fool around with SYSTEM commands. As far as the TRS-80 is concerned, the programs are purely BASIC programs, and they can be treated as such.

There is much more to be said about this matter of linking BASIC and machine-language programs with the USR function. This chapter has presented the most fundamental parts of the matter. The next chapter deals with special procedures and applications.

## CHAPTER 6

# Manipulating BASIC-Loaded, USR-Linked Programs

The USR function allows the user to link together BASIC and machine-language programs. Chapter 5 describes the process in general terms, and the demonstrations show how values can be passed between the BASIC and machine-language portions of the program.

This chapter, like Chapter 5, deals with BASIC-loaded, USR-linked programs, but from a somewhat more sophisticated point of view. Here, for instance, you will see how it is possible to enter the machine-language portion of a program at any desired point, and not just from the lowest address. This principle then leads to the notion of calling any one of a number of different machine-language programs.

To put some finishing touches on the subject you will see how to delete the loader portion of a BASIC program, leaving behind just the operational part of BASIC and the machine-language subroutines.

### **SPECIFYING ENTRY POINTS FOR USR-LINKED PROGRAMS**

When writing the BASIC portion of a program that links BASIC and machine language via the USR statement, it is necessary to specify the starting address of the machine-language routine before executing the USR function. That starting address is POKEd into 16526 and 16527 as a 2-byte version of the address. See Chapter 5 if you have any doubts about the meaning of this important principle.

The POKE statements, however, do not necessarily have to point to the very beginning of a machine-language subroutine. They can point to any valid *entry point*—any valid point in the subroutine that will execute a set of instructions without causing confusion.



The practical significance of this notion is that you can write machine-language subroutines that do a number of different tasks and you can call from BASIC any one you choose. It's all a matter of setting the desired entry point at addresses 16526 and 16527.

Before getting into the details of a program that illustrates this point, it is instructive to view it from a general angle. The program you will enter is a rather simple counting program. It uses a 46-byte machine-language routine to count upward or downward between 0 and 9. It can also preset the count to 0 or 9. All operations are under the control of the BASIC portion of the program—the machine-language portion simply carries out the decisions made at the BASIC level.

Table 6-1 summarizes the four entry points for the machine-language program. The first column describes the operation, the second specifies the starting address of that operation in standard decimal form, and the third column specifies the starting addresses in a 2-byte decimal format.

**Table 6-1. Summary of the Entry Points for the Machine-Language Program in Table 6-2**

Function	Entry-Point Address		
	Standard Decimal	2-Byte Decimal	
		MSB	LSB
ZERO the counter	31488	123	0
Preset to NINE	31493	123	5
Count UP	31501	123	13
Count DOWN	31513	123	25

Getting to any one of the four valid entry points of the machine-language program is a matter of first setting the entry-point address and then doing the USR function. To zero the counter, for example, the BASIC program should include the sequence: POKE 16526,0: POKE 16527,123. A USR function appearing anytime after that will cause the system to enter the machine-language program at address 31488, thus causing the counter to reset to zero.

Presetting the counter to 9 is a matter of specifying the entry point for that particular operation: POKE 16526,5:POKE 16527,123. A USR function appearing anytime after that will call the preset-to-9 operation at the machine-language level.

In a similar fashion, doing a POKE 16526,13:POKE 16527,123 sets the entry point for counting upward, and doing a POKE 16526,25: POKE 16527,123 sets up a down-counting operation.

The standard Radio Shack literature says that Level II BASIC can call just one machine-language subroutine from the USR function.

It might be more precise to say that Level II BASIC allows the assignment of just one USR-linked operation at any given time. There is nothing to prevent you from specifying different starting points as the program progresses, thus achieving the effect of accessing any number of machine-language entry points you choose.

Table 6-2 is the machine-language portion of this counting program. The entry points are at ZERO, NINE, UP, and DOWN. It is a rather tightly structured program that shares two other routines with each of the four entry points: NEXT and DOIT.

Referring to the machine/assembly program in Table 6-2, the NEXT routine saves the current count from the accumulator in memory location 31535, then does a jump to routine DOIT. DOIT converts the count to an ASCII character code, prints the result in the screen upper left-hand corner, and returns operations to BASIC.

NEXT and DOIT are accessed from all four of the program's main entry points: ZERO, NINE, UP, and DOWN. ZERO sets the accumulator to 0, NINE sets it to 9, UP either increments the accumulator or resets it to 0 after showing a 9, and DOWN either decrements the accumulator or resets it to 9 after showing a 0.

Enter the machine-language program at ZERO, and you will see a 0 printed on the screen. Enter at NINE, and you will see a 9. Enter at UP or DOWN, and you will see the character increment or decrement within the counting range of 0 through 9.

The BASIC portion of this program must do two things: it must enter the machine-language portion of the program and control the selection of the entry point. The most important feature, as far as the present discussion is concerned, is the adjustment of the entry point just prior to executing a USR function. Here is the BASIC program:

```

10  FOR I=0 TO 47:READ D:POKE 31488+I,D:NEXT I
20  DATA 62,0,195,7,123
22  DATA 62,9,50,47,123,195,37,123
24  DATA 58,47,123,254,9,210,0,123,60,195,7,123
26  DATA 58,47,123,254,0,202,5,123,61,195,7,123
28  DATA 33,0,60,58,47,123,198,48,119,201,48
100  RL=16526:RH=16527
110  CLS
120  IF PEEK(14463)=0 THEN 120
130  C$=INKEY$
140  IF C$<>"" THEN R$=C$
150  IF R$="Z" THEN L=0
160  IF R$="N" THEN L=5
170  IF R$="U" THEN L=13
180  IF R$="D" THEN L=25
190  IF NOT(R$="Z" OR R$="N" OR R$="U" OR R$="D") THEN 120
200  POKE RL,L:POKE RH,123
210  X=USR(X)
220  FOR T=0 TO 10:NEXT T
230  GOTO 120

```

Table 6-2. Machine-Language Program for the Controlled Counter Routine

Address	Op Code	Mnemonic	Comment
123 0	62 0	LD A,0	;ZERO THE ACCUMULATOR
123 2	195 7 123	JP NEXT	;JUMP TO DO NEXT
123 5	62 9	LD A,9D	;LOAD 9 IN THE ACCUMULATOR
123 7	50 47 123	LD (COUNT),A	;LOAD ACC. TO COUNT
123 10	195 37 123	JP DOIT	;JUMP TO DOIT
123 13	58 47 123	LD A,(COUNT)	;COUNT TO ACCUMULATOR
123 16	254 9	CP 9D	;MAXIMUM COUNT?
123 18	201 0 123	JP NC ZERO	;IF SO, JUMP TO ZERO
123 21	60	INC A	;INCREMENT ACCUMULATOR
123 22	195 7 123	JP NEXT	;JUMP TO NEXT
123 25	58 47 123	LD A,(COUNT)	;FETCH CURRENT COUNT
123 28	254 0	CP 0	;MINIMUM COUNT?
123 30	202 5 123	JP Z NINE	;IF SO, JUMP TO NINE
123 33	61	DEC A	;ELSE DECREMENT ACCUMULATOR
123 34	195 7 123	JP NEXT	;JUMP TO NEXT
123 37	33 0 60	LD HL,VIDEO	;POINT TO START OF VIDEO
123 40	58 47 123	LD A,(COUNT)	;FETCH CURRENT COUNT
123 43	198 48	ADD A,48D	;CONVERT TO ASCII
123 45	119	LD (HL),A	;PRINT CHARACTER
123 46	201	RET	;RETURN TO BASIC

Lines 10 through 28 make up the machine-language loader. The procedure is identical with the one described in Chapter 5, with line 10 doing the actual loading operation and the DATA in lines 20 through 28 carrying the instruction bytes that are justified in Table 6-2.

So the loading of the machine-language portion of the program is completed at line 28.

Line 100 points to the USR entry-point registers. Those points, here assigned variable names RL and RH, are the same ones used for all USR-linked programs. Note, however, that the address they are to specify is not indicated—not yet, anyway.

Lines 120 through 140 make up a key-depression sensing operation: one that responds as long as a key is being depressed. Recall the details from Chapter 3. Whenever a key is being depressed, its string value is assigned to variable R\$ at line 140.

Lines 150 through 180 decode four different key depressions, assigning some numerical values to variable L. As long as the Z key is being depressed, for instance, variable L takes on the value of zero.

Those values of L represent the lsb of the entry-point address for the machine-language subroutines; and they are assigned to the USR's entry-point register at line 200. So if you are depressing the Z key, line 200 performs this operation for you: POKE 16526,0:POKE 16527,123. This is precisely the set of operations required for starting the USR-called machine-language routine at address 31488—the entry point for presetting the counter to zero. Depressing the N, U, or D keys set up things in the same way, pointing to the entry points for NINE, UP counting, and DOWN counting, respectively.

It is not only important to set up the proper entry points to be assigned at line 200, but in a program of this sort it is equally important to eliminate key depressions that will provide other, unwanted values for L. You certainly do not want to specify some undetermined entry points, and that is the purpose of line 190.

Without line 190, you could depress any other key and end up specifying an entry point having an lsb of some undetermined value; and that could be disastrous, sending the program to some strange point in the system and, most likely, crashing the whole program out of existence.

#### **IMPORTANT**

BASIC programs that call multiple entry points in a machine-language program must be structured in such a way that unwanted entry points cannot be specified.

Line 210 calls the specified machine-language routine. When the routine is completed, line 220 does a short time delay. The delay simply makes it easier for you to watch the counting operation. Line 230 returns operations to the key-entry point at line 120.

Summarizing the main points of this discussion:

1. Write a machine-language program, keeping close track of the addresses of all possible entry points.
2. Write a BASIC program that:
  - a. Loads the machine-language portion.
  - b. Controls the entry-point addresses before doing the USR function.
  - c. Eliminates the possibility of specifying undetermined machine-language entry points.

Once this is done, you can enter the machine-language portion of the program at any desired point, creating the effect of having a system that can call any number of entry points or machine-language sub-routines.

### **DELETING THE MACHINE-LANGUAGE LOADER**

Through all of the BASIC/machine-language programs presented thus far, the machine-language portion is loaded from BASIC (via DATA lists). Once the machine language is loaded, the loader line and DATA lists perform no useful function—all subsequent operations at the BASIC level involve doing some controls and implementing the USR function to call the machine-language subroutines that are stored in protected memory space.

Once the machine-language portion of the program is loaded from BASIC, the lines of the BASIC program devoted to that process can be deleted. In the example offered in the previous section of this chapter, for example, you can do a RUN to get it loaded; and once it begins running, you can do a BREAK, and then DELETE 10-28. That DELETE command will eliminate the loading portion of the BASIC program, leaving the operational part intact. And as long as the machine-language program still resides in its place in protected memory, the program can be run at your heart's content.

It is a good idea to retain the loading portion of a BASIC/machine-language program until you have a chance to test and debug it. But once you are confident it is running as it is supposed to run—and the machine-language portion is loaded into protected memory space—you can wipe out the loading portion of the BASIC program without altering the operation of the program in any way.

Suppose that you have worked out the program presented in the previous example. It is working to your satisfaction, but you'd like

to tidy up the matter by deleting the loading portion. RUN the full program one more time, just to make sure the machine-language program is, indeed, properly loaded into protected memory, then do a BREAK to interrupt the program at the BASIC level.

Next, CSAVE the entire program on cassette tape (or disk), and then DELETE 10-28. A LIST will show that the loading part of the program is gone; but when you do a RUN you will find the program running just as well as it did before. (The idea of saving the program on tape before deleting the loader makes it possible for you to enter the entire program—including the machine-language loader—at any later time.)

Doing those BREAKs and DELETEs as described here is rather inelegant, however. An elegant way to get rid of the loader at the BASIC level is to write an appropriate DELETE statement right into the BASIC program. Write in a line that will delete the machine-language loader immediately after that loading operation is done.

In the previous example you can insert a new line:

**30 DELETE 10-30**

Insert that statement into the program, and it will delete the loader and line 30 as well, after the loading operation is done. Try it.

Get the loader back into your system. (Hopefully, you have saved it on tape or disk as suggested earlier; otherwise you will have to type in the loading line and DATA lines again.) Then insert the suggested DELETE 10-30 at line 30.

Save on tape or disk, then do a RUN. After a short delay, caused by the loading operation, the system will return a READY? and, indeed, the program is ready to run. From that point on, the loader will be missing from the BASIC program, but the whole thing does its intended job quite nicely.

In fact, you can do a NEW, and the machine-language portion of the program will be unaffected. It will still be residing there in protected memory space, and you can write another BASIC program that calls it.

DELETEing the machine-language loader, either manually from the keyboard or automatically by a built-in DELETE statement, creates a much neater BASIC program. But there are certainly some more compelling reasons for doing the DELETE.

One good reason for DELETEing the machine-language loader is to make more room in the BASIC text memory for elaborate BASIC operations. This is especially desirable in cases where the machine-language portion of the program is a rather extensive one as well. Why bother keeping all those DATA lists in BASIC memory after they have performed their intended task?

Another reason for getting rid of the BASIC-level, machine-language loader is that some machine-language programs are intended to manipulate any kind of BASIC program you might want to enter. A line-renumbering utility, for example, can be loaded as a machine-language subroutine from BASIC. The program you want to renumber, however, will be an entirely different one. So the idea is to write a machine-language loader in BASIC, do the loading operation, and then completely DELETE the BASIC portion of the program. After that, any USR-oriented set of operations will call up the subroutine for you, presumably doing things such as renumbering the resident BASIC program.

Finally, the ability to DELETE BASIC-level, machine-language loaders makes it possible to create some rather elaborate machine-language programs in a bottom-up fashion. Machine-language routines can be loaded and tested piecemeal, moved around in the protected memory at will, and merged with other machine-language routines. This feature is the subject of the next section in this chapter.

### **BUILDING MACHINE-LANGUAGE PROGRAMS FROM THE BOTTOM UP**

For a beginner, writing machine-language programs can be a difficult and emotionally trying task. It takes a whole lot of experience to feel at ease with this sort of programming, and anyone not feeling at ease ought to use as many programming aids as possible.

The idea of writing USR-called, decimal-oriented programs can be of some help for beginners. And one way to take advantage of USR-called machine-language programming is to write and test basic elements of the routines in a piecemeal fashion, working on one element until it is working and then moving to another element. Once the elements are all entered and tested, they can be strung together to make up one large program—a machine-language program that might appear far too imposing to tackle in one shot.

The following operations demonstrate this particular approach to building a machine-language program. The general idea is to compose a USR-called subroutine that draws a large rectangle on the crt screen. The procedure will be to write, enter, and test the drawing operations for the four sides separately, and then assemble things to come up with one machine-language program for doing the job.

#### **Drawing the Top of the Rectangle**

Here is a machine-language and BASIC program for drawing a white line across the top of the screen:

123 0	33 0 60	BORDER:	LD HL,15360D	;POINT TO START
123 3	62 63		LD A,63D	;SET END

```

123 5      189      TOP:      CP L          ;AT END?
123 6      202 20 123      JP Z,NEXT1      ;JUMP IF SO
123 9      54 131          LD (HL),131D      ;ELSE DRAW LINE
123 11     44            INC L          ;SET NEXT SPOT
123 12     195 5 123      JP TOP          ;DO MORE DRAWING
123 20     201      NEXT1:  RET          ;RETURN TO BASIC

0  REM ** BORDER-DRAWING DEMO **
10 DATA 33,0,60,62,63,189,202,20,123,54,131,44,195,5,123
50 REM ** LOADER FOR TOP **
55 FOR I=0 TO 14:READ D:POKE 31488+I,D:NEXT I
57 POKE 31508,201
100 REM ** CALLING PROGRAM **
110 POKE 16526,0:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140

```

Referring to the machine-language portion of the program, it begins by loading the lowest video memory address into the HL register pair. This points to the upper left-hand corner of the screen, the starting point of the line marking the top of the rectangular figure to be drawn.

Decimal 63, marking the endpoint of the TOP line, is loaded into the accumulator, and then the content of the L register is compared with that number. If the numbers do not match, presumably because the content of L is less than that of the accumulator (and the drawing operation isn't done), graphic code 131 is drawn on the screen at the point indicated by the HL pair. Then the L register is incremented one space, and control returns to the TOP routine.

The routine continues drawing graphic 131 until the L register is incremented to 63. At that point the TOP-line drawing is done, and control is sent to memory location 31508 (or 123 30 as expressed as a 2-byte decimal address).

That NEXT1 location contains a RETurn instruction, but only temporarily. When it is time to add a routine to draw the bottom of the rectangular figure, it will begin at that point. For testing purposes, however, this address returns control to the BASIC calling program.

Note that there are some unused bytes between the last instruction in the TOP routine and the beginning of NEXT1. It is generally considered good programming procedure to allow a few unused bytes between the various elements of a machine-language program. The same idea applies to designating line numbers for BASIC programs. In either case the rationale is to allow some room for expanding the routine without having to adjust addresses (or line numbers in BASIC) from that point through the end of the program.

So it is advisable to leave some unused memory locations in a ma-



chine-language program. Just be sure, however, to use an appropriate JUMP instruction to pass over those blank spaces. That's done in this case with the JP Z, NEXT1 instruction.

The BASIC part of the program shows the object codes as a DATA line. Those codes are entered into protected memory space by line 55. Note that the loading begins at address 31488 (or 123 0 in a 2-byte decimal format).

Line 57 is necessary in order to answer the jump-over-blank-places operation. You could add a bunch of zeros to the DATA list, and adjust the FOR. .NEXT statement in line 55 accordingly. But that would place an unnecessary burden on the DATA list. The statement in line 57 is thus added separately.

Lines 100 through 140 execute the program. If all is going well to this point, doing a RUN will load the machine-language segment and execute it, drawing a clean line across the top of the screen. Because of the tight loop built into line 140, you have to strike the BREAK key to get something else going.

Anything wrong with your program will show up right away. If the program blows up, it is short enough to enter again from scratch—after, of course, you work out the bug that caused the blowup. Or you can save the program on tape *before* doing a RUN. That way you can reload it in a few moments, make the fix and give it another try—all in a fairly short period.

### Drawing the Bottom of the Rectangle

After entering and debugging the first element of the program, it is possible to move to the second phase without disturbing the good stuff that is already in place. Here are the machine-language and BASIC routines for drawing the bottom portion of the rectangle:

123 20	33 192 63	NEXT1:	LD HL,16320D	;POINT TO START
123 23	62 255		LD A,255D	;SET END
123 25	1B9	BOT:	CP L	;AT END?
123 26	202 40 123		JP Z,NEXT 2	;IF SO, GET OUT
123 29	54 176		LD (HL),176D	;ELSE DRAW LINE
123 31	44		INC L	;SET NEXT SPOT
123 32	195 25 123		JP BOT	;DO MORE DRAWING
123 40	201	NEXT2:	RET	;RETURN TO BASIC

```

0 REM ** BORDER-DRAWING DEMO **
10 DATA 33,0,60,62,63,1B9,202,20,123,54,131,44,195,5,123
12 DATA 33,192,63,62,255,1B9,202,40,123,54,176,44,195,25,123
50 REM ** LOADER FOR TOP **
55 FOR I=0 TO 14:READ D:POKE 314BB+I,D:NEXT I
60 REM ** LOADER FOR BOTTOM **
65 FOR I=0 TO 14:READ D:POKE 3150B+I,D:NEXT I
67 POKE 3152B,201
100 REM ** CALLING PROGRAM **

```

```

110 POKE 16526,20:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140

```

The machine-language portion of the bottom-line-drawing routine isn't much different, in principle, from the TOP routine already entered and tested. The only difference here is the position of the line and, of course, the addresses of the machine-language programming.

The BASIC part of the program still contains elements of the TOP-drawing routine, but the DATA line (line 12) is added for the bottom line. Line 57 has been deleted from the program, because the end of TOP now calls the beginning of BOT—or at least it will when the two routines are merged later on.

At any rate the BOT routine begins at address 31508, as indicated by its loader in line 65. A RETURN is inserted at 31528, making it possible to get out of the routine at the point the next element of the program will begin.

Note in line 110 that the starting point of the USR-called routine is different. Instead of starting at address 31488, it begins at 31508. That means the USR function calls the new BOT routine instead of the TOP routine.

Run the program as shown here and you should see a white line across the bottom of the screen.

### Merging the Two Elements of the Routine

The following BASIC program shows how it is possible to merge the TOP and BOT routines into one. The BASIC program has been cleaned up a bit, and the USR-called starting point is set back to the beginning of the machine-language program, to address 31488. So the two sections of the machine-language program automatically run in succession, responding to a single USR function in BASIC.

```

0 REM ** BORDER-DRAWING DEMO **
10 DATA 33,0,60,62,63,189,202,20,123,54,131,44,195,5,123
12 DATA 33,192,63,62,255,189,202,40,123,54,176,44,195,25,123
50 REM ** LOADER FOR TOP AND BOTTOM **
52 FOR I=0 TO 14:READ D:POKE 31488+I,D:NEXT I
54 FOR I=0 TO 14:READ D:POKE 31508+I,D:NEXT I
56 POKE 31528,201
100 REM ** CALLING PROGRAM **
110 POKE 16526,0:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140

```

In a sense this is not a very efficient program scheme. Leaving some unused memory space between the TOP and BOT routines, for

example, makes it necessary to use two different loading operations in the BASIC program. But writing efficient or elegant programs is not the point of the present discussion. The main idea is to demonstrate how USR-called programs can be composed by sections, then worked together to make up a single program scheme.

Now it is time to add in some more programming to draw the two sides of the rectangle.

### Drawing the Sides and Merging Again

Here is the programming for drawing the left side of the rectangle. The machine-language portion deals only with the left-side drawing operation, but the BASIC part of the matter includes the TOP and BOT programming already developed. You will note, however, that the USR function calls only the left-drawing part of the machine-language programming; this lets you concentrate on the new problems that might arise.

```

123 40      30 0 60      NEXT2:  LD HL,15360D      ;POINT TO TOP
123 43      54 191      LEFT:    LD (HL),191D      ;DRAW SEGMENT
123 45      125          ID A,L      ;FETCH LSB
123 46      198 64          ADD A,64D      ;ADD 64
123 48      218 55 123      JP C,MORE1      ;JUMP IF DONE
123 51      111          LD L,A      ;SAVE NEW LSB
123 52      195 43 123      JP LEFT      ;DO AGAIN
123 55      124          MORE1:  LD A,H      ;FETCH MSB
123 56      254 63          CP 63D      ;DONE?
123 58      202 70 123      JP Z NEXT3      ;IF SO, GET OUT
123 61      36          INC H      ;SET NEW MSB
123 62      46 0          LD L,0      ;RESTART LSB
123 64      195 43 123      JP LEFT      ;DO ALL OVER AGAIN

0 REM ** BORDER-DRAWING DEMO **
10 DATA 33,0,60,62,63,189,202,20,123,54,131,44,195,5,123
12 DATA 33,192,63,62,255,189,202,40,123,54,176,44,195,25,123
14 DATA 33,0,60,54,191,125,198,64,218,55,123,111,195,43,123,124,254,63,202,70,
123,36,46,0,195,43,123
50 REM ** LOADER FOR TOP AND BOTTOM **
52 FOR I=0 TO 14:READ D:POKE 31488+I,D:NEXT I
54 FOR I=0 TO 14:READ D:POKE 31508+I,D:NEXT I
60 REM ** LOADER FOR LEFT **
65 FOR I=0 TO 26:READ D:POKE 31528+I,D:NEXT I
67 POKE 31558,201
100 REM ** CALLING PROGRAM **
110 POKE 16526,40:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140

```

DATA line 14 in the BASIC programming contains the machine-code instructions, and they are loaded by the statements in line 65. Line 67 inserts a RETURN into the machine-language program, forc-

ing the scheme to return to BASIC after drawing the left-side line on the screen. If there are any problems with the program, they will certainly show up at this point, giving you a chance to remedy the trouble without affecting the other parts of the program—the parts you have presumably tested before.

Assuming that the left-line drawing operation works, the next step is to add in the right-line drawing sequence.

123	70	33 63 60	NEXT3:	LD HL,15423D	;POINT TO TOP
123	73	54 191	RIGHT:	LD (HL),191	;DRAW SEGMENT
123	75	125		LD A,L	;FETCH LSB
123	76	198 64		ADD A,64D	;ADD 64 TO LSB
123	78	21B 85 123		JP C,MORE2	;JUMP IF DONE
123	81	111		LD L,A	;SAVE NEW LSB
123	82	195 73 123		JP RIGHT	;DO AGAIN
123	85	124	MORE2:	LD A,H	;FETCH MSB
123	86	254 63		CP 63D	;DONE?
123	88	202 100 123		JP Z NEXT4	;IF SO, GET OUT
123	91	36		INC H	;SET NEW MSB
123	92	46 63		LD L,63D	;RESTART LSB
123	94	195 73 123		JP RIGHT	;DO ALL OVER AGAIN

```

0 REM ** BORDER-DRAWING DEMO **
10 DATA 33,0,60,62,63,189,202,20,123,54,131,44,195,5,123
12 DATA 33,192,63,62,255,189,202,40,123,54,176,44,195,25,123
14 DATA 33,0,60,54,191,125,198,64,218,55,123,111,195,43,123,124,254,63,202,70,
    123,36,46,0,195,43,123
16 DATA 33,63,60,54,191,125,198,64,218,85,123,111,195,73,123,124,254,63,202,100,
    123,36,46,63,195,73,123
50 REM ** LOADER FOR TOP AND BOTTOM **
52 FOR I=0 TO 14:READ D:POKE 31488+I,D:NEXT I
54 FOR I=0 TO 14:READ D:POKE 31508+I,D:NEXT I
60 REM ** LOADER FOR LEFT **
65 FOR I=0 TO 26:READ D:POKE 31528+I,D:NEXT I
70 REM ** LOADER FOR RIGHT **
75 FOR I=0 TO 26:READ D:POKE 31558+I,D:NEXT I
77 POKE 31588,201
100 REM ** CALLING PROGRAM **
110 POKE 16526,70:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140

```

The BASIC part of the program is now arranged to call the new section of the machine-language program—the part that is supposed to draw the right-side line on the rectangle. Its function is quite similar to the left-side drawing operation; only the addresses are different.

Running that BASIC program should draw the right-side line for you. And if it works, you are ready to tidy up the BASIC portion of the program. The BASIC program shown next is cleaned up so that

it loads the machine-language programming as efficiently as possible, given the fact that some blank spaces exist in the memory between each of the basic sections. Line 110 sets the entry point of the USR function to the very beginning of the machine-language program. So when you run the program it will draw the complete rectangle for you. As usual, you have to strike the BREAK key to get out of the loop created by the BASIC statement in line 140.

```
0 REM ** BORDER-DRAWING DEMO **
10 DATA 33,0,60,62,63,189,202,20,123,54,131,44,195,5,123
12 DATA 33,192,63,62,255,189,202,40,123,54,176,44,195,25,123
14 DATA 33,0,60,54,191,125,198,64,218,55,123,111,195,43,123,124,254,63,202,70,
123,36,46,0,195,43,123
16 DATA 33,63,60,54,191,125,198,64,218,85,123,111,195,73,123,124,254,63,202,100,
123,36,46,63,195,73,123
50 REM ** LOADERS **
52 FOR I=0 TO 14:READ D:POKE 31488+I,D:NEXT I
54 FOR I=0 TO 14:READ D:POKE 31508+I,D:NEXT I
56 FOR I=0 TO 26:READ D:POKE 31528+I,D:NEXT I
58 FOR I=0 TO 26:READ D:POKE 31558+I,D:NEXT I
60 POKE 31588,201
100 REM ** CALLING PROGRAM **
110 POKE 16526,0:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140
```

If this whole rectangle-drawing routine is to be just one element in a larger program, you can insert a DELETE 0-65 at line 65 in the BASIC program. On running the program, then, the BASIC section will first load the machine-language programming, then delete all the DATA lines and loading statements. You can always call the rectangle-drawing routine at a later time by doing the BASIC sequence in lines 100 through 140.

### **Some Comments About Bottom-Up Programming**

The expression *bottom-up programming* refers to a programming technique whereby the program is developed by one section at a time. Each section, or *module*, is capable of standing alone, and it is thoroughly debugged before the programmer goes on to the matter of developing another module. And when all the modules have been thus written and debugged, they are tied together to make up one big program. Chances are good the final version will work rather well, because the separate modules have been tested beforehand. This was the technique used for building up the rectangle-drawing routine.

While this bottom-up technique has the advantage of letting the programmer work with the program in a piecemeal fashion, dividing

the job into a number of manageable parts, it does have its disadvantages. The most common disadvantage is that the final program is inefficient in terms of the amount of memory it uses. It is possible, for instance, to write the rectangle-drawing program using a whole lot less memory space.

An alternative to bottom-up programming is a technique known as *top-down programming*. In this case the entire program is developed from beginning to end. The procedure calls for some careful planning ahead of time and a form of thinking that is outside the scope of our present discussion. The result, though, is a program that is generally shorter than a version developed by a bottom-up procedure.

But there is something else you can do at this point to tighten up the programming. Suppose you have developed a USSR-called machine-language program, using the bottom-up, module-oriented technique used in the previous example—the one that draws a rectangle on the screen. You have the program working and you are satisfied with the operation of the thing.

In the process of working out the program you have learned some things about getting the job done, perhaps learning some of those things by making some mistakes. That being the case, you are in a good position to rewrite the program in a more efficient manner.

Looking over that program, you should be able to see that the top and bottom drawing sections of the machine-language listings are quite similar. The same happens to be true for the left and right drawing parts of the program. If nothing else, you should be able to condense the program to at least two basic sections, instead of four.

That is just an example of what you can do with programs developed by the bottom-up process. Chances are good that any program developed in this fashion can be tightened up if you apply a bit of thought to the matter.

The basic idea here is that you can take advantage of what you've done while perfecting the bottom-up procedure, thus avoiding a lot of hard-to-find mistakes in a more elegant revision.

### **Tightening Up the Program**

Here is a condensed version of the border-drawing program. The machine-language portion is shown in Table 6-3. It uses only 65 bytes of protected memory space, beginning at address 31488. The earlier version also started at that address, but it ate up 100 bytes of memory.

The revised BASIC portion of the scheme loads this new machine-language program, then executes it from lines 100 through 140.

Load the BASIC program, CSAVE it, and then give it a try. If there are any errors in the scheme, you will have a recorded copy to

**Table 6-3. Machine-Language Program for BORDER-DRAWING DEMO, V.2**

Address	Op Code	Mnemonic	Comment
123 0	33 0 60	BORDER:	
123 3	62 63	LD HL,15360D	;POINT TO TOP CORNER
123 5	6 131	LD A,63D	;SET ENDPOINT
123 7	205 37 123	LD B,131D	;SET CHARACTER
123 10	33 192 63	CALL DRAW1	;DRAW TOP
123 13	62 255	LD HL,16320D	;POINT TO BOTTOM CORNER
123 15	6 176	LD A,255D	;SET ENDPOINT
123 17	205 37 123	LD B,176D	;SET CHARACTER
123 20	33 0 60	CALL DRAW1	;DRAW BOTTOM
123 23	6 0	LD HL,15360D	;POINT TO TOP CORNER
123 25	205 44 123	LD B,0	;SET RESET POINT
123 28	33 63 60	CALL DRAW2	;DRAW RIGHT SIDE
123 31	6 63	LD HL,15423D	;POINT TO TOP RIGHT
123 33	205 44 123	LD B,63D	;SET RESET POINT
123 36	201	CALL DRAW2	;DRAW LEFT SIDE
123 37	189	RET	;RETURN TO BASIC
123 38	200	CP L	;ENDPOINT?
123 39	112	RET Z	;RETURN IF SO
123 40	44	LD (HL),8	;ELSE DRAW SEGMENT
123 41	195 37 123	INC L	;SET NEW POINT
123 44	54 191	JP DRAW1	;DRAW AGAIN
123 46	125	LD (HL),191D	;DRAW CHARACTER
123 47	198 64	LD A,L	;FETCH LSB OF PLACE
123 49	218 56 123	ADD A,64D	;SET NEW PLACE
123 52	111	JP C,MORE	;JUMP IF DONE
123 53	195 44 123	LD L,A	;ELSE SAVE NEW PLACE
123 56	124	JP DRAW2	;AND DRAW AGAIN
123 57	254 63	LD A,H	;FETCH MSB OF PLACE
123 59	200	CP 63D	;DONE?
123 60	36	RET Z	;RETURN IF SO
123 61	104	INC H	;ELSE GET NEW PLACE
123 62	195 44 123	LD L,B	;RESET STARTING POINT
		JP DRAW2	;AND DRAW AGAIN

work with. And even if there aren't any errors, you will have a copy that is ready to go anytime you need it.

```
0 REM ** BORDER-DRAWING DEMO, V.2 **
10 DATA 33,0,60,62,63,6,131,205,37,123,33,192,63,62,255,6,176,205,37,123
12 DATA 33,0,60,6,0,205,44,123,33,63,60,6,63,205,44,123,201
14 DATA 189,200,112,44,195,37,123
16 DATA 54,191,125,198,64,218,56,123,111,195,44,123
18 DATA 124,254,63,200,36,104,195,44,123
50 REM ** LOADER **
60 FOR I=0 TO 64:READ D:POKE 31488+I,D:NEXT I
100 REM ** CALL BORDER **
110 POKE 16526,0:POKE 16527,123
120 CLS
130 X=USR(X)
140 GOTO 140
```

## SUMMARY

The main point of this section is to illustrate how it is possible to use BASIC/machine-language programming to develop useful, USR-called subroutines. The discussion uses one specific example, drawing a rectangular figure, but the general procedures are applicable to any sort of routine.

So you can grasp the general ideas without thinking only in terms of a rectangle-drawing program, here are the basic points:

1. Divide the machine-language task into small elements that are each capable of doing something meaningful.
2. Compose a machine-language program for one of those elements.
3. Compose a BASIC program that both loads the machine-language program and executes it in a meaningful way.
4. Debug that portion of the scheme, getting it to work as you want it to work.
5. Repeat Steps 2 through 4 until all the smaller elements of the machine-language program have been tested.
6. Write a BASIC program that loads all the smaller elements of the machine-language program and executes them in some meaningful way.
7. Work out any bugs and, if possible, tighten up the BASIC portion of the program.
8. If you wish, insert a BASIC line that will DELETE the loader-related lines of the BASIC program, leaving just the USR-function part.
9. If you wish, use the working machine-language program to rewrite a tighter, more elegant version of it. Revise the BASIC portion accordingly.



If you save everything on cassette tape or disk as you go along, a blowup won't cost you a whole lot of time and effort. In the case of very complex machine-language programs you can divide the task into several main sections, and then divide those sections into even smaller subsections. You can build and test the subsections, one at a time, and then begin fitting them into the larger scheme.

The idea is to take advantage of the fact that BASIC-loaded, USR-called subroutines can be saved on tape, disk, or a line printer rather easily. And since the BASIC part of the program coordinates much of the machine-language programming, it is relatively easy to change things around as you go along.

Why not dream up some machine-language tasks of your own, then test your understanding of this procedure? If you have done little or no machine-language programming before, this is the most suitable approach for you. It lets you get around some of the more tedious machine-language tasks, passing them to a more familiar BASIC format.

## CHAPTER 7

# Hexadecimal Programming With T-BUG

All of the discussions and demonstrations in Chapters 2 through 6 deal with the TRS-80 system in a decimal format. Chapters 2, 3, and 4 consider some BASIC operations with the video memory, keyboard “memory,” and user’s memory—all in the decimal-oriented TRS-80 BASIC. Chapters 5 and 6 blend BASIC and some machine language, but since the machine-language programs are entered from BASIC in those two chapters, the whole business is decimal oriented.

You probably are aware, however, that the decimal number system is not the simplest system for writing machine-language programs, and that fact made the BASIC-loaded programs in Chapters 5 and 6 a bit tedious to write. All those conversions and 2-byte decimal numbers are quite troublesome. The hexadecimal number system is far more compatible with machine-language programming.

Unfortunately, the TRS-80 keyboard and display mechanisms are not set up to handle hexadecimal notation directly; they are set up for decimal-oriented BASIC programming. For this reason some other software mechanism has to be added to provide a hexadecimal i/o format. One such mechanism is Radio Shack’s T-BUG monitor and debugging tool—the subject of this chapter.

T-BUG is available from Radio Shack (catalog No. 26-2001) at a reasonable price. One side of the T-BUG cassette tape loads for Level I machines, and the other loads it for Level II versions. Both work with any system of 4K or more. The T-BUG discussions in this book, however, assume that you are using a Level II system with 16K of RAM or more.

## THE T-BUG ENVIRONMENT

The T-BUG monitor/debugging tool is, itself, a machine-language program. It is loaded into the TRS-80 via the SYSTEM command, and any memory space you have protected with MEMORY SIZE? is not relevant—T-BUG loads right through anything in its path.

### Loading T-BUG From Its Tape

Set up the T-BUG tape for loading; then, from BASIC, type and ENTER the command SYSTEM. This will bring up the \*? prompt symbols. Respond to them by typing and ENTERing the file name TBUG. The tape will load, showing the usual flashing asterisk near the upper right-hand corner of the screen.

When T-BUG is fully loaded, you will see the \*? symbols again. Respond by typing and ENTERing the / (slash) symbol. This should bring up the prompt symbol for the T-BUG monitor: a pound sign.

T-BUG is thus ready to go. The useful work you can do from that point will be described shortly.

### Going Between T-BUG and BASIC

T-BUG users often find it necessary to go from the T-BUG monitor to BASIC, or vice versa. Sometimes this is done on purpose; often it happens accidentally.

Going from T-BUG to BASIC can be useful at times. As you will see in Chapter 8 it is possible to write very limited BASIC programs for working the two in conjunction with one another. Other times a mistake in the machine-language program written through T-BUG will cause a "crash" that returns the system to BASIC (as indicated by the READY message and less-than prompt symbol).

There are two basic ways to get out of the resident T-BUG monitor and back to BASIC. One way is to depress the RESET push button on the back of the TRS-80 keyboard unit. Another way to accomplish the transfer is by answering the T-BUG prompt symbol by typing J0000 (it is not necessary to strike the ENTER key at the end of this entry).

In either case the system is returned to the BASIC monitor without disturbing the T-BUG programming in the least.

But once you get to BASIC from T-BUG you will probably want to get back to T-BUG again. To get from BASIC to T-BUG (assuming T-BUG is still resident from a previous tape-loading operation), simply type and ENTER the SYSTEM command, and then respond to the \*? characters by typing and ENTERing /17312. That will bring up the T-BUG monitor's prompt symbol—a pound sign. Then you're back in business with T-BUG.

There are two circumstances, however, that prevent a return to T-BUG from BASIC. One is a catastrophic "blowup" of the program at the T-BUG level. Such a catastrophe generally loads the screen with a lot of meaningless characters and often ends up displaying the MEMORY SIZE? message. In such a case the T-BUG monitor, itself, is messed up, and trying to get back to it by doing a /17312 from the SYSTEM mode will not improve matters at all.

A second set of circumstances that will prevent a return from BASIC to T-BUG is one where you write some BASIC programs into the memory that is normally allocated for the T-BUG monitor. As described in Chapter 8 it is possible to write short BASIC programs of a certain nature without disturbing the T-BUG monitor, but you have to know what you are doing. Generally speaking, writing BASIC programs while the T-BUG monitor is loaded will mess up T-BUG; and doing a /17312 from SYSTEM will cause a "crash."

Naturally, turning the TRS-80 power switch off and on will kill the T-BUG monitor, too. But perhaps this should go without having to say it.

For the time being, at least, avoid going from T-BUG to BASIC and writing some BASIC programs. And if a T-BUG generated program should blow up and the system returns to BASIC, doing a /17312 from SYSTEM might get you back into T-BUG quite nicely. But if T-BUG doesn't come back or it begins doing strange things for you, you will have to load the T-BUG monitor from the tape again.

#### SUMMARY

To get from T-BUG to BASIC:

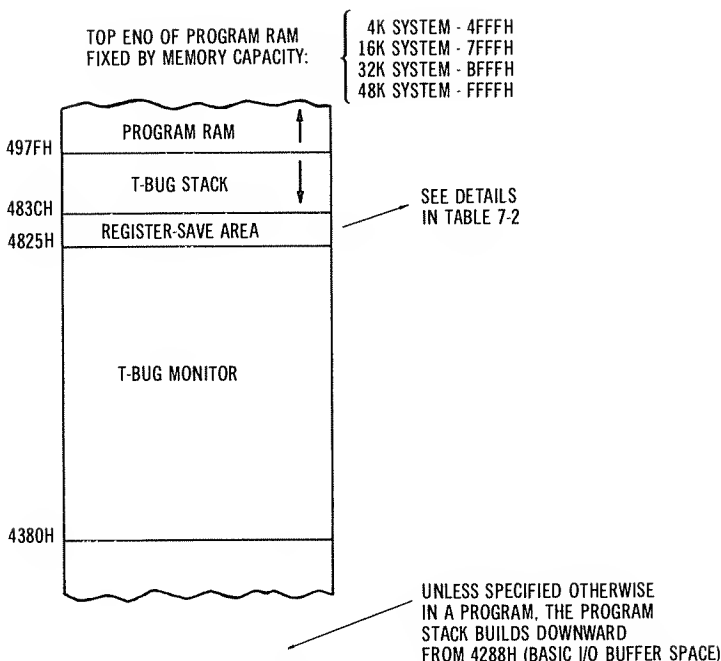
- a. Work the RESET push button
- or*
- b. Answer the # prompt symbol with a J0000.

To get from BASIC to T-BUG:

>SYSTEM	<do ENTER>
*?/17312	<do ENTER>

#### Organization of the T-BUG Monitor

The T-BUG monitor occupies the user's memory space, between 4380H and 497FH (decimal 17280 through 18815). Comparing these figures with a Level 11 memory map, you will find T-BUG is resident some 50 bytes above the beginning of the space that is normally allocated for BASIC program text. Then it runs 15360



**Fig. 7-1. General memory map of T-BUG monitor.**

bytes upward from there. See the T-BUG monitor's memory map in Fig. 7-1.

That entire memory space for T-BUG is not wholly occupied with the machine-language programming for it. Near the top of that space, locations 4825H through 483CH, is the so-called register-save area. These 18 addresses are used for saving the content of the Z-80's 16 working registers whenever a T-BUG operation does things calling for saving them on a temporary basis. The two additional addresses are not used, appearing at either end of the register-save area.

The memory space at the very top of the T-BUG's memory map is devoted to its own stack. When T-BUG is loaded into the TRS-80 it automatically sets the monitor's stack pointer to address 4980H. And when the stack is used during the running of the T-BUG monitor it builds downward toward the top end of the register-save area.

Incidentally, the T-BUG monitor stack and the SYSTEM stack are two entirely different things. The monitor stack, just described, is used exclusively by the monitor; you, the programmer, should not attempt to work with it directly.

The SYSTEM stack, on the other hand, is the one available to the programmer during the execution of a T-BUG generated, machine-language program. This is the one that builds downward when your program uses instructions such as nested CALLs, PUSHes, and POPs.

The T-BUG monitor automatically places the SYSTEM stack at 4288H—right smack in the TRS-80's i/o buffer for BASIC operations. This is one reason why you must be careful about trying to use BASIC and a SYSTEM-oriented machine language at the same time. Again, Chapter 8 will show how you can get out of the situation by relocating the SYSTEM stack: getting it out of the i/o-buffer space.

## T-BUG OPERATIONS

Most introductory discussions of the T-BUG monitor describe the eight available commands as separate identities. Most T-BUG users, especially beginners, are far more interested in knowing how to do things than brief descriptions of the commands. So the following discussion tells how to do things with the T-BUG monitor, leaving basic command definitions to a summary on p. 140.

### Examining Memory Locations

It is possible to examine the contents of all the TRS-80 memory locations from BASIC. Something such as PRINT PEEK (*address*) will do the job, where *address* is a decimal number within the system's memory map. The result is another decimal number that represents the 1-byte contents of that address.

You can do the same thing much easier, and in a hexadecimal format, from the T-BUG monitor. Just answer the # prompt symbol by striking the *M* key, followed by four hexadecimal characters representing the address location. The content of that address, in a hexadecimal format, will then appear to the right of the number you just entered. (You will not have to strike the ENTER key. The ENTER key is rarely needed for entering information in T-BUG.)

Suppose you want to examine the contents of memory location 0000H. Answer the # prompt symbol by striking keys M0000. Striking the *M* key puts the monitor into its Memory mode, and striking the zero key four times in succession specifies address 0000. Do that, and you will see this sort of display on the screen:

```
# M 0000 F3
```

The machine has printed the pound sign, F3, and spaces. The main idea, however, is that the content of memory address 0000H is F3 (for whatever that might be worth to you).

Now strike the ENTER key, and the display will look like this:

```
# M 0000 F3  
0001 AF
```

The second line is showing address 0001H and its contents—AF. Strike the ENTER key several more times in succession, and the display grows in this fashion:

```
# M 0000 F3  
0001 AF  
0002 C3  
0003 74  
0004 06  
0005 C3
```

What you are doing here is examining the contents of successive memory locations, beginning with the one specified after striking the *M* key. The project can run indefinitely this way; if your hand holds up long enough, you can theoretically examine the entire ROM and RAM in your TRS-80.

If you want to skip to another place in memory and examine the contents from that point, you must first get out of the memory mode and back to the monitor. To get to the monitor, strike the *X* key, instead of the ENTER key. That will bring up the pound-sign prompt symbol again. Then strike the *M* key, followed by the four-digit hexadecimal address of the location you want to examine. Looking at the contents of successively higher locations is a simple matter of striking the ENTER key for each one.

Incidentally, you might notice that the backspace (left-arrow) key doesn't work as it does in BASIC. If you strike a wrong key while specifying an address for examining memory locations, striking the backspace key doesn't change anything. The only way to correct a typing error at this point is to strike the *X* key to return to the T-BUG's monitor mode, and then start over by striking the *M* key and, hopefully, the correct four-place hexadecimal address.

Here are a few notes that summarize the procedure for examining the contents of any memory location in the TRS-80:

1. Answering the # prompt symbol by striking the *M* key puts T-BUG into its *Memory* mode.
2. The *Memory* mode can be entered only from the monitor mode. To get to the monitor mode from any other command mode, including *Memory*, strike the *X* key.
3. To specify the starting address of a memory-examining sequence, strike the *M* key, followed by a 4-byte hexadecimal address:

Maaaa

where the *a* terms represent the four address bytes or hexadecimal characters.

4. To view the contents of successive memory locations, strike the ENTER key for each one.
5. Escape from the *Memory* mode and return to the monitor mode at any time by striking the *X* key.
6. Correct address-entry errors by returning to the monitor mode (striking the *X* key) and starting over again.

### **Altering the Contents of RAM Locations**

From BASIC it is possible to alter the contents of any RAM location by doing something such as *POKE address, data*, where *address* is the decimal-format address of the byte to be changed and *data* is the byte to be entered into that address. The *data* must also be in a decimal format.

In T-BUG the content of any RAM address can be changed from the *Memory* mode just described. Your thinking of this sort ought to be limited to address locations above those of the T-BUG monitor: from 4980H to the top of the RAM space for your system. Altering the contents of memory locations in the T-BUG monitor or below risks some program catastrophes.

At any rate, get the system into the T-BUG *Memory* mode by striking the *M* key and typing the address of the location to be altered. The system will respond as though you want to simply examine the contents of that location.

Instead of simply striking the ENTER key to view the contents of the next-higher address, strike two keys that specify the one byte of hexadecimal data you want to enter into that address. If you have thus typed some valid data, the system will print the new data to the right of the old data, and then automatically examine the contents of the next memory location for you.

Here is an example:

```
# M 4A00 2E CD
4A01 FE
```

In this example the machine printed the pound sign to indicate the system was in the monitor mode. The user struck the *M* key to get into the *Memory* mode, then struck the 4A00 keys in succession to specify the starting address. The system then inserted the spaces and printed the current contents of address location 4A00. In this example the contents happen to be 2E, but you might see any two-character code when you try this for yourself.

At that point the display looked like this:

```
# M 4A00 2E
```



But then the user struck the *CD* keys to change the contents of that address from 2E to CD. The system responded by inserting a space and the revised contents of that same memory location. After that the system automatically displayed the next-higher address and its current data. There is no need to strike the ENTER key after altering the contents of a memory location—the T-BUG does that for you.

Now, if you want to alter the contents of address 4A01, just strike the two keys representing the new contents. The system will print your new data to the right of the old (to the right of the FE in this case) and jump down to display the next address location and its current contents.

But if you want to leave the contents of address 4A01 unchanged, strike the ENTER key to get to the next address location.

This is how machine-language programs are written from T-BUG. They are written from the *Memory* mode, examining successive memory locations and altering them as necessary to build up the program you want.

The following is a short programming sequence that loads this set of Z-80 instructions:

4A00	CD C9 01	CALL 01C9H	;CALL C1S FUNCTION
4A03	C3 03 4A	JP 4A03H	;LOOP TO SELF

```
# M 4A00 01 CD
4A02 FF C9
4A03 5E 01
4A04 88 C3
4A05 1A 03
4A06 09 4A
4A01 21
```

Answering the # prompt symbol with M4A00 put the system into the *Memory* mode, beginning at address 4A00H. The current content was 01, but I wanted it to be CD, the first instruction in the program. So I struck the *CD* keys, and the system automatically displayed the contents of address 4A01—a 21. I wanted a C9 in that location, so I struck those keys. I continued the process until the last instruction (4A) had been entered into address location 4A05. That completed the programming phase of the operation. It marked the end of my 6-byte program.

To see whether or not that program has been entered as expected, you can examine the six locations again from the *Memory* mode. But in order to get back to address 4A00 to do that sort of job, you must first exit the ongoing scheme by striking the X key, then answering the # prompt with an M and address 4A00. By doing

that and striking the ENTER key several times in succession, the display will look like this:

```
# M 4A00 CD
4A01 C9
4A02 01
4A03 C3
4A04 03
4A05 4A
```

Sure enough, this *Memory-examining* operation confirms that the 6-byte, machine-language program has been entered in addresses 4A00 through 4A05.

In summary, to write a machine-language program from T-BUG:

1. Get into the *Memory* mode by striking the *M* key, followed by the address (in hexadecimal) of the starting point of your program.
2. Respond to the machine's printing of the address and its contents by striking the two keys representing the data you want to enter at that address.
3. If you make a mistake in specifying the data, strike the *X* key to get out of the current *Memory* mode, then reenter the *Memory* mode by striking the *M* key and specifying the address where the mistake occurred. Alter the data accordingly.

### **Executing Programs From T-BUG**

One of the most attractive features of T-BUG is that you can actually run the programs you write with it. The more sophisticated Editor/Assembler (described in later chapters) cannot do that. You can thus compose a machine-language program as described in the previous section and execute it on the spot, testing for any possible errors. And if you handle the execution properly, T-BUG will still be resident in your system, and you can go to the *Memory* mode to correct any errors you discover.

But just as a person ought to learn how to stop a car before learning how to make it go, a T-BUG user ought to know how to stop the execution of a machine-language program before trying to run one. Without knowing how to "put on the brakes," the program can zip through its end and wander into sections of memory where there is no rational programming, thus crashing the whole thing.

When running the T-BUG monitor, the least troublesome way to conclude a program is by inserting a *Break* at the address marking the point where you want the program to stop its execution. Suppose, for instance, you want a program to run between addresses

4A00 and 4AF6, inclusively. Before executing that program (a procedure described shortly), you should insert a T-BUG *Break* at the address immediately following the final instruction in the program. In this case the final instruction resides at address 4AF6; so the *Break* should be located at 4AF7.

How do you specify a *Breakpoint*? Get into the T-BUG's monitor mode, as indicated by the pound-sign prompt symbol, then strike the *B* key, followed immediately by the *Breakpoint* address. To insert the *Break* at address 4AF7, the presentation on the screen looks something like this:

```
# B 4AF7
#
```

The T-BUG monitor prints the pound signs and inserts the spaces. All you have to do is answer that first pound sign with B4AF7.

Now, when you execute the program from the beginning (address 4A00 in this example), it will run until it reaches your specified *Breakpoint*—address 4AF7. At that time program control will break away from your machine-language program and return to the T-BUG monitor. The system will return the # symbol, waiting for your next instruction.

Here is an example to illustrate the point: get into the Memory mode, and load a program that turns out looking like this when you examine it . . .

```
# M 4A00 CD
4A01 C9
4A02 01
4A03 21
4A04 20
4A05 3C
4A06 36
4A07 58
```

Strike the *X* key to get into the monitor mode, then strike B4A08 to place a *Break* at address 4A08—the address following the end of the program residing in locations 4A00 through 4A07.

The main point of this particular discussion is the placing of the breakpoint. But to test its effectiveness you ought to execute the program. Do that by answering the # with J4A00.

Immediately after typing that particular sequence of characters, the screen will clear, an *X* character will appear near the middle of the top line on the screen, and the T-BUG's prompt symbol will appear at the upper left-hand corner of the screen.

If things look this way on your screen, you know that the program has been executed to the breakpoint and that it returned to the T-BUG's monitor for you.

The program in this example is responsible for clearing the screen and printing the X character. T-BUG executed those instructions when you typed in J4A00, then the *Break* set at address 4A08 marked the end of the execution and returned the system to its monitor mode.

The breakpoint will remain in that same place until you either write some new instructions into location 4A08 or do a *Fix* operation (an operation described a bit later).

In the process of illustrating the use of T-BUG's *Break* function you have also initiated the execution of a program. That was the J4A00 keyboard operation.

So the simplest way to begin execution of a program from T-BUG is by answering the monitor's # symbol with a J, followed by the starting address, in a hexadecimal format, of course. To execute the previous demonstration program, for instance, you did a J4A00—a *Jump* to the starting address of the program.

#### SUMMARY

Jsss begins execution of a program at address ssss.

Beee specifies the end, or breakpoint, of a program at address eeee.

The breakpoint must be specified *before* the program is executed, and it must be specified at the address following the last instruction to be executed.

The system automatically returns to the T-BUG monitor on finding the breakpoint.

Actually, the *Break* feature of T-BUG is intended to be a debugging tool and not an endpoint marker for all machine-language programs. Machine-language programs are normally self-contained; they do not break away to some other monitor such as BASIC or T-BUG. Finished machine-language programs, in other words, have loops that keep them running indefinitely.

Machine-language programs that run as endless loops do not have to conclude with a *Break* function in T-BUG. A *Break*, under those conditions, is really irrelevant as far as marking the end of the program is concerned—the program has no end. The very nature of endlessly looping machine-language programs prevents operations from zooming into some unspecified memory space.

It is possible and, indeed, a good idea to *Break* machine-language loops for troubleshooting purposes; but what if you want to run the program in its normal fashion? What if you want to run the

program without *Breaking* it? And, finally, how do you get back to the T-BUG monitor from an endlessly looping machine-language program?

Well, the only way to get out of a machine-language loop is by working the RESET push button on the back of the TRS-80 keyboard unit. When you are running a machine-language program in the SYSTEM mode, working the RESET push button immediately breaks up the program and returns control to the BASIC monitor.

Working the RESET push button while running programs in T-BUG also returns control to BASIC. As described earlier in this chapter you can then get back to the T-BUG monitor by entering SYSTEM, followed by /17312.

Here is a program that enters an endless loop:

```
# M 4A00 CD
4A01 C9
4A02 01
4A03 21
4A04 20
4A05 3C
4A06 36
4A07 58
4A08 C3
4A09 08
4A0A 4A
```

You can specify a *Break* at the next-higher address, 4A0B, but it won't do any good. The last three bytes of the program cause it to loop around to itself, effectively creating a tight program loop that can run endlessly. There is no way the program can jump out to any unspecified memory locations. In short, the program will never see any *Break* that might be specified at address 4A0B.

Execute that program from its beginning by doing a J4A00. You will see the X character printed near the middle of the top line on the screen, but this time you won't see the T-BUG's pound-sign prompt character. Also notice that the keyboard is "dead." Nothing that you try at the keyboard, including striking the BREAK key, has any effect whatsoever.

The program is tied up in an endless loop. The only way out of the situation is to work the RESET push button. That will return the system to BASIC; but then you can get back to the T-BUG monitor by entering SYSTEM and answering the \*? symbols by entering /17312. The appearance of the pound sign near the upper left-hand corner of the screen confirms a return to T-BUG. Now you can run the program again, modify it, write a new program, or do whatever you want from T-BUG.

## Debugging Programs From T-BUG

So far in this chapter you've seen how to load the T-BUG monitor, examine the contents of memory locations, alter the contents to write a program, execute a program, and get out of a program that is being executed from T-BUG. These principles are adequate for writing just about any kind of machine-language program, but they assume any troubles are fairly easy to find and remedy. Unfortunately, machine-language bugs are not always so easy to locate; that's where the debugging power of T-BUG really pays off.

One debugging tool already described is the *Break* operation. You can insert a *Break* at the end of any complete instruction sequence in a machine-language program. Then, when you execute the program, it will run until it encounters your specified *Break* address. At that time the system returns to the T-BUG's monitor mode, and you are free to examine or alter the contents of some memory locations. More importantly in many instances, you can also examine the content of the Z-80's internal registers. More about that in a moment.

The main point is that you can stop and freeze a machine-language program at any point specified by the *Break* operation, thus giving you an opportunity to check out the program's current status.

To get rid of a *Break* you have specified at some earlier time, simply get into the T-BUG's monitor mode and strike the *F* key. This does a "Fix" on the *Breakpoint*, restoring any legitimate machine-language instructions that should normally appear there.

T-BUG can handle just one *Break* operation at a time. If you specify more than one *Breakpoint*, only the last one will work, and you will end up with meaningless instructions inserted at the memory locations for previously specified *Breaks*.

### IMPORTANT

After specifying a *Break* at one point in a machine-language program, always do a *Fix* before specifying a different *Breakpoint*.

You don't have to keep track of the address specified for an earlier *Break* operation. Since the system is supposed to handle just one *Breakpoint* at a time, it keeps track of the address in the T-BUG stack. Doing a *Fix* (striking the *F* key) automatically calls up the *Breakpoint* address and inserts the legitimate instructions it replaced when you originally specified the *Break*.

As implied earlier in this discussion, one of the good reasons for

inserting a *Break* into a machine-language program is to give you a chance to examine the contents of the Z-80's internal registers. On encountering the *Break* the system will leave the machine-language program and return to the T-BUG monitor mode, as indicated by the # prompt symbol. To view the contents of the Z-80's internal registers, simply respond to the # by striking the *R* (*Register*) key.

On striking the *R* key, a set of 12 hexadecimal figures pop onto the screen. Those figures represent the contents of the 12 main Z-80 registers as shown in Fig. 7-2.

A'	F'
D'	E'

B'	C'
H'	L'

A	F
D	E

B	C
H	L

IX (MSB)	IX (LSB)
SP (MSB)	SP (LSB)

IY (MSB)	IY (LSB)
PC (MSB)	PC (LSB)

**Fig. 7-2. Screen format for Register operation.**

After displaying the contents of the registers in this particular fashion, the system returns to the T-BUG monitor mode, awaiting further instructions from you.

The following demonstration is a rather extensive one, but it walks you through just about all of the T-BUG principles described thus far.

The assembly-language version of the program for this demonstration appears in Table 7-1. It is basically a time-delay program. On running it you should see a white rectangle printed immediately near the upper left-hand corner of the screen. Then after a brief time delay (about  $\frac{3}{4}$  of a second) a second white rectangle appears near the middle of the top line on the screen.

The first three instructions initialize the program, clearing the screen with a CALL 01C9H, loading the rectangle character code

**Table 7-1. Assembly-Language Version of the Demonstration Program**

Address	Contents	Mnemonics	Comment
4A00	CD C9 01	START: CALL 01C9H	;CLEAR THE SCREEN
4A03	3E 8F	LD A,8FH	;CHARACTER TO ACC
4A05	32 00 3C	LD (3C00H),A	;PRINT FIRST SPOT
4A08	06 FF	LD B,0FFH	;SET B TO FF
4A0A	0E FF	SETC: LD C,0FFH	;SET C TO FF
4A0C	0D	DECC: DEC C	;COUNTDOWN C
4A0D	20 FD	JR NZ,DECC	;IF NOT DONE, COUNT AGAIN
4A0F	05	DEC 8	;COUNTDOWN 8
4A10	20 F8	JR NZ,SETC	;IF NOT DONE, START AGAIN
4A12	32 20 3C	LD (3C20H),A	;PRINT SECOND SPOT
4A15	18 FE	SELF: JR SELF	;JUMP TO SELF

(BF) into the A register, and printing that character at video memory location 3C00H.

The next six lines, addresses 4A08 through 4A11, perform the short time delay. Those lines simply count down the BC register pair from FFFF to 0000, where the B register is the lsb for the countdown operation.

The penultimate instruction prints the second white rectangle at video memory location 3C20. And the last line simply “buzzes” to itself, thus creating an endless loop at the end of the program.

Load the program from T-BUG, beginning at address 4A00. After that, examining the program should turn up this sort of display:

```
# M 4A00 CD
4A01 C9
4A02 01
4A03 3E
4A04 BF
4A05 32
4A06 00
4A07 3C
4A08 06
4A09 FF
4A0A 0E
4A0B FF
4A0C 0D
4A0D 20
4A0E FD
4A0F 05
4A10 20
4A11 F8
4A12 32
4A13 20
4A14 3C
4A15 18
4A16 FE
```



This is the 17-byte program as it appears in T-BUG. Any data in address locations above 4A16 or below 4A00 is not relevant.

Get back into the T-BUG monitor mode by striking the X key, then run the program by doing a J4A00. This starts the program from the very beginning.

If all is going well, it should print that white rectangle near the upper left-hand corner of the screen, do a  $\frac{3}{4}$ -second time delay, and then print the second rectangle near the middle of the top line on the screen. The T-BUG prompt symbol should *not* appear when the program is done, for the simple reason that the program, as presented here, is never done. The looping statement at the end of the program is an endless one.

The only effective way to break out of the program (without destroying it) and return to the T-BUG monitor is by working the RESET push button and then entering SYSTEM, followed by entering /17312.

For testing purposes, the matter of RESETing to BASIC and then calling SYSTEM again is a very troublesome affair. To make things easier, specify a T-BUG Break at the beginning of the final looping instruction in the program—at address 4A15.

After inserting the Break at that point, execute the entire program with a J4A00. The program will run as before, but it automatically Breaks to the T-BUG monitor after printing the second rectangle. The first rectangle will be cleared from the screen, but only because the monitor does that clearing operation every time it is called. The fact that the first rectangle figure disappears has nothing to do with your program, itself.

To get rid of the Break at the end of the program. Strike the F key to do a Fix. Now, when you run the program with a J4A00, it will run and enter the endless loop again. Yes, you have to do a RESET, SYSTEM, /17312 to get back to T-BUG again.

Referring to the assembly-language version of the program in Table 7-1, you can see that there is nothing but register-loading operations taking place between the instructions in locations 4A03 and 4A0A. So by the time the program reaches the instruction in 4A0C, register A should have a BF in it, while registers B and C should be loaded with FF.

To test the operation of the program to that point, specify a Break at 4A0C, and then do a J4A00. The program will execute right up to address 4A0C, then return to the T-BUG monitor. Then, to see the contents of the registers, strike the R key.

Comparing the register display with the format shown in Fig. 7-2, you should see a BF in the A register and FFFF in the BC register pair. The content of the other registers shown on the screen is not relevant at this point.

If anything has been wrong with the initialization phase of the program, you would be able to spot the trouble by examining the registers at this place in the program. That being the case, you would make the necessary changes in the program and try the operation again.

Do a *Fix* to remove the *Break* when you are ready to proceed to the next phase of this demonstration. If you wish, you can double-check the overall operation by running it from the start again—by doing a J4A00.

Now, here is something rather tricky, but meaningful: insert a *Break* at 4A0D. This will stop the program just after the C register is decremented at 4A0C. Do a J4A00; and when the system returns to the T-BUG monitor, do an *R* to view the contents of the registers. Now the C register should show an FE, instead of an FF. Why? Because the *Break* is located after the point in the program that calls for decrementing the C register by 1. It has counted down from FF to FE.

Here's the fun part. Do a J4A0C. That will rerun the program from the point where the C register will be decremented again. When the system returns to the monitor mode (the *Break* is still set at 4A0D), examining the registers will show that the C register has counted down again—counted down to FD.

Run through that sequence several times: doing a J4A0C and examining the content of the C register. Each time you do that, the number in that register should get one unit smaller. In effect, you are watching the C-register down-counting operation. If you do a J4A00—start the program from the beginning—you will find the C register starting again from FE.

Do a *Fix* to get rid of the current breakpoint, and set a new *Break* at 4A12. Run the whole program by doing a J4A00; and when the system breaks to the T-BUG monitor, strike the *R* key to look at the registers. Can you explain why the BC register pair now contains 0000?

Doing a *Fix* gets everything back to normal.

Through all of this, you should be getting at least an inkling of how T-BUG can be used for composing, debugging, and running machine-language programs.

### **Altering the Contents of the Registers**

A more subtle debugging tool involves altering the contents of the Z-80's registers at some point during the execution of a program. You have already seen how it is possible to interrupt the execution of a program with a *Break* function and then view the contents of the registers by doing a *Register* function.

Before running the program again, it is sometimes useful to alter

the contents of the registers first. Recall that the contents of the Z-80's registers are saved in the T-BUG's register-save area. Every time the system is returned to the monitor mode, the current contents of the Z-80's registers are saved in that area—and that's what you see displayed on the screen whenever you do the Register function.

Fig. 7-2 outlines the memory map for the register-save area in greater detail than the general memory map in Fig. 7-1 does.

**Table 7-2. Register-Save Area Addresses**

Register	Address	Register	Address
A	482EH	A'	4826H
B	4830H	B'	4828H
C	482FH	C'	4827H
D	4832H	D'	482AH
E	4831H	E'	4829H
H	4834H	H'	482CH
L	4833H	L'	4828H
IX (msb)	4836H	IY (msb)	4838H
IX (lsb)	4837H	IY (lsb)	3837H
SP (msb)	483AH	PC (msb)	483CH
SP (lsb)	4839H	PC (lsb)	483BH

Suppose, for some reason, you have interrupted the execution of a machine-language program and you want to set the HL register pair to BF10, where BF is the msb. That can be done as though you were writing a program at the addresses reserved for the HL pair in the register-save area. To get the BF into the H register, do an M4834 (note from Fig. 7-2 that 4834H is the address of the H register-save area). And after the system prints the current contents of 4834H, strike the BF keys. After that, the system will bring up the contents of the next address location, 4835. But you want to access 4833—the register-save location for the L register. So strike the X key to get back to the monitor, do an M4833, followed by a 10.

Now, when you do a Register function, you will see that the HL pair is set to BF10.

Of course, it would be easier to access the lower-addressed L-register location first, then let the system access the higher H-register location for you. But the main point is that you are free to alter the contents of the registers, using the ordinary T-BUG addressing and data-changing technique. The addresses of the register-save area, specified in Fig. 7-2, are your guide.

One of the most compelling reasons for tinkering around in the register-save area is to set the program counter (pc) registers to some specified point. Such an operation is indispensable when you

want to resume the execution of a program without resorting to a Jump function.

The pc register in the Z-80 carries the address of the next instruction to be executed. Thus, tinkering with the pc register's contents, it is possible to specify some new starting point in the program—then use the Go function to begin execution.

The Go function is called by striking the G key, but it is offered in the Radio Shack literature with some clearly stated reservations: You must know what you are doing to use it. You must know where the pc register is pointing at the time. That you can determine by examining the pc register via a Register function; and if you want to alter the starting point, you can use the Memory function for the pc register addresses to do the job. Then you can execute the operations by striking the G key.

### **Saving T-BUG-Generated Programs On Tape**

Machine-language programs generated by T-BUG functions can be saved on cassette tape by doing:

```
P bbbb pppp eeee name
```

where

*bbbb* is the beginning address—the lowest address of the machine-language program,

*eeee* is the highest-used address in the program,

*pppp* is the entry point for the machine-language program—usually, but not necessarily, the same address as the beginning address,

*name* is the file name, composed of one to six alphanumeric characters.

To save the program in Table 7-1 under the file name of DELAY, for example, set the recorder to its RECORD mode and strike the following keys:

```
P4A004A164A00DELAY
```

and strike the ENTER key. The system will respond to this Punch operation by inserting the appropriate spaces for you. The display on the screen will thus look something like this:

```
# P 4A00 4A16 4A00 DELAY
```

The recorder will begin running as soon as you terminate the instruction with an ENTER. If your file name happens to have six or more characters, the recorder will start after you strike the key for

the sixth character in the file name—you won't have to strike the ENTER key in that case, and any characters beyond the sixth one in the file name will be ignored.

If, during the process of entering the *Punch* sequence, you decide something is wrong with it, you can abort it by striking the *X* key. But if you've got as far as the file name, you can abort the *Punch* operation only by striking the *BREAK* key. (The system would otherwise think the *X* is part of the file name.)

To load a machine-language program previously saved from the T-BUG monitor, set the recorder to the beginning of the program, set it for the *PLAY* mode, and strike the *L* (*Load*) key. The usual flashing asterisk indicates the loading operation is taking place.

It is also possible to work with machine-language programs not specifically generated via T-BUG, but such programs must fit into memory spaces that aren't occupied by any portion of the T-BUG monitor (between 4380H and 497FH). Any *SYSTEM*-oriented tape can be resident with the T-BUG monitor as long as it loads above 497FH.

### SUMMARY OF T-BUG OPERATIONS

**Memory (*M* key)**—*M aaaa* displays address *aaaa* and its contents in a hexadecimal format. The data at that address can be altered by striking the two appropriate keys.

**Jump (*J* key)**—*J aaaa* begins the execution of a machine-language program at address *aaaa*.

**Break (*B* key)**—*B aaaa* sets the breakpoint at address *aaaa*.

**Fix (*F* key)**—Striking the *J* key restores the instructions that were set aside by the last-specified *Break* function.

**Register (*R* key)**—Striking the *R* key displays the contents of the Z-80's internal registers as shown in Fig. 7-1.

**Go (*G* key)**—Striking the *G* key begins execution of a program at the address contained in the Z-80's *pc* registers.

**Punch (*P* key)**—*P aaaa pppp eeee name* saves a machine-language program on cassette tape. The program begins at *aaaa*, ends at *pppp*, has an entry point specified by *eeee*, and goes by the file name *name*.

**Load (*L* key)**—Striking the *L* key loads a tape version of a machine-language program that was previously saved by a *Punch* operation.

## CHAPTER 8

# Exploring the TRS-80 With T-BUG

Using the T-BUG monitor brings up the possibility of working with the TRS-80/Z-80 system in a hexadecimal format—a scheme that is much easier to use than BASIC and its decimal numbers for machine-language programming. This chapter considers the TRS-80 video, keyboard, and user memory from a hexadecimal point of view.

It is assumed that you already understand some of the fundamentals of these topics from a decimal viewpoint, as presented in Chapters 2, 3, and 4. Here, then, it is possible to concentrate on T-BUG and hexadecimal notation without having to cover a lot of material about the structure of the video, keyboard, and user memory.

### THE HEXADECIMAL VIDEO ENVIRONMENT

The TRS-80's video memory occupies RAM addresses 3C00H through 3FFFH. Each of those addresses contains a byte of data that represents an alphanumeric or TRS-80 graphics code, and each address “points” to a well-defined position on the crt screen.

Fig. 8-1 shows the addresses in the video memory line up with endpoint positions on the screen. Address 3C00H, for example, corresponds to the first character position on the top line; 3C3FH corresponds to the last character position on that top line; and 3FFFH represents the last character position on the last line.

There are 40H character positions on each line; so if you want to position a character near the middle of a given line, you can figure it is about 20H positions to the right of the first character position. What, then, is the address in video memory for a character

ADDRESS	LINE	ADDRESS
3C00H	----->	3C3FH
3C40H	----->	3C7FH
3C80H	----->	3CBFH
3CC0H	----->	3CFFH
3D00H	----->	3D3FH
3D40H	----->	3D7FH
3D80H	----->	3DBFH
3DC0H	----->	3DFFH
3E00H	----->	3E3FH
3E40H	----->	3E7FH
3E80H	----->	3EBFH
3EC0H	----->	3EFFH
3F00H	----->	3F3FH
3F40H	----->	3F7FH
3F80H	----->	3FBFH
3FC0H	----->	3FFFH

**Fig. 8-1. Hexadecimal video memory/screen format.**

to be positioned near the middle of the screen? Well, a line near the middle of the screen begins with address 3E00H, and moving 20H positions to the right of 3E00H gives us 3E20H. Try it yourself with Program 8-1.

**Program 8-1. Assembly language for placing a rectangle of light near the middle of the video screen/memory.**

```

4A00 CD C9 01    START:    CALL 01C9H    ;CLEAR THE SCREEN
4A03 21 20 3E        LD HL,3E20H    ;SET THE VIDEO POINTER
4A06 36 BF        LD(HL),191D    ;PRINT A CHARACTER
4A08 18 FE        DONE:    JR DONE    ;LOOP TO SELF

```

Since Program 8-1 concludes with a “loop-to-self” operation, you will have to work the RESET push button to get out of it.

If you want to explore the video memory on your own, comparing the positions indicated in Fig. 8-1 with your own plots, just change the VIDEO POINTER address for the HL pair in the second line of Program 8-1. Use any valid video address between 3C00H and 3FFFH.

Program 8-2 demonstrates the addressing makeup of the video memory/display in a more dynamic fashion. The first line sets the starting point of the video memory/display, and the second line marks the end of it. Load and run the program from T-BUG and you will see the first group of four lines filled with 0s, the second group of four filled with 1s, the third group with 2s, and the bottom group of four lines filled with 3s.

Besides showing that the video memory starts at 3C00H and ends at 3FFFH, Program 8-2 shows how the lsb of the 2-byte memory address cycles between 00H and FFH four different times through the video memory/display.

### Program 8-2. Assembly listing for video memory demonstration.

```

4A00 21 00 3C          LD HL,3C00H      ;POINT TO START OF VIDEO
4A03 01 FF 3F          LD BC,3FFFH      ;POINT TO END OF VIDEO
4A06 16 30             LD D,48D        ;SET FIRST CHARACTER
4A08 72                LD(HL),D        ;PRINT CHARACTER
4A09 79                LD A,C          ;FETCH LSB OF VIDEO
4A0A BD                CP L            ;END OF QUADRANT?
4A0B 28 03             JR Z,MSB        ;IF SO,DO MSB TEST
4A0D 23                NXT:            INC HL      ;SET NEXT VIDEO POS.
4A0E 18 F8             JR AGN          ;AND JUMP TO AGN
4A10 78                MSB:            LD A,B      ;FETCH MSB OF VIDEO
4A11 BC                CP H            ;END OF VIDEO?
4A12 28 03             JR Z,DUN        ;IF SO,JUMP TO DUN
4A14 14                INC D           ;ELSE SET NEW CHAR.
4A15 18 F6             JR NXT          ;AND DO AGAIN
4A17 18 FE                DUN:         JR DUN      ;LOOP TO SELF

```

Then try the program specified here as Program 8-3. This one is an example of some machine-language animation graphics. It generates a small white rectangle that bounces back and forth across the middle of the screen.

The first line of Program 8-3 clears the screen. The second line sets the initial position of the rectangle at 3E01H—to a position

### Program 8-3. Assembly listing for a bouncing rectangle of light.

```

4A00 CD C9 01          START:          CALL 01C9H      ;CLEAR THE SCREEN
4A03 21 01 3E          LD HL,3E01H      ;SET INITIAL POINT
4A06 11 00 3E          LD DE,3E00H      ;SET 'OLD' POSITION
4A09 06 FF             LD B,0FFH        ;SET RIGHT PHASE
4A0B 3E 20             PLOT:            LD A,32D      ;SET ERASE CHARACTER
4A0D 12                LD(DE),A        ;ERASE OLD POINT
4A0E 5D                LD E,L          ;SAVE NEW POINT
4A0F 36 BF             LD(HL),191D      ;PLOT NEW POINT
4A11 D5                TIME:            PUSH DE      ;SAVE 'OLD' POSITION
4A12 16 08             LD D,08H        ;SET MSB OF DELAY
4A14 1E FF             SETE:            LD E,0FFH     ;SET LSB OF DELAY
4A16 1D                DECE:            DEC E        ;COUNTDOWN LSB
4A17 20 FD             JR NZ,DECE       ;IF NOT DONE,COUNT AGAIN
4A19 15                DEC D            ;ELSE COUNTDOWN MSB
4A1A 20 F8             JR NZ,SETE       ;IF NOT DONE,COUNT MORE
4A1C D1                POP DE          ;ELSE FETCH 'OLD' POS.
4A1D CB 40             RUN:             BIT 0,B      ;RIGHT PHASE?
4A1F 28 0E             JR Z,LTLIM      ;IF NOT,CHECK LEFT LIMIT
4A21 3E 2F             LD A,3FH        ;GET RIGHT LIMIT
4A23 BD                CP L            ;AT RIGHT LIMIT?
4A24 20 U6             JR NZ,INC       ;IF NOT,THEN INCREMENT
4A26 06 00             LD B,0          ;ELSE SET FOR LEFT PHASE
4A28 2D                DEC:            DEC L        ;DECREMENT TO LEFT
4A29 C3 0B 4A          JP PLOT        ;PLOT NEW POINT
4A2C 2C                INC:            INC L        ;INCREMENT TO RIGHT
4A2D 18 FA             JR PLOT        ;PLOI NEW POINI
4A2F AF                LTLIM:          XOR A        ;SEI LEFT LIMIT
4A30 BD                CP L            ;AT LEFT LIMIT?
4A31 20 F5             JR NZ,DEC       ;IF NOT,DECREMENT
4A33 06 FF             LD B,0FFH      ;ELSE SET RIGHT PHASE
4A35 18 F5             JR INC          ;AND JUMP TO INCREMENT

```



near the beginning of the middle line on the screen. The HL register pair is the video pointer, where H is always at 3E, but L is incremented and decremented between 00H and FFH to make the figure move back and forth across the screen.

The ASCII character set, already described in Chapter 2 in a decimal format, applies equally well to machine-language and T-BUG operations; but, of course, the characters have to be specified in a hexadecimal form. (See Appendix C.)

The same idea applies to the special TRS-80 graphics characters. These, too, can be applied to machine-language and T-BUG programs, in a hexadecimal version, of course. (See Appendix D.)

The main difference between BASIC decimal-oriented video operations and the T-BUG hexadecimal-oriented video is the fact that the latter cannot deal with cursor-control operations in a meaningful way. Much of the material in Chapter 2 deals with cursor-control operations, including those specified with decimal numbers 0 through 31, and 192 through 255. None of those are relevant in machine-language programming. You have to write in your own cursor controls now.

If, for instance, you print a character at one position on the screen, and you want to print another character at the next position to the right, you must insert a machine-language instruction that moves the character pointer—increments it. And if you want to do a line-feed/carriage-return operation, you have to write machine-language instructions that set the character pointer (whatever register pair you are using to indicate video addressing) to the beginning of the next line.

This does not necessarily mean that every video printing or graphic operation has to be written as a tedious, step-by-step, machine-language program. The TRS-80 contains a number of ROM functions that can help you with the programming task. It is, in other words, possible to do some CALLs from machine language that will cause interesting and highly useful things to happen in the video environment. Those CALLs, however, mean a lot more when they are used in conjunction with some keyboard operations.

So the next section in this chapter takes up the subject of the keyboard environment, as viewed in a hexadecimal format. And in the process of demonstrating some of the keyboard principles, you will also uncover some useful video “tricks.”

## **THE HEXADECIMAL KEYBOARD ENVIRONMENT**

Whether working in BASIC or machine language, the TRS-80 keyboard is the system's primary input control device. The following material expands on the basic keyboard principles outlined in Chap-

ter 3, putting things into a machine-language-oriented, hexadecimal format.

## The Keyboard Matrix

Anytime a single key is depressed, at least one address in the “keyboard memory” takes on a well-defined, predictable hexadecimal number. Table 8-1 illustrates this particular keyboard principle.

**Table 8-1. The Hexadecimal Keyboard Matrix Format**

Data	Keyboard Addresses							
	3801H	3802H	3804H	3808H	3810H	3820H	3840H	3880H
01H	@	H	P	X	0	(	8	ENTER
02H	A	I	Q	Y	1	)	9	CLEAR
04H	B	J	R	Z	2	*	:	BREAK
08H	C	K	S		#	+		
					3	;		↑
10H	D	L	T		\$	<		
					4	,		↓
20H	E	M	U		%	=		
					5	—		←
40H	F	N	V		&	>		
					6	.		→
80H	F	O	W		,	?		
					7	/		SPACE

The keyboard addresses, shown along the top of the table, indicate the primary keyboard “memory” locations. The corresponding data to appear at those addresses is shown running down along the left side of the table.

So if you happen to be depressing the *H* key, you will find that address 3802H contains data 01H. Depressing the *V* key, however, turns up data 40H at address 3804H. Each key thus has a specific address/data combination assigned to it.

You can check out the table for yourself, using the program listed as Program 8-4A. That program fetches the content of address 3801H, tests for any key depression at that address, and if there is a key depression involving any of the keys under address 3801H, the program returns to the T-BUG monitor. At that time you can do a *Register* function to view the contents of the A register—the one carrying the data byte from 3801H.

## Keyboard Demonstrations

### Program 8-4A. Checking for key depression in row 3801H.

```

4A00 3A 01 38    FETCH:    LD A,(3801H)    ;FEICH BYTE FROM KB
4A03 B7          OR A      ;CHECK FOR KEY DEPRESSION
4A04 28 FA          JR Z,FETCH    ;IF NOT, FETCH AGAIN
4A06 C3 A0 43      JP 43A0H      ;ELSE RETURN TO T-BUG

```

### Program 8-4B. Checking for key depression anywhere on the keyboard.

```

4A00 21 7F 38          LD HL,387FH    ;POINT TO WHOLE KB
4A03 7E          WAIT:    LD A,(HL)    ;FETCH KB BYTE
4A04 B7          OR A      ;CHECK FOR CLEAR
4A05 20 FC          JR NZ,WAIT    ;IF NOT, WAIT AGAIN
4A07 7E          FETCH:    LD A,(HL)    ;ELSE FETCH KB BYTE
4A08 B7          OR A      ;CHECK FOR ANY KEY
4A09 28 FC          JR Z,FETCH    ;IF NOT, FETCH AGAIN
4A0B C3 A0 43      JP 43A0H      ;ELSE RETURN TO T-BUG

```

### Program 8-4C. Checking for depression of the A key.

```

4A00 3A 7F 38          WAIT:    LD A,(387FH)    ;FETCH BYTE FROM WHOLE KB
4A03 B7          OR A      ;CHECK FOR ANY KEY DEPR.
4A04 20 FA          JR NZ,WAIT    ;IF SO,WAIT AGAIN
4A06 3A 01 38          FETCH:    LD A,(3801H)    ;GET BYTE FROM A ROW
4A09 FE 02          CP 2H        ;IS IT 'A'?
4A0B 20 F9          JR NZ,FETCH    ;IF NOT, FETCH AGAIN
4A0D C3 A0 43      JP 43A0H      ;ELSE RETURN TO T-BUG

```

You can use Program 8-4A to check the operation of any of the addresses on the basic keyboard matrix table. Simply change the memory address in the first line to the one you want to inspect.

Perhaps you are getting an inkling of what can be done here with respect to controlling the operation of machine-language programs from the keyboard. This is one kind of “wait for keyboard response” operation. The only problem is that the program is sensitive to any one of a number of keys within the specified “keyboard memory” address. Usually a wait-for-key-response function works with two extremes: do something when *any* key is depressed, or do something when one specific key is depressed. Programs 8-4B and 8-4C illustrate these two extremes of keyboard controls.

Program 8-4B is sensitive to any key depression. The point in the “keyboard memory” to search for any key depression is 387FH. That address isn’t shown on the matrix chart in Table 8-1, but it simply represents a logical oring of all the addresses that are there. So if a key depression occurs at any of the prime “keyboard memory” addresses, address 387FH will take on a hexadecimal value other than zero.

Now, there is a small problem inherent in the idea of writing a program that senses any key depression: depressing some key in

order to initiate the program can be sensed and considered a key depression for the operation of the program itself. Because of that effect Programs 8-4B and 8-4C begin with a WAIT operation—one that does not begin looking for new key depressions until a key (the key that initiates the program) is released.

Load the program in Program 8-4B and you will see it does nothing until you depress any *key* on the keyboard. The last line in the program then brings up the T-BUG monitor, making it convenient for you to do a *Register* function and inspect the content of the A register.

Going to the opposite extreme, it is often important to make the scheme sensitive to the depression of just one particular key. You must determine which key in advance and then write the program accordingly.

The listing for Program 8-4C is sensitive only to the A key. It appears to do nothing at all until you strike that particular key; then the program returns to the T-BUG monitor.

The WAIT phase of Program 8-4C “buzzes” the keyboard until it is completely free of any key depressions, then the FETCH phase begins looking for the depression of the A key. According to the matrix table, the A key is characterized by data 02H at address 3801H. Thus the FETCH portion of the program is written to inspect the content of address 3801H, compare it with 02H, and come to an end if the comparison is true.

### Three Classes of Keyboard Entries

In-line, keyboard-entry situations generally fall into one of three categories:

- Interrupt the program and wait for a certain keystroke (similar to BASIC's INPUT statement).
- Interact with the program via a keystroke but do not interrupt ongoing activity (similar to BASIC's INKEY\$ operations).
- Interact with, but do not interrupt, an ongoing program only as long as a key is depressed (similar to the PEEK(1446) operations in Chapter 3).

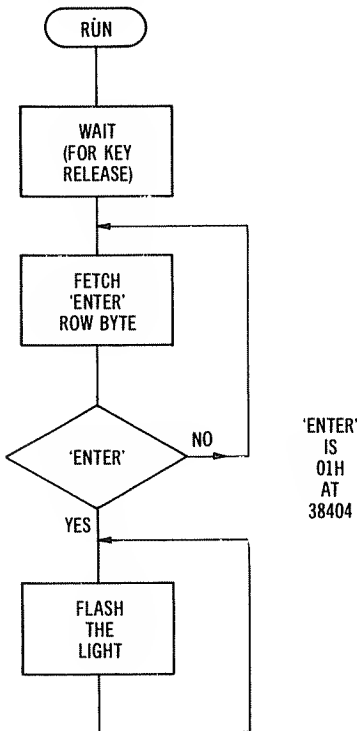
Program 8-5, flowcharted in Fig. 8-2, is an example of interrupting the execution of a program until a keystroke is made. In this case, striking the ENTER key resumes program operation, causing a rectangle of light to flash in the upper left-hand corner of the screen.

The program continuously looks for an ENTER keystroke (01H at address 3840H) until that condition is satisfied. This amounts to an interruption of the overall program. Once the operator strikes the ENTER key, the program breaks out of the FETCH loop and begins flashing the light on the screen.

**Program 8-5. A keyboard-entry operation that interrupts the program until the ENTER key is depressed.**

4A00	3A 7F 38	WAIT:	LO A,(387FH)	; FETCH BYTE FROM WHOLE KB
4A03	B7		OR A	; CHECK FOR KEY DEPRESSION
4A04	20 FA		JR NZ, WAIT	; IF SO, WAIT AGAIN
4A06	3A 40 38	FETCH:	LD A,(3840H)	; GET BYTE FROM 'ENTER' ROW
4A09	FE 01		CP 1	; IS IT 'ENTER'?
4A0B	20 F9		JR NZ, FETCH	; IF NOT, FETCH AGAIN
4A00	C0 C9 01		CALL 01C9H	; ELSE CLEAR THE SCREEN
4A10	C0 20 4A		CALL FLASH	; AND CALL FLASH
4A20	21 00 3C	FLASH:	LO HL, 3C00H	; POINT TO VIDEO
4A23	10	OECE:	DEC E	; COUNT DOWN
4A24	20 F0		JR NZ, OECE	; IF NOT DONE, COUNT AGAIN
4A26	15		DEC 0	; ELSE COUNT DOWN 0
4A27	CB 6A		BIT 5, 0	; CHECK IF 'OFF' CYCLE
4A29	28 04		JR Z, OFF	; IF NOT, DO 'OFF' CYCLE
4A2B	36 BF		LO (HL), 1910	; ELSE TURN ON THE SPOT
4A20	18 4F		JR OECE	; AND TIME AGAIN
4A2F	36 20	OFF:	LO (HL), 320	; TURN OFF THE SPOT
4A31	18 FA		JR OECE	; AND TIME AGAIN

Since FLASH THE LIGHT is an endless-loop operation in this example, you must operate the RESET key to get out of it. The important point, however, is that the program holds that FETCH



**Fig. 8-2. Flowchart for Program 8-5.**

loop—interrupting progress—until a certain keystroke occurs.

The program flowcharted in Fig. 8-3 is an entirely different sort of key-entry situation. The listing in Program 8-6 allows the light to flash continuously, whether the operator is striking a key or not. But every FLASH ONE CYCLE operation ends by polling the keyboard—polling specifically for an *L* or *R* keystroke.

If an *L* keystroke occurs during that keyboard-polling interval, the spot of light flashes at the upper left-hand corner of the screen. But if an *R* keystroke occurs, the position of the flashing light changes to the upper right-hand corner of the screen.

**Program 8-6. A keyboard-entry operation that accepts keyboard controls without interrupting the ongoing program.**

```

4A00 3A 7F 38  WAIT:      LD A,(387FH)      ;POINT TO KB
4A03 B7                OR A          ;ANY KEY DEPRESSED?
4A04 20 FA                JR NZ,WAIT    ;IF SO, WAIT AGAIN
4A06 CD C9 01          CALL 01C9H      ;ELSE CLEAR THE SCREEN
4A09 21 00 3C          LD HL,3C00H     ;AND INITIALIZE VIDEO
4A0C 3A 06 38  FETCH:    LD A,(3806H)   ;CHECK ROWS FOR 'L','R'
4A0F B7                OR A          ;EITHER KEY DEPRESSED?
4A10 20 05                JR NZ,CODE    ;IF SO, DO CODE
4A12 CD 30 4A  DOIT:     CALL FLASH     ;ELSE FLASH AS IS
4A15 18 F5                JR FETCH     ;AND FETCH AGAIN
4A17 FE 10  CODE:       CP A,10H       ;IS IT 'L'?
4A19 20 04                JR NZ,RIGHT   ;IF NOT,SET RIGHT
4A1B 2E 00                LD L,0        ;ELSE SET LEFT
4A1D 18 F3                JR DOIT       ;AND FLASH
4A1F FE 04  RIGHT:     CP A,4          ;IS IT 'R'?
4A21 20 FA                JR NZ,DOIT    ;IF NOT, FLASH AS IS
4A23 2E 3F                LD L,3FH      ;ELSE SET RIGHT
4A25 18 F6                JR DOIT       ;AND FLASH
4A30 1D  FLASH:       DEC E            ;COUNT DOWN E
4A31 20 FD                JR NZ,FLASH    ;IF NOT DONE,DO AGAIN
4A33 15                DEC D            ;ELSE COUNT DOWN D
4A34 CB 6A                BIT 5,D        ;CHECK IF 'OFF' CYCLE
4A36 28 04                JR Z,OFF      ;IF NOT,THEN DO 'OFF'
4A38 36 8F                LD (HL),191D  ;ELSE TURN ON THE SPOT
4A3A 18 F4                JR FLASH     ;AND DO TIMING AGAIN
4A3C 36 20  OFF:       LD (HL),32D     ;TURN OFF THE SPOT
4A3E C9                RET             ;CHECK KEYBOARD AGAIN

```

So keyboard operations do not interrupt the flashing light. Striking the *L* or *R* keys influences the program, however.

Finally, Program 8-7 and Fig. 8-4 illustrate a situation where depressing and holding down a certain key influences, but does not interrupt, the behavior of a program. In this case the flashing rectangle of light appears on the screen as long as the ENTER key is depressed. Otherwise, the light is turned off.

The program polls the keyboard after each light-flashing cycle or after finding that the ENTER key is not depressed. When the ENTER key is depressed, the program does a light-flashing cycle.

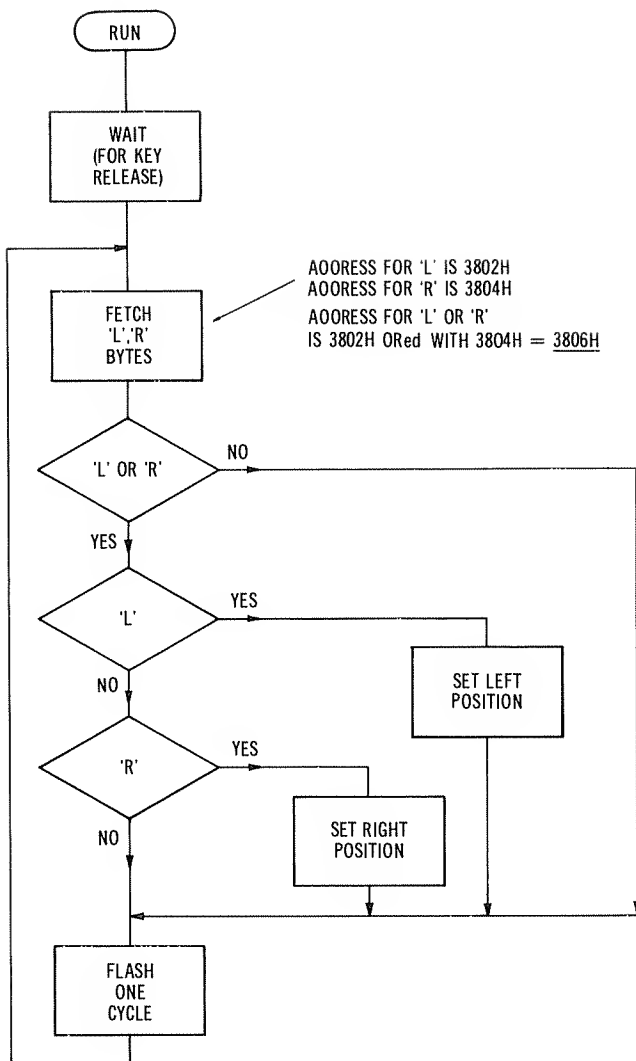


Fig. 8-3. Flowchart for Program 8-6.

### SOME SPECIAL VIDEO/KEYBOARD FUNCTIONS

Your TRS-80 is provided with a number of useful functions that are stored in the ROM space. Such functions can be CALLED from a machine-language program, thus performing important video/key-

**Program 8-7. A keyboard-entry operation that is sensitive to holding down a key.**

```

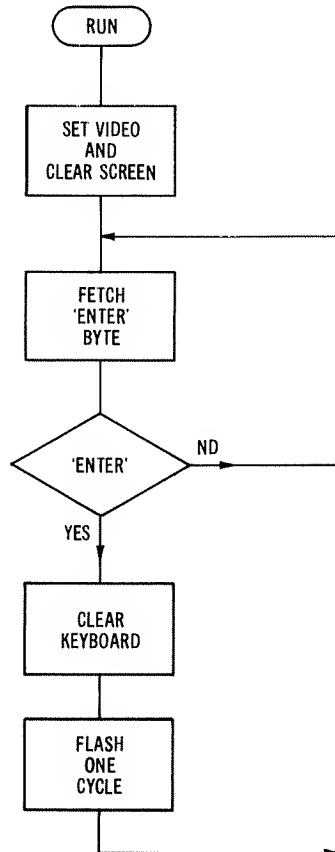
4A00 21 00 3C   TOGL:      LO HL,3C00H       ;POINT TO VIDEO
4A03 CO C9 01       CALL 01C9       ;CLEAR THE SCREEN
4A06 3A 40 38   FETCH:    LD AA,(3840H)    ;GET BYTE FROM 'ENTER' ROW
4A09 FE 01       CP 1               ;IS IT 'ENTER'?
4A0B 20 F9       JR NZ,FETCH        ;IF NOT,FETCH AGAIN
4A00 AF         XOR A               ;ZERO THE ACCUMULATOR
4A0E CO 30 4A     CALL FLASH        ;ANO CALL FLASH
4A11 C3 06 4A     JP FETCH          ;JUMP TO FETCH

```

Also add FLASH subroutine, addresses 4A30 through 4A3E, from Program 8-6

board operations without your having to do a great deal of tedious programming—it has been done for you.

Table 8-2 lists some of the most useful kinds of ROM CALLs. A simple CALL 01C9H, for example, does all the work necessary



**Fig. 8-4. Flowchart for Program 8-7.**



for clearing the screen and sending the cursor to home. That CALL does not turn on the cursor, however; as shown in Table 8-2, turning on the cursor is a two-step operation: LD A,14D, followed by CALL 033AH.

**Table 8-2. Some Useful ROM-Based Cursor-Control Routines**

Contents	Mnemonic	Comment
CD C9 01	CALL 01C9H	;CLEAR THE SCREEN AND HOME THE CURSOR
CD 49 00	CALL 0049H	;FULLY DECODED (ASCII) K8 CHARACTER ;TO THE ACCUMULATOR
CD 3A 03	CALL 033AH	;PRINT ASCII CHARACTER FROM ACCUMULATOR, ;ADVANCE THE CURSOR POSITION, AND ;SCROLL THE DISPLAY IF NECESSARY
3E 0E CD 3A 03	LD A,14D CALL 033AH	;SET CURSOR 'TURN ON' TO ACCUMULATOR ;PRINT CURSOR 'ON'
3E 0F CD 3A 03	LD A,15D CALL 033AH	;SET CURSOR 'TURN OFF' TO ACCUMULATOR ;PRINT CURSOR 'OFF'
3E 0D CD 3A 03	LD A,13D CALL 033AH	;SET 'LINE FEED/CARRIAGE RETURN' TO ACC. ;'PRINT' IT
3E 1C CD 3A 03	LD A,28D CALL 033AH	;SET 'HOME CURSOR' TO ACCUMULATOR ;'PRINT' IT
3E 1D CD 3A 03	LD A,29D CALL 033AH	SET 'CURSOR TO BEGINNING OF CURRENT LINE' ;TO ACCUMULATOR ;'PRINT' IT
3E 1E CD 3A 03	LD A,30D CALL 033AH	;SET 'ERASE TO END OF CURRENT LINE' TO ;ACCUMULATOR ;'PRINT' IT
3E 1F CD 3A 03	LD A,31D CALL 033AH	;SET 'CLEAR TO END OF FRAME' TO ACC. ;'PRINT' IT

Notice that a CALL 0049H places a fully decoded (hexadecimal ASCII) keyboard character into the accumulator. There is no need to write lengthy keyboard-decoding programs—this one does it for you.

A CALL 033AH prints an ASCII character from the accumulator to the screen. It mimics a single-character PRINT operation from BASIC. Like the PRINT command in BASIC, this CALL executes all the cursor functions. The procedure is to load the character or function into the accumulator, then do the CALL 033AH. See

more examples in Table 8-2 and some working demonstrations in Program 8-8.

Program 8-8A uses the CALL 0049H and CALL 033AH combination to get a keyboard character and print it on the screen. Those two operations, included in the KB sequence, are part of a loop that allows you to fill the screen with characters if you choose to do so.

## A Simple Text-Writing Program

### Program 8-8A. Printing and editing screen text without a cursor.

```

4A00 3A 7F 38    WAIT:      LD A,(387FH)    ;LOOK AT KB
4A03 B7          DR A          ;ANY KEY DEPRESSED?
4A04 20 FA          JR Z,WAIT    ;IF SO,WAIT AGAIN
4A06 CD C9 01    CLS:      CALL 01C9H    ;ELSE CLEAR AND HDME
4A09 CD 49 00    KB:       CALL 0049H    ;KB CHAR TO A
4A0C FE 2A          CP 42D      ;IS IT '*'?
4A0E 28 F6          JR Z,CLS    ;IF SD,CLEAR AND HDME
4A10 CD 3A 03    CALL 033AH    ;ELSE PRINT CHAR
4A13 18 F4          JR KB      ;AND GET THE NEXT DNE

```

### Program 8-8B. Printing and editing screen text with the help of a cursor.

```

4A00 3A 7F 38    WAIT:      LO A,(387FH)    ;LOOK AT KB
4A03 B7          DR A          ;ANY KEY DEPRESSED?
4A04 20 FA          JR Z,WAIT    ;IF SD,WAIT AGAIN
4A06 CD C9 01    CLS:      CALL 01C9H    ;ELSE CLEAR AND HDME
4A09 3E DE          LD A,14D    ;SET 'CURSDR ON' CDDE
4A0B CD 3A 03    CALL 033AH    ;PRINT THE CURSOR
4A0E CD 49 00    KB:       CALL 0049H    ;KB CHAR TO A
4A11 FE 2A          CP 42D      ;IS IT '*'?
4A13 28 F1          JR Z,CLS    ;IF SD, CLEAR AND HDME
4A15 CD 3A 03    CALL 033AH    ;ELSE PRINT CHARACTER
4A18 18 F4          JR KB      ;AND GET THE NEXT DNE

```

### Program 8-8C. Text writing and editing with a cursor/no-cursor control option.

```

4A00 3A 7F 38    WAIT:      LD A,(387FH)    ;LOOK AT WHDLE KB
4A03 B7          DR A          ;ANY KEY DEPRESSED?
4A04 20 FA          JR Z,WAIT    ;IF SO,WAIT AGAIN
4A06 CD C9 01    CALL 01C9H    ;ELSE CLEAR AND HDME
4A09 3E 3F          SET:      LD A,63D      ;'? ' TD A
4A0B CD 3A 03    CALL 033AH    ;PRINT IT
4A0E CD 49 00    KBI:      CALL 0049H    ;FETCH KB CHARACTER
4A11 FE 43          CP 67D      ;IS IT 'C'?
4A13 28 0B          JR Z,CURON    ;IF SD,TURN DN CURSOR
4A15 FE 4E          CP 78D      ;IS IT 'N'?
4A17 20 F5          JR NZ,KBI    ;IF NOT,GET CHAR AGAIN
4A19 3E 0F          LD A,0FH    ;ELSE TURN OFF CURSDR
4A1B CD 3A 03    CALL 033AH    ;PRINT
4A1E 18 05          JR CLS2     ;AND START TYPING
4A20 3E 0E          CURON:    LD A,0EH    ;SET CURSDR ON CODE
4A22 CD 3A 03    CALL 033AH    ;PRINT IT
4A25 CD C9 01    CLS2:      CALL 01C9H    ;CLS AND HDME THE CURSOR
4A28 3A 40 38    KB2:      LD A,(3840H)    ;GET 'CLEAR' RDW FRDM KB
4A2B FE 02          CP 2        ;IS IT 'CLEAR'?
4A2D 28 DC          JR Z,SET    ;IF SO,DO SET AGAIN
4A2F CD 49 00    CALL 0049H    ;FETCH KB CHAR TO A
4A32 CD 3A 03    CALL 003AH    ;PRINT IT
4A35 18 F1          JR KB2     ;AND GET NEXT CHARACTER

```

But the KB sequence is also sensitive to a striking of the asterisk key. Whenever that happens, operations are sent to CLS—a CALL 01C9H that clears the screen and homes the cursor. So you can type a lot of characters, then strike the \* key to clear the screen, and start all over again. It is possible to edit the text by taking advantage of the SHIFT *arrow key* functions.

The WAIT sequence at the beginning of all three programs is necessary only for these demonstrations. It makes certain nothing is printed on the screen until you release the key that starts the programs in the first place.

When running Program 8-8A you will find things a bit awkward because the cursor symbol isn't there to let you know where the next character will be printed—a feature that is especially troublesome when one is attempting to move the cursor to a position in a blank portion of the screen. So Program 8-8B fixes that.

Program 8-8B runs the same way as the first one in the series, but it includes the LD A,14D and CALL 033AH sequence to turn on the cursor symbol. Those two instructions, included in the CLS operations, perform the same task as BASIC's PRINT CHR\$(14) as described in Chapter 2.

The notions are further refined in Program 8-8C. The WAIT sequence ends with a CALL 01C9H to clear the screen and home the cursor. But then the SET sequence loads the hexadecimal code for a question mark into the accumulator and uses a CALL 033AH to print it. You thus see a clear screen with a question mark appearing in the upper left-hand corner.

The KB1 sequence gives you an opportunity to see or not see the cursor symbol. Using a CALL 0049H to pick up a keyboard character, this sequence is sensitive to both the C and N keys (CP 67D and CP 78D, respectively). If you choose to see the cursor and respond to the question mark by striking the C key, operations jump down to CURON—a set of two instructions that turn on the cursor symbol. But if you answer the question mark with a N, the last three lines in the KB1 sequence turn off the cursor and jump operations to CLS2.

CLS2 clears the screen and homes the cursor, and then the sequence in KB2 allows you to print characters at your heart's content. That particular video/keyboard sequence, incidentally, improves on the first two programs by doing away with the idea of clearing the screen by striking the asterisk key. Instead, the LD A,(3840) and CP 02H sequence makes the system sensitive to the CLEAR key.

So when you strike the CLEAR key while the system is running the screen-printing KB2 sequence, the system returns to SET. You then have a chance to select seeing or not-seeing the cursor, and

on making your selection the screen clears and you are back in the typing business again.

The main point of Program 8-8C is to illustrate the operation of the cursor-on/cursor-off operations—examples of cursor controls CALLED from a machine-language program. In the process of illustrating the point, however, you will find a review of much of the material already described in this chapter.

## **T-BUG AND THE MEMORY ENVIRONMENT**

The memory environment for T-BUG isn't nearly as complicated as that of BASIC. T-BUG, itself, is loaded as a SYSTEM-oriented, machine-language program between 4380H and 497FH (see Fig. 7-1). Obviously, any program written from T-BUG ought not alter the content of any address in that area—otherwise T-BUG will be destroyed.

All addresses from 4980H through the top of your RAM space are thus free for programming applications. Note, for instance, that most of the previous demonstrations begin at address 4A00H. This leaves plenty of space between the top of T-BUG and the beginning of those programs.

There is, however, some additional RAM space below the T-BUG monitor area: everything from 41E6H (the beginning of BASIC's i/o buffer) to 437FH (the address just below T-BUG). T-BUG users generally avoid this RAM space for a couple of reasons. First, it is a relatively small amount of space compared to that available above T-BUG. Second, any SYSTEM program, unless instructed otherwise, will set the program stack at 4288H—in the BASIC i/o space.

Nevertheless, the RAM space below T-BUG can prove quite useful, as shown later in this chapter. The fact that it is a fairly small amount of memory space isn't relevant to the applications that will be cited at that time, and the stack can always be pointed elsewhere, say to the top of your available memory space.

Now, let us take a look at some memory operations that take advantage of the wealth of RAM above T-BUG.

### **Working With User's RAM Above T-BUG**

Program 8-9 allows you to compose text on the video screen and then transfer it to some user's RAM space where it can be saved for an indefinite period. In the meantime, you can tinker around in video memory, recalling the original text whenever you choose.

One of the most impressive features of this program is the speed at which saved text can be restored to the screen. BASIC users are accustomed to a fairly slow printing rate, even when using direct

**Program 8-9. Assembly listing for a program that allows you to transfer video memory/screen information to a block of higher memory space.**

```

4A00 CD C9 01    NEXT:      CALL 01C9H      ;CLS AND HOME THE CURSOR
4A03 3E 3F      QUERY:     LD A,63D        ;'? ' TO ACCUMULATOR
4A05 CO 3A 03    CALL 033AH      ;PRINT IT
4A08 CD 49 00    CALL 0049H      ;KB CHAR TO ACCUMULATDR
4A0B FE 52      CP 82D          ;IS IT 'R'?
4A0D CA 60 4A    JP Z,DISP      ;IF SO,RECALL DISPLAY
4A10 FE 43      CP 67D          ;IS IT 'C'?
4A12 CA 20 4A    JP Z,COMP      ;IF SD, COMPOSE
4A15 C3 03 4A    JP QUERY      ;ELSE QUERY AGAIN
                                   ;
4A20 CO C9 01    COMP:      CALL 01C9H      ;CLS AND HOME THE CURSDR
4A23 3E 0E      LD A,140        ;SET 'CURSOR ON' CODE
4A25 CD 3A 03    CALL 033AH      ;PRINT IT
4A28 CD 49 00    KB:        CALL 0049H      ;GET KB CHARACTER
4A2B FE 2A      CP 42D          ;IS IT '*'?
4A2D CA 40 4A    JP Z,SAVE      ;IF SO,SAVE OISP IN MEM
4A30 CD 3A 03    CALL 033AH      ;ELSE PRINT THE CHAR.
4A33 18 F3      JR KB          ;AND GET THE NEXT ONE
                                   ;
4A40 3E 0F      SAVE:      LD A,15D        ;SET 'CURSOR OFF' COOE
4A42 CD 3A 03    CALL 033AH      ;PRINT II
4A45 21 00 3C    LO HL,3C00H    ;PDINT TO VIDEO
4A48 11 00 4B    LD DE,4B00H    ;POINT TO MEMORY
4A4B 01 00 04    LD BC,1024D    ;SET NUMBER OF BYTES
4A4E ED B0      LDIR           ;XFER SCREEN TO MEMORY
4A50 C3 00 4A    JP NEXT      ;AND BACK TO BECINNINC
                                   ;
4A60 CD C9 01    DISP:      CALL 01C9H      ;CLS AND HDME THE CURSDR
4A63 21 00 4B    LD HL,4B00H    ;PDINT TO MEMORY
4A66 11 00 3C    LO DE,3C00H    ;POINT TO VIDEO
4A69 01 00 04    LD BC,1024D    ;SET NUMBER DF BYTES
4A6C ED B0      LDIR           ;XFER MEMDRY TO SCREEN
4A6E CO 49 00    HOLD:      CALL 0049H      ;KB CHAR TO ACCUMULATDR
4A71 FE 2A      LD A,42D        ;IS IT '*'?
4A73 20 F9      JR NZ,HOLD      ;IF NOT, HDLD AGAIN
4A75 C3 00 4A    JP NEXT      ;ELSE BACK TD BEGINNING

```

POKE and PEEK operations. But this machine-language program flashes a full screen of text into place in a very short time—certainly too short a time to see the characters being printed individually.

Load the program and run it from 4A00H. The screen will clear and a question mark will appear in the upper left-hand corner of the screen. Reply by striking the *C* (Compose) key. That will replace the question mark with the typing cursor.

Now type in any sort of text you want, using any editing keys as you wish. When done; strike the \* key. As soon as you do that the screen will clear and the question mark will appear again.

This time, respond by striking the *R* (Recall) key. Presto! there is your original text returned to the screen.

The general idea is to compose new text (and save it by striking the \* key) or recall old text from memory.

Note the use of special video/keyboard CALLs in the listing, as well as the Z-80's block-transfer instruction, LDIR. In view of

past demonstrations you should be able to figure out the operating details yourself. Hint: All 1024 video bytes, beginning at 3C00H are saved in RAM, beginning at 4B00H.

Program 8-10 is a somewhat more sophisticated version of the one just described. This program allows you to generate the effect of doubling the video screen capacity; it allows you to work with two different "pages" of text.

You can, for instance, compose or edit some text appearing on the screen, and then exchange it with a full screen of text previously saved in memory. That text can also be deleted or edited and exchanged again with the original. The effect is that of working with two separate crt screens (although you can really work with just one at a time).

**Program 8-10. Listing for a memory exchange program gives the user two "pages" of text to work with.**

```

4A00 CO 49 00    NEXT:    CALL 0049H    ;FETCH KB CHARACTER
4A03 FE 58      CP 88D      ;IS IT 'X'?
4A05 28 14      JR Z,CEX    ;IF SO, OO EXCHANGE
4A07 FE 4E      CP 78D      ;IS IT 'N'?
4A09 28 0B      JR Z,CLR    ;IF SO,CLEAR THE SCREEN
4A0B FE 45      CP 690      ;IS IT 'E'?
4A00 28 02      JR Z,CED    ;IF SO, OO EDITING
4A0F 18 DF      JR NEXT     ;ELSE LOOK AT KB AGAIN
4A11 CD 30 4A    CED:      CALL COMP    ;CALL COMPOSE MODE
4A14 18 F9      JR NEXT     ;THEN START AGAIN
4A16 CD C9 01    CLR:      CALL 01C9H    ;CLS AND HDME THE CURSOR
4A19 18 F9      JR NEXT     ;THEN START AGAIN
4A1B CD 50 4A    CEX:      CALL EXCH    ;CALL EXCHANGE OPS
4A1E 18 F9      JR NEXT     ;THEN START AGAIN
;
4A30 3E 0E      COMP:      LO A,14D      ;SET 'CURSOR ON' COOE
4A32 CD 3A 03    CALL 033AH    ;PRINT IT
4A35 CO 49 00    KB:       CALL 0049H    ;FETCH KB CHARACTER
4A38 FE 2A      CP 42D      ;IS IT '*'?
4A3A 28 05      JR Z,CROFF   ;IF SD, WRAP IT UP
4A3C CD 33 00    CALL 033AH    ;ELSE PRINT THE CHAR.
4A3F 18 F4      JR KB       ;AND CET THE NEXT DNE
4A41 3E 0F      LD A,15D     ;SET 'CURSDR OFF' CODE
4A43 CD 3A 03    CALL 033AH    ;PRINT IT
4A36 C9         RET         ;RETURN TO NEXT
;
4A50 21 00 4B    EXCH:      LD HL,4B00H    ;POINT TO MEMDRY
4A53 11 00 3C    LD DE,3C00H    ;POINT TD VIDEO
4A56 01 00 04    LD BC,1024D    ;SET NUMBER OF BYTES
4A59 1A         LD A,(DE)      ;VIDEO BYTE TO ACC.
4A5A EO AO      LDI          ;MEMORY BYTE TO VIDEO
4A5C 2B         DEC HL        ;POINT BACK MEMORY ADDR.
4A5D 77         LO (HL),A     ;ACC. TO MEMORY
4A5E 23         INC HL        ;POINT AHEAD MEMDRY ADDR.
4A5F AF         XDR A         ;CLEAR ACCUMULATDR
4A60 B9         CP C         ;LSB DF COUNTER DDNE?
4A61 20 F6      JR NZ,SWAP    ;IF NOT,SWAP ANDTHER BYTE
4A63 B8         CP B         ;MSB OF COUNT DONE?
4A64 C8         RET Z         ;IF SD, RETURN TO NEXT
4A65 18 F2      JR SWAP      ;ELSE SWAP ANOTHER BYTE

```

The program uses three control characters from its monitor mode: *N* clears the screen and the “page” residing there, *E* sets up the Edit mode that allows you to type new information or alter some old text on the displayed page, and *X* exchanges the two pages—transfers the one currently on the screen to memory, and displays the page previously saved in memory.

Striking the \* key gets the system out of the Edit or exchange mode and returns it to the monitor. The only way out of the program is to work the RESET push button.

So load Program 8-10 and give it a try. Play with the program for a bit and then you will be in a better position to appreciate the internal workings.

With the program loaded, run it by doing a J4A00. Then delete the text on the screen by striking the *N* key. After that, get into the Edit mode by striking the *E* key. At that time the cursor symbol should appear in the upper left-hand corner of the screen.

Now type in some text—anything will do. Notice you can move around the cursor, without disturbing the text, by working the arrow keys when the SHIFT key is depressed.

When you are satisfied with the text on this first page, strike the \* key to return to the monitor. To save that page and call up its alternate, strike the *X* key.

There’s no telling what the alternate page will look like at this time, but you can always clear it up by striking the *N* key and then the *E* key to type some new stuff. And when you are done with that page, strike the \* key.

To bring back the original page, strike the *X* key. Everything should appear as it did earlier in the demonstration. If you want to see your second page again, simply strike the *X* key again.

The main point of the demonstration is that exchange operation—swapping the contents of the video memory with some user’s memory of equal size elsewhere in the system. That “hidden-page” memory, incidentally, begins at 4B00H, and occupies 1024 (decimal) bytes from there. In a manner of speaking, this program creates an alternate video memory.

Referring to the listing for Program 8-10, labels NEXT, CED, CLR, and CEX make up the “monitor.” This group of instructions uses a CALL 0049H to pick up a keystroke and put its ASCII character code into the accumulator. After that, a series of compare instructions look for the control characters and call the appropriate routines.

The COMP routine, for instance, is called whenever you want to Edit, or compose, some new text on the screen. It begins by turning on the cursor (so you will know where the screen printing will begin) and CALLing 0049H to pick up an ASCII character code

from the keyboard. The routine uses a CALL 033AH to print the characters and advance the cursor.

If, while running the COMP routine, the system sees an asterisk character, it turns off the cursor (so that it will not be saved in memory with the other text) and returns operations to the monitor.

Striking the N key from the NEXT routine does a CALL 01C9H to clear the screen and home the cursor.

The EXCH routine is the one responsible for exchanging the contents of the video memory with the alternate video memory space. It simply runs through both memory blocks, passing a byte at a time through the accumulator. When the exchange is completed, operations automatically return to the monitor.

You can change the address of the alternate video memory by changing the hexadecimal address loaded into the HL register pair at the beginning of the EXCH routine.

### **Using the I/O Buffer Space**

The i/o buffer is located between 41E6 and 42E7H on the TRS-80 memory map. That isn't very much RAM space, compared to that above the T-BUG's space, but it can be useful. The i/o buffer space is especially useful when one is attempting to control machine-language programs, using keyboard entries having more than one character in them. Thus far all keyboard controls have used just one character.

A CALL 0361H connects the keyboard to the i/o buffer. That CALL does a lot of routine housekeeping for you. For instance, it automatically begins loading keyboard characters from the first character location in the buffer; there is no need to initialize it yourself. That CALL also turns on the cursor symbol for you, and, after each keyboard entry, it advances the cursor position and i/o pointer.

What's more, the system remains in that i/o loop until you strike the ENTER key. There is no need for dreaming up a special control character for breaking out of the routine.

What's the price to pay for all of these nice features? It turns out to be rather small. Remember that any SYSTEM tape will set its program stack into the i/o memory space, unless instructed to do otherwise. So if you fail to relocate the program stack before doing a CALL 0361H, you run the risk of losing the program.

Before doing a CALL 0361H, then, you should relocate the stack with a LD SP,NN type of instruction. LD SP,7FFF for example. (That particular address represents the top of RAM space for a 16K machine.) You can put the stack anywhere in the user's RAM space, remembering that the stack will always build downward.

To get a feeling for how this keyboard i/o scheme works, try this:



4A00 31 FF 4F	LD SP,4FFFH	;GET STACK OUT OF I/O
4A03 CD C9 01	CALL 01C9H	;CLEAR THE SCREEN
4A06 CD 61 03	CALL 0361H	;KB TO I/O

and insert a Break at 4A09.

Running from 4A00 then lets you type in all sorts of characters—much like typing a new program line in BASIC or answering an INPUT instruction. When you are through entering some characters, strike the ENTER key. If you have inserted the Break as suggested, the system will return safely to the T-BUG monitor.

After leaving a CALL 0361H routine, your text will be saved in the i/o buffer space. Adding the following instructions to the previous ones will show you the contents of the entire i/o buffer. That, of course, should confirm that your information was indeed entered there from the keyboard.

4A09 21 E7 41	LD HL,41E7H	;POINT TO START OF I/O
4A0C 11 00 3C	LD DE,3C00H	;POINT TO START OF VIDEO
4A0F 01 FF 00	LD BC,00FFH	;SET BYTE COUNT
4A12 ED 80	LDIR	;SHOW I/O BUFFER

and insert a Break at 4A14.

Now you can run from 4A00, load some characters into the i/o buffer and strike the ENTER key to end the loading operation and display the contents of the i/o buffer.

There will most likely be a lot of garbage in the display, but your keyboard input will always appear at the beginning. Also note that your characters are bracketed by @ symbols. Those represent the nulls that are automatically inserted into a buffer string for BASIC interpretation purposes. (See “The I/O Buffer” in Chapter 4.)

Program 8-11 represents an application of this i/o buffer routine. The point of the demonstration is to show how the buffer can be used as an entry point for control expressions having more than one character in them. The control expressions in this case are *BL*, *EX*, and *EN*. Doing a *BL* makes a small rectangle of light blink on and off. Entering an *EX* makes that rectangle exchange positions—from the left to right or from the right to left corners of the screen. Entering expression *EN* ends the program and returns operations to the T-BUG monitor.

The program begins by setting the stack of the i/o buffer and to address 4FFFH. Then it clears the screen and prints a question-mark prompt symbol. The idea is to let you know it is time to designate one of the three basic functions: *BLink*, *EXchange*, or *END*.

Instruction CALL 0361H calls the i/o routine, and it remains in that loop until you strike the ENTER key. The system's sensitivity to the ENTER key is built into the routine—you don't have to write it into the program itself.

**Program 8-11. Using the i/o buffer as an entry point for multiple-character keyboard commands.**

```

4A00 31 FF 4F      STAK:      LD SP,4FFFH      ;SET STACK OUT OF I/O
4A03 CD C9 01      START:     CALL 01C9H      ;CLEAR THE SCREEN
4A06 3E 3F          LD A,3FH      ;SET UP '?'
4A08 CD 3A 03      CALL 033AH      ;PRINT IT
4A0B CD 61 03      CALL 0361H      ;KB INPUT TO I/O BUFFER
4A0E 21 00 3C      LD HL,3C00H      ;SET LEFT VIDEO POINTER
4A11 11 3F 3C      LD OE,3C3FH      ;SET RIGHT VIDEO POINTER
4A14 01 E7 41      LD BC,41E7H      ;SET I/O POINTER
4A17 0A          GETC:         LD A,(BC)      ;GET CHAR FROM I/O
4A18 FE 20          CP 20H          ;IS IT ALPHABETICAL?
4A1A 30 08          JR NC,TRY      ;IF SO,GIVE IT A TRY
4A1C FE EF          CP 0EFH          ;END OF USEFUL I/O SPACE?
4A1E CA 03 4A      JP Z,START      ;IF SO, START AGAIN
4A21 03          INC BC          ;ELSE LOOK AT NEXT I/O
4A22 18 F3          JR GETC        ;AND GET IT
4A24 FE 42          TRY:         CP 42H          ;IS 1ST CHAR 'B'?
4A26 20 00          JR NZ,NEXT      ;IF NOT,CHECK FOR 'E'
4A28 03          INC BC          ;ELSE LOOK AT 2ND CHAR
4A29 0A          LD A,(BC)        ;2ND CHAR TO A
4A2A FE 4C          CP 4CH          ;IS IT 'L'?
4A2C C2 03 4A      JP NZ,START      ;IF NOT,START ALL OVER
4A2F CD 00 4C      CALL BLINK      ;ELSE BLINK THE SPOT
4A32 C3 03 4A      JP START        ;AND START ALL OVER
4A35 FE 45          NEXT:        CP 45H          ;IS 1ST CHAR 'E'?
4A37 C2 00 4A      JP NZ,START      ;IF NOT,START ALL OVER
4A3A 03          INC BC          ;ELSE POINT TO NEXT I/O
4A3B 0A          LD A,(BC)        ;AND PUT IT INTO A
4A3C FE 58          CP 58H          ;IS 2ND CHAR 'X'?
4A3E 20 06          JR NZ,NEXT2     ;IF NOT, TRY FOR 'N'
4A40 C0 20 4C      CALL EXCH        ;ELSE CALL EXCHANGE
4A43 C3 03 4A      JP START        ;AND START ALL OVER
4A46 FE 4E          NEXT2:       CP 4EH          ;IS 2ND CHAR 'N'?
4A48 C2 03 4A      JP NZ,START      ;IF NOT,START ALL OVER
4A4B CD 80 43      CALL 4380H      ;CALL T-BUG TO GET AWAY
;
4C00 CD C9 01      BLINK:        CALL 01C9H      ;CLEAR THE SCREEN
4C03 36 20          RBLINK:      LD(HL),20H      ;CLEAR SPOT
4C05 CD 15 4C      CALL DELAY      ;DO TIME DELAY
4C08 36 BF          LD (HL),191D      ;PRINT SPOT
4C0A CD 15 4C      CALL DELAY      ;DO TIME DELAY AGAIN
4C0D 3A 40 38      LD A,(3840H)      ;GET 'BREAK' KB ROW
4C10 FE 04          CP 4H          ;IS IT 'BREAK' KEY?
4C12 20 EF          JR NZ,RBLINK    ;IF NOT,BLINK AGAIN
4C14 C9          RET              ;ELSE RETURN
4C15 06 1F          DELAY:        LD B,1FH        ;SET MSB OF DELAY
4C17 0E FF          SETC:         LD C,OFFH        ;SET LSB OF DELAY
4C19 0D          DECC:          DEC C          ;COUNT DOWN LSB
4C1A 20 FD          JR NZ,DECC      ;IF NOT DONE,DO MORE
4C1C 05          DEC B          ;ESLE COUNT DOWN MSB
4C1D 20 F8          JR NZ,SETC      ;IF NOT DONE,DO LSB AGAIN
4C1F C9          RET              ;ELSE RETURN TO BLINK
4C20 EB          EXCH:          EX DE,HL        ;SWAP VIDEO LOCATIONS
4C21 C3 00 4C      JP BLINK        ;AND MAKE IT BLINK

```

Thus, the instruction following CALL 0361H begins a decoding operation. That decoding operation runs through the i/o buffer, looking for character patterns BL, EX, and EN. On finding one of those sequences the program takes appropriate action, using one of the subroutines beginning at address 4C00H.

If the program does not find a BL, EX, or EN combination in the i/o buffer, it defaults to START, giving you a chance to enter the characters again. Incidentally, the program works equally well if you enter the full words, BLINK, EXCHANGE, or END. The i/o decoding scheme just looks for the first two characters in each case—if other characters follow, that's of no consequence.

#### *SUMMARY*

CALL 0361H connects the keyboard to the i/o buffer.

It also:

- Turns on the cursor symbol.
- Enters keyboard characters from the beginning of the buffer space.
- Automatically advances the cursor position.
- Remains in effect until the user strikes the ENTER key.

A CALL 0361H must be preceded by an LD SP,NN operation to get the SYSTEM program stack out of the i/o space.

## CHAPTER 9

# Introducing the TRS-80 Editor/Assembler

At first thought, Radio Shack's Editor/Assembler software (Catalog No. 26-2002) might seem to be little more than a convenient language translator—a cassette-based system that translates standard Z-80 mnemonics into machine code. Certainly Editor/Assembler (or EDTASM) does that, and that is a nice feature. But there is a lot more to it.

In fact, if you have never used an assembler system before, you can be easily confused by all the special features that carry it far beyond being a simple language translator. This assembler is a system that is best understood by working into it gradually, starting with some familiar ideas and letting first-hand experience lead you to a real appreciation of the power of EDTASM.

This is the approach used in this chapter and the one that follows.

### A FEW PRELIMINARY NOTES

The main purpose of EDTASM is to let you type in assembly or source-code instructions from the keyboard. You still have to design the overall program yourself, but EDTASM does all the legwork involved in translating the mnemonics into machine-language, or object-code, instructions. EDTASM also saves a lot of time by letting you specify variables in terms of decimal numbers and specify jumps and calls in terms of labels, and by generally letting you avoid a lot of routine and often tedious figuring and table searching.

In short, EDTASM takes advantage of the computing power of your machine, doing a lot of clerical work for you. And you will find that using EDTASM makes hand assembling seem primitive.

So EDTASM lets you type in machine-language programs in an assembly-language format. When you are done with that phase of the task, you enter a command that tells the system to assemble the program into machine code. It is possible to view your assembly or source code listing right along with the assembled machine-code listing.

If you make any errors in the assembly listing, EDTASM will let you know about it. Of course, EDTASM cannot be held responsible for mistakes in the overall program, but it does seek out syntax errors—improperly specified or misspelled source instructions.

Programs generated with EDTASM can be saved on cassette tape and written out to a line printer.

The only drawback is that you cannot actually test your program while the EDTASM program is residing in the user's memory space. In order to check the operation of an EDTASM-generated, machine-language program, you must first save the program on cassette tape. Then the program has to be entered as an ordinary machine-language program under the SYSTEM command. EDTASM is thus lost in the process, but you can check the operation of the program.

If you don't like the way the program runs, you have the option of debugging it with T-BUG or by reloading EDTASM and fixing the problem in assembly language.

There is a very practical reason for this shortcoming of EDTASM—not being able to test a machine-language program without destroying EDTASM. The reason is that most TRS-80 users do not have enough RAM to support both EDTASM and a machine-language program of any reasonable size.

The EDTASM program occupies some 6300 bytes of RAM. The entry point is 18058 (decimal), but there are special registers located at lower memory addresses and, indeed, in the i/o buffer space as well.

A program of this sort that occupies only about 6K of RAM would seem to leave a lot of RAM for machine-language programs you develop with EDTASM—about 7K of memory for a 16K system. But there's a snag at this point.

As you type in the assembly codes they are stored in RAM, beginning at the point where the assembler program, itself, leaves off. So as you type in the assembly program it grows upward in the user's RAM space, taking up one byte of memory for every character you enter. And you will find that even a modest assembly listing will use up at least a couple of thousand bytes of memory. But that isn't all.

When you tell EDTASM to assemble the program the object codes start filling up more RAM space, beginning where the source codes left off. So there goes another big chunk of RAM space.

This all adds up to the fact that EDTASM eats up a whole lot of RAM, leaving little space for entering the machine-language program you are trying to develop. What's more, EDTASM uses up the prime memory space that you normally want for the program you are writing.

No, we cannot test a program developed by EDTASM without losing or writing over EDTASM. And that is why we have to live with the awkward process of first saving the program on tape and then entering it again under the SYSTEM command.

There is another trouble with EDTASM. It is a minor, but sometimes annoying, problem. Recall that it is possible to bomb out of T-BUG and end up in the BASIC monitor. Whenever that happens you can return to T-BUG by calling SYSTEM and entering /17312. But that cannot be done as effectively with EDTASM.

Once you return to BASIC from EDTASM, some valuable registers are disrupted. It is possible to call SYSTEM and return to EDTASM by entering /18058; and it will seem that EDTASM is working again. But it won't take long to find that things don't work out very well—you will sooner or later notice some strange effects, such as seeing the video display suddenly change to a 32-character line format.

So if, for any reason, you find the system going from EDTASM to BASIC, prepare yourself to lose everything you've done so far and load the EDTASM system again from scratch.

It is not the purpose of this discussion to cast a bad light on EDTASM. Rather, the idea is to point out a couple of problems that must be clearly understood from the start. Difficulties aside, EDTASM is a powerful programming tool.

## **BRINGING UP EDTASM**

Obviously, the first step in using EDTASM is to acquire a copy of the cassette tape and load it into your system. Since it is a machine-language program, it is loaded from the SYSTEM command.

So enter SYSTEM, set up the cassette player for playing the EDTASM tape, and then enter the file name EDTASM. The program is long, so it takes a few minutes to load. The familiar flashing asterisk in the upper left-hand corner of the screen confirms a good loading operation.

When the EDTASM tape is fully loaded into your system, you will see the machine-language prompt symbol (an asterisk), followed by a question mark and cursor symbol. Respond by entering a slash (/).

The screen will then clear and you will see a title message and an asterisk, indicating that EDTASM is up and ready to go.

You can respond at this point by typing anything you choose from the keyboard. But if the assembler is to be of any use, the things you type and enter ought to make sense in the context of a few rules.

### **LINE NUMBERS, ORG, AND END**

One important rule is that every assembly-language program must begin by specifying the origin of that program. The origin actually spells out the address at which the machine-language program you develop should begin loading—when, indeed, you load it from cassette tape later on in the process.

A good many T-BUG programs cited in earlier chapters began at address 4A00H. That is the idea here. You can set the origin of your machine-language program anywhere you choose within the system's available memory space.

Before learning how to specify the origin of the program, it is important to understand that assembly listings you enter from the keyboard must have line numbers assigned to them. These line numbers serve as convenient identifiers for each line of assembly text you enter from the keyboard. The numbers are not saved on tape with the finished machine-language program.

When you are ready to begin entering an assembly program under EDTASM, the first task is to specify a line number and an interval between line numbers. It works something like BASIC's AUTO command.

Suppose you have just brought up EDTASM as described in the previous section. To specify the starting line number and interval, respond to the asterisk by entering something such as I100,10. This is telling the assembler to begin numbering the text lines at 100 and to use increments of 10. After that, the assembler takes care of numbering the lines for you.

As another example, entering I10,5 tells the numbering scheme to begin at line 10 and increment in steps of 5.

Here is how the first couple of steps in an assembly programming operation could appear on the screen:

```
*I100,10
00100      ORG      4A00H
00110__
```

In this case the user responded to the asterisk by entering I100,10. That set up the line-numbering scheme to begin at line 100 and increment in steps of 10.

The assembler printed the 00100 following that particular entry. The 00100 is a five-digit version of the starting line number, 100.

You need not specify line numbers with the leading zeros, but the assembler will always insert them as it goes along.

As mentioned earlier, an assembly program should begin by specifying the origin of the machine-language program you will ultimately develop. That is the purpose of the text in line 00100. Here, the user responded by striking the right-arrow key ( $\rightarrow$ ), typing ORG, striking the right-arrow key again, typing 4A00H, and, finally, striking the ENTER key. This sequence of operations completes the necessary process of specifying the program's origin. In this case the origin is set at address 4A00H, and when the finished program is loaded into the system at some later time the machine codes will begin loading at address 4A00H.

Once the user completes that ORG line and strikes the ENTER key, the assembler accepts the information and responds by printing out the next line number.

#### NOTES

The ultimate origin of a machine-language program generated under EDTASM must be specified by ORG, followed by the origin address.

The ORG code must be written at the beginning of the second vertical field on the screen—an effect accomplished easily by striking the right-arrow key just one time.

If the ORG address is specified as a hexadecimal number, it must be followed by an H. Otherwise, the assembler will attempt to interpret the number as a decimal version of the address.

Every assembly program must conclude with an END operation. Specifically, it must conclude with END, followed by an address. The END line specifies the entry point of the machine-language program that you have prepared.

Often the ORG address and END are the same, thus indicating that the execution of the machine-language program begins at the lowest address it occupies in RAM. But it is often desirable to enter a program at an address that is somewhat higher than the origin. This being the case, the END address will be different from the one specified in the ORG line.

So here is the simplest possible assembly program as it appears on the screen:

```
*100,10
00100      ORG      4A00H
00110      END      4A00H
00120...
```



The program specifies an origin and an entry point. Note that the END operation, like ORG, must be written at the beginning of the second vertical field.

This program won't do anything, because there are no real machine instructions appearing between ORG and END. The example merely illustrates how an assembly program is started and concluded.

#### NOTES

The END, followed by an address, at the conclusion of an assembly listing accomplishes two things:

1. The END marks the end of the assembly text.
2. The END address marks the entry point of the machine-language program being generated under EDTASM.

EDTASM will accept initial line numbers and generate line numbers automatically thereafter in the range of 00000 through 65529.

### DELETING AND WRITING PROGRAM TEXT

Every assembly program must begin with an ORG and address and conclude with an END and address. The actual program text is fit in between.

If you have actually entered the simple, do-nothing assembly text illustrated in the previous section, it is necessary to delete it before beginning the next example. (Program text could be *inserted* between those two ORG and END operations, but that sort of editing function is reserved for a later discussion.)

To delete an entire assembly listing, first strike the BREAK key to return to EDTASM's monitor mode. This will break up the automatic line-numbering operation and return an asterisk. The asterisk is the prompt symbol for the monitor mode.

Once you are back into the monitor mode you can delete the entire text by answering the asterisk with D#:\* and entering it. This particular set of characters literally means delete (the D) from the beginning of the text (#) to (:) the end of the text (\*). This is much the same thing as doing the NEW command in BASIC.

After entering D#:\* the system returns an asterisk prompt symbol, signaling the fact that the system is ready for the next command. So do the following to get another program-writing operation underway:

```
*1100,10
00100      ORG      4A00H
00110_
```

With the origin of the program thus specified, you can begin writing some program text. To keep things simple, suppose you are working with a one-instruction program that loads 3C00H to the HL register pair.

So strike the right-arrow key to set the beginning of the instruction at the start of the second vertical field, then type in the appropriate instruction, and strike the ENTER key to signal the end of the line of text. By doing that, the display should look something like this:

```
*1100,10
00100      ORG  4A00H
00110      LD  HL,3C00H
00120_
```

Next, signal the end of the listing by entering END 4A00H:

```
*1100,10
00100      ORG  4A00H
00110      LD  HL,3C00H
00120      END 4A00H
00130_
```

This example represents a complete assembly program. It certainly isn't a very sophisticated program, but it illustrates the general procedures described thus far.

Notice that the Z-80 instruction is typed as it appears on the instruction-set listings. The H appearing at the end of the 3C00 address indicates that the number is in hexadecimal. If you omit the H, the assembler will attempt to interpret it as a decimal number.

Obviously you can enter more than one instruction line. The upper limit on the number of instructions in the assembly listing depends on the memory capacity of your TRS-80 system: the larger the memory capacity, the greater the number of instructions the assembler can handle.

## **ASSEMBLING THE SOURCE PROGRAM**

After entering the assembly program text—the source program—it is time to let EDTASM assemble it for you. The assembly operation automatically translates the instructions into machine code and assigns those codes to the proper address locations.

Assuming you have entered the simple one-line program cited in the previous section, first return to EDTASM's monitor mode by striking the BREAK key. The appearance of the asterisk prompt character confirms that the system has returned to the monitor.

To assemble the program, answer the asterisk by entering an A from the keyboard. After that, the screen should show this sort of display:

```

*1100,10
00100          ORG  4A00H
00110          LD  HL,3C00H
00120          END  4A00H
00130
*A
4A00           00100      ORG  4A00H
4A00 21003C    00110      LD  HL,3C00H
4A00           00120      END  4A00H

```

The first phase of the display—the portion showing line numbers at the left-hand end—is the part you entered as assembly program text. After doing the BREAK and answering the asterisk with an A, the listing shows the assembled version of the program.

The assembled portion shows addresses in the first vertical field, followed by hexadecimal versions of the machine code. The second field shows the original line numbers for convenient reference, and the third vertical field shows the original assembly text.

The important feature is the next-to-the-last line that reads: 4A00 21003C. This shows the address location of the beginning of a 3-byte, machine-language instruction, 21003C. And what does 21003C mean? It means load 3C00H to the HL register pair—an operation originally specified with the mnemonics LD HL,3C00H. Yes, indeed, EDTASM assembled LD HL,3C00H into machine code and assigned it to the ORG address 4A00H. This is the same result you would get by hand assembling that particular instruction.

Incidentally, if you are actually trying this experiment on your machine, you will see some additional material, namely:

```

00000 TOTAL ERRORS
READY CASSETTE

```

Those messages inform you that you have made no errors in the assembly listing and that the program has been successfully assembled into machine code. It is then time to save the machine code, or object code, on cassette tape. If you want to save that one-instruction, machine-language program on tape, all you have to do is set up the cassette player for recording a program and then answer the cursor symbol by striking the ENTER key.

When the assembled program is thus recorded, EDTASM will return an asterisk, indicating it is back to its monitor mode and awaiting the next command. But if you don't want to record the assembled program, simply strike the BREAK key to return to the monitor mode.

Now the system is ready to do something else. Perhaps you want to delete the old program. In that case answer the asterisk by entering D#:\*. Then you are ready to start all over again, doing

something such as I100,10 to begin writing a new assembly, source-code program.

### SUMMARY

- Striking the BREAK key interrupts the current activity and returns EDTASM to its monitor mode.
- *Inn,i* sets up the Insert mode, beginning reference line numbers at *nn* and incrementing each of them at intervals of *i* line numbers.
- ORG *addr* specifies the address, *addr*, where the machine-language program is to begin loading. Every assembly program must begin with that pseudo-operation.
- END *addr* specifies the address, *addr*, where the machine-language program is to begin execution. It is often the same as the ORG address, but not necessarily. As demonstrated later, the *addr* in this pseudo-op can be replaced by a line label.
- Entering an A from the monitor mode causes EDTASM to assemble your source-code listing, to convert the mnemonic-oriented source code into machine addresses and instructions.

This time, write a program that will actually do something. Begin by answering the asterisk with D#:\* and entering it. This ensures that the old program is wiped out. Then answer the asterisk with an I100,10 to begin at line number 100 and set up line increments of 10. Type in the assembly listing as shown in Program 9-1.

Chances are quite good you won't be able to type in the entire listing without making at least one typing error. If that happens, there are several procedures for correcting it. The following error-correcting procedure is not the simplest one available, but it is the

**Program 9-1. A typical source code listing for EDTASM.**

```

00100          ORG 4A00H
00110 TON      CALL 01C9H          ;CLEAR THE SCREEN
00120 PON      LD HL,3C00H         ;SET VIDEO POINTER
00130          LD (HL),42D         ;PRINT ASTERISK
00140          LD A,H              ;MSB OF POINTER TO A
00150          CP 3DH              ;END OF QUADRANT?
00160          JR Z,CTD            ;IF SO,DO DELAY
00170          INC HL              ;ELSE INCREMENT POINTER
00180          JR PON              ;AND PRINT MORE
00190 CTD      CALL TD              ;CALL TIME DELAY
00200          CALL 01C9H          ;CLEAR THE SCREEN
00210          CALL TD              ;CALL TIME DELAY
00220          JR TON              ;AND DRAW AGAIN
00230 TD       LD B,0CH            ;SET MSB OF TD
00240 SETC     LD C,0FFH           ;SET LSB OF TD
00250 DECC     DEC C               ;COUNT LSB
00260          JR NZ,DECC          ;IF NOT DONE, DO MORE
00270          DEC B               ;ELSE COUNT MSB
00280          JR NZ,SETC          ;IF NOT DONE,DO LSB AGAIN
00290          RET                ;ELSE RETURN
00300          END                4A00H

```

most straightforward in the context of what has been discussed.

To correct an error in the assembly listing:

1. Strike the BREAK key to return to EDTASM's monitor mode.
2. Delete the defective line by answering the asterisk with a *D*, followed by the appropriate line number. If the problem is in line 00180, for example, enter D180 (leading zeros need not be typed in).
3. Insert a corrected version of the line after answering the asterisk with an *I*, followed by the appropriate line number, for example, I180. The system is now in the Insert mode and at the line number you just entered. Type in the corrected version of the entire line. Inserting a corrected line will usually turn up a NO ROOM BETWEEN LINES message. Don't worry about that now; it is simply telling you that the next line in the sequence has some information in it.
4. With the system back in the monitor mode, as signaled by the asterisk prompt symbol, you can resume programming by entering *Inn*, where *nn* is the next line number you want to work with.

The error is corrected, and you are back on the right track for completing the job.

The program in Program 9-1 ultimately causes the system to draw four lines of asterisks across the top of the screen, making them flash on and off. But it is far too soon to check its operation.

Line 100 specifies the origin of the machine-language program at address 4A00H. Line 300, the last line, specifies the entry point at that same address. In other words, the first address for the program is the place where execution is to begin.

Line 110 calls the ROM-based operation for clearing the screen. Notice that the address concludes with the letter H. An H must follow any data or address that is specified in a hexadecimal format.

#### NOTE

Following a data or address number with an H specifies a hexadecimal number.

Following a data or address number with a D, or no letter at all, specifies a decimal format.

Following a data or address number with an O specifies an octal number.

Hexadecimal numbers must begin with a number, and not a letter. EDTASM, for example, will not accept data FFH; but it will understand its equivalent 0FFH.

So line 110 specifies a screen-clearing operation. It concludes with a comment. In assembly listings, comments serve the same purpose as REMs in BASIC; they are a convenient way to make some notes concerning the operations at hand. *Comments must be preceded by a semicolon.*

Line 120 sets the video pointer to the first video memory location, 3C00H. Now that line begins with a label. You can make up your own labels, following a few simple rules that are spelled out in the next NOTE. Their true purpose and assembly power won't be apparent, however, until we get further into this analysis.

#### NOTE

Labels are used to mark critical points in an assembly program.

A label must follow the reference line number.

A label can be any combination of alphanumeric characters, up to six and always beginning with an alphabetical character. A \$ must not be used anywhere in a label.

There are some reserved labels that must not be used. See Chart 9-1. This list of reserved labels corresponds to BASIC's list of reserved words that cannot be used as variable names.

**Chart 9-1. Reserved Labels for EDTASM\***

A	IX
B	IY
C	SP
D	PC
E	AF
H	BC
L	DE
I	HL
R	

\*These letters may be used within a label, but they must not stand alone as a label.

Line 130 prints the asterisk at the point indicated by the HL pair on the screen. Note that the ASCII character is specified in a decimal format in this case. A 2AH would work out equally well.

Skipping down to line 160, you find the instruction JR Z,CTD. That is saying: jump relative if zero to the line labeled CTD. You don't have to count the number of bytes to be jumped—just specify the label. This is at least an introductory note concerning the power of labels in EDTASM. The computer is going to figure out how many program bytes to be jumped.

See the same sort of things taking place in lines 180, 190, 210, 220, 260, and 280. Knowing how to use the labels, you should be able to work your way through the entire program to see how it works. Hopefully, the comments will be of great help, too.

In the process of discussing the assembly listing in Program 9-1, you should be getting some idea about how to use labels and comments. Both serve to simplify the programming operation and clarify what is going on.

Once the listing is in place as illustrated in Program 9-1, the next job is to let EDTASM assemble it for you. First, get back to the monitor by striking the BREAK key; then assemble the program by entering an A.

The assembled version of the program will be dumped onto the screen—far too quickly for you to study it. The important thing now, however, is the number of TOTAL ERRORS. If there are none, you've done a good job. But if that counter shows some errors, you have some editing work to do.

Program 9-2 is a hardcopy version of the assembled program. The program is too long to be seen in its entirety on the screen, so you will have to use Program 9-2 as a guide for the time being.

#### Program 9-2. Assembled version of Program 9-1.

```

4A00          00100      ORG      4A00H
4A00 CDC901  00110      CALL 01C9H      ;CLEAR THE SCREEN
4A03 21003C  00120 TON   LD HL,3C00H      ;SET VIDEO POINTER
4A06 362A    00130 PON   LD (HL),42D      ;PRINT ASTERISK
4A08 7C      00140      LD A,H           ;MSB OF POINTER TO A
4A09 FE3D    00150      CP 3DH           ;END OF QUADRANT?
4A0B 2803    00160      JR Z,CTD         ;IF SO, DO DELAY
4A0D 23      00170      INC HL           ;ELSE INCREMENT POINTER
4A0E 18F6    00180      JR PON          ;AND PRINT MORE
4A10 CD1B4A  00190 CTD   CALL TD         ;CALL TIME DELAY
4A13 CDC901  00200      CALL 01C9H      ;CLEAR THE SCREEN
4A16 CD1B4A  00210      CALL TD         ;CALL TIME DELAY
4A19 18E8    00220      JR TON          ;AND DRAW AGAIN
4A1B 060C    00230 TD    LD B,0CH        ;SET MSB OF TD
4A1D 0EFF    00240 SETC  LD C,OFFH       ;SET LSB OF TD
4A1F 0D      00250 DECC  DEC C           ;COUNT LSB
4A20 20FD    00260      JR NZ,DECC      ;IF NOT DONE,DO MORE
4A22 05      00270      DEC B           ;ELSE COUNT MSB
4A23 20F8    00280      JR NZ,SETC      ;IF NOT DONE, DO LSB AGAIN
4A25 C9      00290      RET            ;ELSE RETURN
4A00          00300      END      4A00H
00000 TOTAL ERRORS

DECC      4A1F
SETC      4A1D
TD         4A1B
CTD        4A10
PON        4A06
TON        4A03

```

Let's suppose that it turns out that you have made a syntactical error somewhere in the program; the TOTAL ERRORS message shows something other than 00000. This means you have to do some editing. But first you have to find the error.

If you don't know what the error is, the following procedure helps you locate it:

1. Get to the monitor by striking the BREAK key.
2. Answer the asterisk by entering A/WE.

Doing this, the system will begin assembling the program again. This time, however, the /WE part of the command tells the system to stop assembling when it encounters an error and write a message describing its nature. You can then figure out what the error is and correct the error by referring to the appropriate line number.

One method for correcting errors was described earlier. Basically it amounted to deleting the entire line of text containing the error and then going to the Insert mode to reenter the entire line from scratch. There is, however, a simpler way to do the same thing.

EDTASM has a Replace command. Answer the asterisk with an Rnn, where nn specifies the line number of the line to be replaced. The system automatically deletes the line for you and waits for you to type in the corrected version of the line. The R command takes the place of the Delete/Insert sequence.

The best way to edit a line having just one or two minor errors, however, is to take advantage of EDTASM's Edit mode of operation. This Edit feature works almost exactly the same way as BASIC's EDIT mode.

To get into the Edit mode, answer the asterisk with Enn, where nn is the line number of the text to be edited. Then use the BASIC-like EDIT commands to fix the trouble. The Edit instructions in your *Editor/Assembler Instruction Manual* adequately describe all the Edit commands you will ever need.

Continue doing the A/WE process until there are no longer any errors in the text. At that time the listing on the screen will resemble that in Program 9-2.

Notice, incidentally, that the assembled listing ends by specifying the absolute addresses of all the labeled instructions. Label DECC, for example, refers to the machine instruction beginning at address 4A1F. Such a listing can be invaluable when it comes to debugging a very long program.

## WORKING WITH THE OBJECT CODE VERSION

After going through the processes described in the previous section, the program is still not ready to be tested on line. It must be



transferred to cassette tape first and then loaded back into the machine under the SYSTEM command.

Assembled programs should be saved on tape under a file name. If you do not specify a file name, EDTASM will automatically assign NONAME as a file name. But as long as you have control over matters, it is a good idea to assign a name that means something significant to you. The following NOTES show a few rules for composing file names for machine programs written under EDTASM.

#### NOTES

A file name can be composed of combinations of alphanumeric characters, not exceeding six characters in length.

The first character must be an alphabetical character.

To assign a file name to a program, get into the monitor mode and then answer the asterisk by typing an A, followed by one space and your chosen file name.

The A instructs the system to assemble the program again, but under the specified *file name*. A space between the A and *file name* is mandatory.

After assembling the program under the designated file name, the system brings up the message READY CASSETTE. Set up your tape machine for recording, then strike the ENTER key. Within just a few moments the system will load the machine-language portion of the program to the tape. *None of the assembly text is loaded this way.*

Now, if you are perfectly confident that the program will run as you expect, and if you are sure you will never want to edit it from EDTASM at any time in the future, you are done with the assembly operations.

But you might want to save the assembly text for future reference. To save the source-code part of the program—the stuff you typed into the machine originally—do this:

1. Set up the tape machine for recording a program.
2. From the monitor mode, enter W *file name*, where *file name* is a name composed according to the rules outlined earlier.
3. Answer READY CASSETTE by striking the ENTER key.

When the recording is done, you have the assembly text saved on tape.

The object code listing—the actual machine-language program—is always recorded immediately after doing an assembly operation.

The source code listing can be saved on tape only by going to the Write mode—entering *W file name*.

Suppose, for instance, you want to save the object code from Program 9-2 under the file name OBLINK. To do that, enter A OBLINK; and when you see the READY CASSETTE, record the program by striking the ENTER key.

If you want to save the source listing under the file name SBLINK, enter W SBLINK, and record it by striking the ENTER key in response to READY CASSETTE.

By this time you have composed and assembled a program with EDTASM. The machine-language version is saved on cassette tape as OBLINK, and the source version is saved as SBLINK. Now it is time to test the program's actual operation—finally.

Unfortunately, EDTASM has to be destroyed. To get out of EDTASM, answer the asterisk by entering a *B*. Entering a *B* from EDTASM returns things to the power-up sequence for BASIC. You will thus see MEMORY SIZE?

Then set up the cassette player for entering the machine-language version of the program—the one saved as OBLINK. Enter SYSTEM, and respond to the asterisk by typing OBLINK. The tape will load the program into the machine and signal the end of the loading process by showing the usual SYSTEM symbols \*?\_... Respond by entering the slash (/)—and, presto, there are the flashing lines of asterisks.

But wait! There are a couple of problems. EDTASM checked for syntactical error in the assembly listing, but it did not (and cannot) check the validity of the program, itself. One problem is that the asterisks appear to be flashing too rapidly; maybe you would like to see them flashing at a somewhat lower rate. Then, too, there is a programming problem that causes an extra asterisk to appear at the beginning of the fifth line on the screen—it's a "small" problem, but one of a type that could cause some real trouble under different circumstances.

So there is a need to debug the program; the original version isn't quite good enough.

Exactly how to go about fixing those problems is often a matter of opinion. One approach is to load T-BUG and fix the troubles, using the original EDTASM as a guide to the relevant address locations. This approach is a good one if the "fixes" aren't very extensive (as is the case here). Another nice thing about fixing things with T-BUG is that you can test the program on the spot. The only disadvantage is that your original EDTASM listings are no longer valid.

The second approach is to reload EDTASM, load the source version of the program into it, and fix it at the assembly level. That

way, you end up with fresh listings; but then you have to do a lot of recording and reloading to see whether or not the “fixes” really do the job.

In this particular case it seems better to load T-BUG. Load T-BUG and refer to the listing in Program 9-2 to see what has to be done to fix the troubles.

To make the time delay longer, and hence get a lower flashing rate, register B in line 230 can be loaded with CFH, instead of 0CH. So change address 4A1C from 0C to CF. Do a J 4A00, and you will see that did, indeed, lower the flashing rate.

The fact that an extra asterisk appears at the beginning of the fifth line comes about as the result of a common sort of programming error. Referring to Program 9-2, incrementing the HL register pair should take place *before* the H register is tested for a value of 3D. Thus the operations beginning at line 130 should read:

4A06	362A	00130	PON	LD (HL),42D	;PRINT ASTERISK
4A08	23	00140		INC HL	;INCREMENT POINTER
4A09	7C	00150		LD A,H	;MSB OF POINTER TO A
4A0A	FE 3D	00160		CP 3DH	;END OF QUADRANT?
4A0C	28 02	00170		JR Z,CTD	;IF SO, DO DELAY
4A0E	18 F6	00180		JR PON	;ELSE PRINT MORE

Everything else is the same before and after this sequence.

In this case we were lucky enough to have the “fix” fit exactly within the memory space already allocated for the drawing sequence. Had there not been enough room, it would be necessary to do the fix from EDTASM.

So enter the “fix” from T-BUG, changing the data content of addresses 4A06 through 4A0F as necessary. Doing a J 4A00 will show that the program works better this time.

The machine-code version can be saved on cassette tape from T-BUG by using the Punch command. In many instances this marks the end of the programming job. But, for the sake of illustration, suppose you want to make the changes from EDTASM instead of T-BUG. You might use T-BUG to confirm the validity of your changes, but you want to make corrected versions of the source listings. Here is how that is done.

First, reload EDTASM from scratch, as described in the opening part of this chapter. Then, to save some time, load the defective source code listing.

To load a source listing under EDTASM, answer the asterisk with L *file name*. If you specify no *file name*, the system will accept the first program loaded from cassette tape. You could, in this instance, load under the file name, SBLINK—that was the name attached to the original source listing.

When the loading is done, as signaled by a new asterisk prompt symbol, you can confirm entry of the right program by entering P#:\* . That will display the entire source listing on the screen.

#### NOTE

Under the EDTASM monitor, a P command causes a printing of the source listing. See the *Editor/Assembler Instruction Manual* for some useful variations of the Print command.

The system is now set up so that you can make the necessary changes in the source listing, using Deletes, Inserts, Replaces, and Edit commands. A corrected and assembled version of the program appears as Program 9-3.

### OTHER EDTASM COMMANDS AND VARIATIONS

With this sort of experience behind you, you should be able to learn a lot more from the *Editor/Assembler Instruction Manual*.

**Program 9-3. Assembled and corrected version of Program 9-2.**

4A00		00100	ORG	4A00H	
4A00	CDC901	00110	CALL	01C9H	;CLEAR THE SCREEN
4A03	21003C	00120	LD	HL,3C00H	;SET VIDEO POINTER
4A06	362A	00130	PON	LD (HL),42D	;PRINT ASTERISK
4A08	23	00140	INC	HL	;INCREMENT POINTER
4A09	7C	00150	LD	A,H	;MSB OF POINTER TO A
4A0A	FE3D	00160	CP	3DH	;END OF QUADRANT?
4A0C	2802	00170	JR	Z,CTD	;IF SO,DO DELAY
4A0E	18F6	00180	JR	PON	;AND PRINT MORE
4A10	CD1B4A	00190	CTD	CALL TD	;CALL TIME DELAY
4A13	CDC901	00200		CALL 01C9H	;CLEAR THE SCREEN
4A16	CD1B4A	00210		CALL TD	;CALL TIME DELAY
4A19	18E8	00220		JR TON	;AND DRAW ACAIN
4A1B	060C	00230	TD	LD B,0CH	;SET MSB OF TD
4A1D	0EFF	00240	SETC	LD C,0FFH	;SET LSB OF TD
4A1F	0D	00250	DECC	DEC C	;COUNT LSB
4A20	20FD	00260		JR NZ,DECC	;IF NOT DONE,DO MORE
4A22	05	00270		DEC B	;ELSE COUNT MSB
4A23	20F8	00280		JR NZ,SETC	;IF NOT DONE, DO LSB AGAIN
4A25	C9	00290		RET	;ELSE RETURN
4A00		00300	END	4A00H	
00000	TOTAL	ERRORS			
DECC	4A1F				
SE IC	4A1D				
TD	4A1B				
CTD	4A10				
PON	4A06				
TON	4A03				

See what you can learn on your own from the following sections of that manual:

- Assemble (A)
- BASIC (B)
- Delete (D)
- Edit (E)
- Find (F)
- Hardcopy (H)—line-printer applications only
- Insert (I)
- Load (L)
- Number (N)
- Print (P)
- Replace (R)
- Type (T)—line-printer applications only
- Scroll up (↑)
- Scroll down (↓)
- Tab (→)
- Delete character (←)
- Delete line (shift ←)
- Pause (shift @)
- Write (W)

The best way to get familiar with EDTASM is by actually working with it, composing interesting programs of your own design. But if you think EDTASM is a powerful machine-language programming tool at this point, wait until you have a chance to get into the discussions in the next chapter.

## CHAPTER 10

# Real Assembly Power With Pseudo-Ops

In the context of writing assembly programs, pseudo-ops are operations that affect the assembly operation but do not appear as legitimate machine-language instructions. Chapter 9 specified two pseudo-ops: ORG and END.

ORG and END are operations used by the assembler, but they do not appear in the assembled, object-code version of the listing. In a sense they are housekeeping operations for the assembler; they are important, to be sure, but they are not part of the finished product—the fully assembled, ready-to-run, machine program.

EDTASM supports a number of other pseudo-ops. They are not mandatory, as are the ORG and END operations; rather, they save the programmer a great deal of time and effort.

### SIMPLE EQU PSEUDO-OPS

One of the simplest, yet most used, optional pseudo-ops is EQU. It is used to assign a specific data or address value to a label of your own invention. Consider the following line of assembly text:

```
00100   CLS   EQU   01C9H   ;DEFINE CLS
```

In this example, label CLS is assigned the 2-byte, hexadecimal address 01C9. That particular address, you may recall, is the entry point for a ROM-based TRS-80 operation that clears the screen and homes the cursor.

Without taking advantage of the EQU pseudo-op, setting up an instruction for clearing the screen meant doing a CALL 01C9H

in the assembly program. And each time you wanted to use that instruction, you would have to type it in full.

But by doing a CLS EQU 01C9H at some early point in the assembly listing, writing an instruction for clearing the screen is a simple matter of doing CALL CLS. Once a number is assigned to a label via an EQU pseudo-op, any subsequent references to that label automatically bring up the value originally assigned to it. One advantage is that a label such as CLS is certainly easier to remember than 01C9H when you are in the middle of writing an assembly program.

**Program 10-1. Source and assembled versions of a program that uses simple EQU operations to define labels.**

```

00100          ORG          4A00H
00110 CLS      EQU          01C9H
00120 KBIN     EQU          0049H
00130 PCHR     EQU          033AH
00140 ;BEGINNING OF THE PROGRAM
00150          CALL CLS      ;CLEAR THE SCREEN
00160 CET      CALL KBIN    ;CET A KEYBOARD CHARACTER
00170          CALL PCHR     ;AND PRINT IT
00180          JR CET       ;CET NEXT CHARACTER
00190          ENO          4A00H

4A00          00100      ORG          4A00H
01C9          00110 CLS  EQU          01C9H
0049          00120 KBIN EQU          0049H
033A          00130 PCHR EQU          033AH
00140 ;BEGINNING OF THE PROGRAM
4A00 C0C901    00150      CALL CLS      ;CLEAR THE SCREEN
4A03 C04900    00160 CET  CALL KBIN    ;CET A KEYBOARD CHARACTER
4A06 C03A03    00170      CALL PCHR     ;AND PRINT IT
4A09 18F8      00180      JR CET       ;CET NEXT CHARACTER
4A00          00190      ENO          4A00H
00000 TOTAL ERRORS

CET  4A03
PCHR 033A
KBIN  0049
CLS   01C9

```

By way of a convincing demonstration, type in the assembly listing in the first part of Program 10-1. That listing begins with the mandatory ORG pseudo-op, but then there are three optional EQU pseudo-ops.

The EQU operation in line 110 defines label CLS as 01C9H, the KBIN label in line 120 is defined as 0049H (entry point for a keyboard character to the accumulator), and PCHR in line 130 is a

label assigned to 033AH (entry point for printing a character from the accumulator and advancing the cursor).

With the critical entry addresses thus defined as labels, the working part of the program, lines 150 through 180, uses the labels to work out a simple assembly program: one that clears the video screen and then allows you to type and print characters from the keyboard.

The exciting part of the job doesn't become apparent, however, until the listing is assembled. See the second part of Program 10-1.

In line 150, for example, CALL CLS is assembled into machine code CDC901. Sure enough, there is the actual address that was assigned to label CLS in an earlier EQU operation (in line 110). And in lines 160 and 170 you can see that the assembler kept track of the other addresses as well.

Now there is no longer any need to cite particular addresses or data values over and over again in an assembly listing. Just assign the values to a label with an EQU pseudo-op and then cite that label whenever the number is needed at any later time.

Incidentally, EDTASM allows you to assign decimal numbers to labels via the EQU operation. This is a handy feature in instances where you happen to know a value in decimal terms but don't want to make the effort to translate it into hexadecimal notation. See Program 10-2.

**Program 10-2. Assembled listing for a program that illustrates the fact that EQU can define labels in hexadecimal or decimal terms.**

7000	00100	DRG	7000H	
01C9	00110 CLS	EQU	01C9H	
00BF	00120 CHAR	EQU	191	
	00130 ;BECINN INC OF PRDGRAM			
7000	CDC901	00140	CALL CLS	;CLEAR THE SCREEN
7003	060A	00150	LD B,10	;SET COUNTER
7005	21003C	00160	LD HL,3C00H	;INITIALIZE VIDEO POINTER
7008	36BF	00170 PLOT	LD (HL),CHAR	;PLOT A RECTANGLE
700A	05	00180	DEC B	;CDUNTDOWN
700B	20FB	00190	JR NZ,PLOT	;IF NOT DONE, PLDT MORE
700D	18FE	00200 LOOP	JR LOOP	;ELSE LOOP TD SELF
7000		00210	ENO	7000H
00000	TOTAL ERRORS			
LOOP	700D			
PLOT	7008			
CHAR	00BF			
CLS	01C9			

In line 120, label CHAR is assigned, by an EQU operation, the decimal value 191. Radio Shack lists the graphics codes in decimal, and this sort of assignment operation avoids the task of determining



the hexadecimal version. But notice how the assembler does that job for you in the machine code in line 170.

## EQU OPERATIONS WITH MATH EXPRESSIONS

Simple EQU pseudo-ops let you define a label as some specific numerical value, but there is more to it. EQU also allows assignments of values that are expressed in a mathematical form. In BASIC, for example, you can do something such as PRINT A, but if A has some preassigned value, you can also do PRINT A+2. That same sort of thing is possible with the EQU pseudo-op.

Suppose you are writing an assembly program that calls for plotting a character at the beginning and end of the top line on the video screen. You probably know the beginning address, perhaps because you've used it so many times in the past. Thus a line such as

```
TSTART    EQU    3C00H
```

doesn't put much of strain on your brain. But then you must also specify the ending address of that line—and you cannot remember it. Of course, you can figure it out on paper or look it up in a table of video addresses; it is easier, however, to do something such as this:

```
TSTART    EQU    3C00H
TEND      EQU    TSTART+63
```

Recalling that the end of a horizontal line on the screen is displaced 63 (decimal) character spaces above the starting address of that line, you can define TEND—the end of the first line—as TSTART plus 63. And after defining TEND in that fashion, citing it later in the program will call up the proper address, 3C3F.

Yes, indeed, the EQU pseudo-op can attach the result of a math expression to a label name, and, what's more, part of that expression can be another label defined elsewhere in the program.

The EQU pseudo-op can handle both addition and subtraction expressions. It is possible to use any number and combinations of + and - operators. The assembler always executes the operators in a strictly left-to-right fashion, but it cannot handle a grouping of operations by parentheses.

See an application in Program 10-3.

The purpose of Program 10-3 is to plot a rectangle of light in the corners of the screen. The programmer happened to know the address of the start of video memory (3C00H) and the end (3FFFH) but not the addresses for the points at the upper-right and lower-left corners. Rather than taking the time and effort to work it out on paper, the programmer defined those points with EQU math expressions. See lines 120 and 140.

**Program 10-3. Source and assembled versions of a program that shows how labels can be defined with EQU and math expressions.**

```

00100          ORG          4A00H
00110 TSTART   EQU          3C00H  ;DEFINE START OF TOP LINE
00120 TEND     EQU          TSTART+63  ;DEFINE END OF TOP LINE
00130 BEND     EQU          3FFFH  ;DEFINE END OF BOTTOM LINE
00140 BSTART   EQU          BEND-63  ;DEFINE START OF BOTTOM LINE
00150 CHAR     EQU          191      ;DEFINE CHARACTER
00160 ;BEGINNING OF PROGRAM
00170          CALL 01C9H  ;CLEAR THE SCREEN
00180          LD HL,TSTART ;POINT TO UPPER LEFT
00190          LD (HL),CHAR ;PLOT CHAR
00200          LD HL,TEND   ;POINT TO UPPER RIGHT
00210          LD (HL),CHAR ;PLOT CHAR
00220          LD HL,BSTART ;POINT TO LOWER LEFT
00230          LD (HL),CHAR ;PLOT CHAR
00240          LD HL,BEND   ;POINT TO LOWER RIGHT
00250          LD (HL),CHAR ;PLOT CHAR
00260 LOOP     JR LOOP      ;LOOP TO SELF
00270          END          4A00H

4A00          00100      ORG          4A00H
3C00          00110 TSTART EQU          3C00H  ;DEFINE START OF TOP LINE
3C3F          00120 TEND  EQU          TSTART+63  ;DEFINE END OF TOP LINE
3FFF          00130 BEND  EQU          3FFFH  ;DEFINE END OF BOTTOM LINE
3FC0          00140 BSTART EQU          BEND-63  ;DEFINE START OF BOTTOM LINE
00BF          00150 CHAR  EQU          191      ;DEFINE CHARACTER
              00160 ;BEGINNING OF PROGRAM
4A00 CDC901  00170          CALL 01C9H  ;CLEAR THE SCREEN
4A03 21003C  00180          LD HL,TSTART ;POINT TO UPPER LEFT
4AD6 36BF    00190          LD (HL),CHAR ;PLOT CHAR
4A08 213F3C  00200          LD HL,TEND   ;POINT TO UPPER RIGHT
4A0B 36BF    00210          LD (HL),CHAR ;PLOT CHAR
4A0D 21C03F  00220          LD HL,BSTART ;POINT TO LOWER LEFT
4A10 36BF    00230          LD (HL),CHAR ;PLOT CHAR
4A12 21FF3F  00240          LD HL,BEND   ;POINT TO LOWER RIGHT
4A15 36BF    00250          LD (HL),CHAR ;PLOT CHAR
4A17 18FE    00260 LOOP     JR LOOP      ;LOOP TO SELF
4A00          00270          END          4A00H
00000 TOTAL ERRORS

LOOP  4A17
CHAR  00BF
BSTART 3FC0
BEND 3FFF
TEND 3C3F
TSTART 3C00

```

The second part of the listing is the assembled version of the same program. Note the object codes in lines 200 and 220—those lines calling for the math-generated TEND and BSTART labels. You can see that the assembler did the figuring for the programmer, coming up with a hexadecimal value of 3C3F for the upper-right corner and 3FC0 for the lower-left corner.

## REDEFINING A LABEL WITH DEFL

A label that is defined by an EQU pseudo-op is committed to the designated value through the entire program listing. Any attempt to change the value assigned to a label will result in a multiple-definition assembly error message.

There are occasions, however, when it is desirable to change the definition of a label during the course of the assembly process. Assigning values to labels, using the DEFL pseudo-op instead of EQU, allows this sort of redefinition to take place.

The example in Program 10-4 plots a character at the four corners of the screen, just as Program 10-3 does. In this example, however, the character is changed each time.

**Program 10-4. Assembled program showing an application of the DEFL pseudo-op.**

```

4A00          00100      ORG      4A00H
3C00          00110 TSTART EQU      3C00H ;DEFINE START OF TOP LINE
3C3F          00120 TEND  EQU      TSTART+63 ;DEFINE END OF TOP LINE
3FFF          00130 BEND  EQU      3FFFH ;DEFINE END OF BOTTOM LINES
3FC0          00140 BSTART EQU      BEND-63 ;DEFINE START OF BOTTOM LINE
              00150 ;BEGINNING OF THE PROGRAM
4A00 C0C901 00160      CALL 01CH9 ;CLEAR THE SCREEN
4A03 21003C 00170      LD HL,TSTART ;POINT TO UPPER LEFT
0041          00180 CHAR   DEFL 65 ;DEFINE 'A'
4A06 3641 00190      LD(HL),CHAR ;PLOT FIRST CHARACTER
4A08 213F3C 00200      LD HL,TEND ;POINT TO UPPER RIGHT
0042          00210 CHAR   OEFL 66 ;DEFINE 'B'
4A0B 3642 00220      LD (HL),CHAR ;PLOT SECOND CHARACTER
4A0D 21C03F 00230      LD HL,BSTART ;POINT TO LOWER LEFT
0043          00240 CHAR   DEFL 67 ;DEFINE 'C'
4A10 3643 00250      LD (HL),CHAR ;PLOT THIRD CHARACTER
4A12 21FF3F 00260      LD HL,BEND ;POINT TO LOWER RIGHT
0044          00270 CHAR   DEFL 68 ;DEFINE 'D'
4A15 3644 00280      LD (HL),CHAR ;PLOT FOURTH CHARACTER
4A17 18FE 00290 LOOP   JR LOOP ;LOOP TO SELF
4A00          00300      END      4A00H
00000 TOTAL ERRORS

LOOP 4A17
CHAR 0044
BSTART 3FC0
BEND 3FFF
TEND 3C3F
TSTART 3C00

```

In line 180, the DEFL pseudo-op defines CHAR as decimal 65, the letter A. Line 210 changes the definition to decimal 66, line 240 uses another DEFL to change CHAR to 67, and finally line 270 sets up the final value of 68.

Of course, it would be easier in this case to omit the CHAR definitions altogether, loading the desired value in an immediate fashion—LD (HL),65, for example. But this is merely an illustrative

example. And, indeed, the assembler accepts the redefinition of CHAR in each case.

The DEFL pseudo-op also accepts simple math expressions as shown in Program 10-5. It accepts the same expressions as the EQU pseudo-op, and the rules and limitations are the same as well.

**Program 10-5. Assembled listing of a program that demonstrates the use of math expressions with DEFL pseudo-ops.**

```

4A00          00100      ORC      4A00H
3C00          00110 TSART  EQU      3C00H ;DEFINE START OF TOP LINE
3C3F          00120 TEND  EQU      TSART+63 ;DEFINE END OF TOP LINE
3FFF          00130 BEND  EQU      3FFFH ;DEFINE END OF BOTTOM LINE
3FC0          00140 BSTART EQU      BEND-63 ;DEFINE START OF BOTTOM LINE
              00150 ;BEGINNING OF THE PROGRAM
4A00 CDC901    00160          CALL 01C9H ;CLEAR THE SCREEN
4A03 21003C    00170          LD HL,TSART ;POINT TO UPPER LEFT
0041          00180 CHAR  DEFL 65 ;DEFINE 'A' CHARACTER
4A06 3641      00190          LD (HL),CHAR ;PLOT FIRST CHARACTER
4A08 213F3C    00200          LD HL,TEND ;POINT TO UPPER RIGHT
0042          00210 CHAR  DEFL CHAR+1 ;REDEFINE CHAR
4A0B 3642      00220          LD (HL),CHAR ;PLOT SECOND CHARACTER
4A0D 21FF3F    00230          LD HL,BSTART ;POINT TO LOWER LEFT
0043          00240 CHAR  DEFL CHAR+1 ;REDEFINE CHAR
4A10 3643      00250          LD (HL),CHAR ;PLOT THIRD CHARACTER
4A12 21FF3F    00260          LD HL,BEND ;POINT TO LOWER RIGHT
0044          00270 CHAR  DEFL CHAR+1 ;REDEFINE CHAR
4A15 3644      00280          LD (HL),CHAR ;PLOT FOURTH CHARACTER
4A17 18FE      00290 LOOP  JR LOOP ;LOOP TO SELF
4A00          00300          END      4A00H
00000 TOTAL  ERRORS

LOOP 4A17
CHAR 0044
BSTART 3FC0
BEND 3FFF
TEND 3C3F
TSART 3C00

```

It is important to bear in mind that the DEFL pseudo-op alters the assembly process and not the execution of the machine-language program being generated. A label that is defined by a DEFL will carry the assigned value through every line of the assembly listing, at least until the label is redefined by another DEFL. Note in Programs 10-4 and 10-5 that CHAR in the label lists carries a value of 0044. Values assigned by DEFL earlier in the program are not listed.

## LEAVING MEMORY SPACE WITH DEFS

In earlier chapters dealing with T-BUG programming, it was explained why it is desirable to leave small blocks of unused memory at strategic places in a machine-language program. The general idea was to leave some room for making revisions that might call

for fitting in a few extra instructions. As described thus far, EDTASM does not leave any room between instructions. But that can be done with the DEFS pseudo-op.

Suppose you are writing a program under EDTASM and you want to leave eight successive memory locations blank. One part of the program is to end with eight address locations that are uncommitted and pick up with another part of the program after that. A line such as SPACE DEFL 8 will do the job. See the illustration of this effect in Program 10-6.

**Program 10-6. Assembled listing demonstrating the use of DEFS to allocate some uncommitted memory bytes.**

```

4A00          00100      DRG      4A00H
4A00  21003C  00110      LD HL,3C00H      ;SET VIDEO POINTER
4A03  CDC901  00120 TON  CALL  01C9H      ;CLEAR THE SCREEN
4A06  CD1A4A  00130      CALL TIME      ;DD TIME DELAY
4A09  36BF    00140      LD (HL),191     ;TURN DN THE LIGHT
4A0B  CD1A4A  00150      CALL TIME      ;DO TIME DELAY
4A0E  18F3    00160      JR TDN         ;AND START ALL OVER
000A          00170 SPACE DEFS 10        ;LEAVE SDME MEMORY SPACE
          00180 ;TIME DELAY SUBROUTINE
4A1A  06CF    00190 TIME LD B,0CFH      ;SET MSB OF TIME DELAY
4A1C  0EFF    00200 SETC LD C,0FFH      ;SET LSB OF TIME DELAY
4A1E  0D      00210 DECC DEC C         ;COUNT LSB
4A1F  20FD    00220      JR NZ,DECC     ;IF NDT DDNE, CDUNT AGAIN
4A21  05      00230      DEC B         ;ELSE CDUNT MSB
4A22  20F8    00240      JR NZ,SETC     ;IF NOT DONE,CDUNT MORE
4A24  C9      00250      RET           ;ELSE RETURN
4A00          0026D      END      4ADDH
00000  TDTAL ERRORS

DECC  4A1E
SETC  4A1C
SPACE 4A10
TIME  4A1A
TON   4A03

```

This program causes a small rectangle of light to flash on and off in the upper left-hand corner of the screen. A TIME subroutine, occupying assembler lines 190 through 250, sets the flashing rate.

Line 120 effectively turns off the rectangle of light by clearing the entire screen, and then line 130 calls the TIME delay subroutine. After that, line 140 turns on the light, and then line 150 calls the TIME delay again. The instruction in line 160 returns operations back to line 120, thus keeping the program running in an endless loop.

The whole point of the program, however, is to demonstrate the application of DEFS. See the assembly instruction in line 170—SPACE DEFS 10. That assembly instruction tells the system to leave 10 consecutive memory locations uncommitted. The subsequent TIME subroutine thus begins at address 4A1AH, instead of

4A10H—an address that would immediately follow the 2-byte instruction in line 160.

A DEFS pseudo-op must be preceded by a label in the first vertical field. In the example cited here, the label happens to be SPACE. It can be any valid label name that isn't used elsewhere in the program.

DEFS also supports simple math expressions as described for EQU and DEFL. This provides a means for leaving some uncommitted memory and picking up the next part of the program at some well-defined address. Suppose, for example, a section of the program leaves off at address 4A0FH. You don't care how many successive memory locations are left unused after that, just so the next part of the program picks up at address 4A20H. Letting the assembler figure the number of blank memory locations for you, you can enter SPACE DEFL 4A20H-4A0FH-1. That will do it.

#### NOTE

To leave some uncommitted memory space and pick up the program again at a well-defined address, use a DEFL with the expression:

new starting address - old ending address - 1.

You can use DEFL to leave any number of gaps in the machine-language addressing. The labels preceding the DEFL pseudo-op must be different in each case, however.

### DEFINING MEMORY CONTENTS WITH DEFB AND DEFW

Certain classes of machine-language programs require fetching data from a table of some sort. Each item in the table has a specific address, and it is fetched from the table by loading from its address. EDTASM's DEFB and DEFW pseudo-ops allow you to build tables of 1- and 2-byte data, respectively.

Program 10-7 is another program that really doesn't do anything useful, but it does illustrate the building of a small data table.

The data table in this instance consists of the ASCII codes for letters A, B, and C. The table is built by means of the DEFB pseudo-ops in lines 110, 120, and 130. For purposes of further demonstration the three characters are defined in three different ways: as a decimal version of A (DEFB 65), a hexadecimal version of character B (DEFB 42H), and as a single-character string enclosed in apostrophes (DEFB 'C').

**Program 10-7. Assembled listing that shows the use of the DEFB pseudo-op.**

```

4A00          00100      ORG      4A00H
4A00 41       00110      DEFB     65      ; 'A' CHARACTER
4A01 42       00120      DEFB     42H     ; 'B' CHARACTER
4A02 43       00130      DEFB     'C'     ; 'C' CHARACTER
000D         00140 BLAND DEFS  4A10H-4A02H-1 ; LEAVE SOME SPACE
          00150 ; BEGINNING OF 'PROGRAM'
4A10 3A004A   00160 PROG  LD A,(4A00H)    ; 'A' TO ACCUMULATOR
4A13 3A014A   00170      LD A,(4A01H)    ; 'B' TO ACCUMULATOR
4A16 3A024A   00180      LD A,(4A02H)    ; 'C' TO ACCUMULATOR
4A19 18FE     00190 LOOP  JR LOOP        ; LOOP TO SELF
4A10         00200      END      PROG
00000 TOTAL ERRORS

LOOP 4A19
PROG 4A10
BLAND 4A03

```

What is more important, however, is the fact that the assembler assigns the ASCII codes to address locations 4A00H through 4A02H. *A DEFB operation inserts the data byte at the current address.* In this particular instance the ORG sets the beginning of the program at 4A00H, and since the following operation is a DEFB the data byte thus defined is automatically assigned to address 4A00H. The second byte then goes to 4A01H, and the third to 4A02H.

This simple, three-character table concludes with some uncommitted memory space that is generated by the DEFS operation in line 140.

The actual operating program picks up at line 160—at the PROG label. It then proceeds to cite the addresses of the data in the table, calling the content of those three addresses to the accumulator. So once the data table is built, the contents can be called by merely citing the appropriate addresses. The only trick, as far as the programmer is concerned, is to keep track of where the desired data byte is located in the table.

The data table does not have to be located at the beginning of the machine-language listing. It can be inserted anywhere you choose. But since it is located at the beginning of the listing in Program 10-7, it is absolutely necessary to use the type of END statement specified in line 200.

When the TRS-80 is executing a program generated under EDTASM, it must not attempt to execute the contents of a data table. The table of data must not be treated as instructions—doing that ensures a total blowup of the program. So to avoid this situation the END statement specifies the entry point of the actual, executable part of the program at label PROG, which happens to be at address 4A10H in this case. This procedure is based on the notion that the END statement specifies the entry point of the program; and this

happens to be the first example in the current series of discussions where it is necessary to start the machine-language listing and execution of the program at two distinctly different places.

The DEFW pseudo-op works just like the DEFB operation, except that the former allocates two successive byte locations in the data table. DEFW is especially useful for building a table of 2-byte address locations.

Program 10-8 uses DEFW pseudo-ops to build a table of 2-byte addresses. The table begins at address 4A00H and runs through 4A07H. See lines 120 through 150.

The four addresses entered into that table point to the corners of the video screen. The first 2 bytes point to the upper left-hand corner, the next 2 point to the upper right-hand corner, then 2 point to the lower left corner, and the last one points to the lower right. Note that it is possible to define a "word" with a simple math expression (line 140).

After building this sort of address table the program can call the 2-byte contents as illustrated in lines 210 through 240. In this particular example the program plots small rectangles (TRS-80 graphic 191) in the four corners.

## BUILDING MESSAGE TABLES WITH DEFM

DEFB allocates a byte of memory for a data byte, DEFW allocates two successive bytes of memory for 2-byte data or address "words," and DEFM sets aside an indefinitely large section of

**Program 10-8. Assembled listing demonstrating the use of DEFW.**

```

4A00          00100      ORG      4A00H          ;SET ORIGIN
              00110 ;VIDE0 PDINT ADDRESS TABLE
4A00 003C     00120      DEFW     3C00H          ;UPPER LEFT
4A02 3F3C     00130      DEFW     3C3FH          ;UPPER RIGHT
4A04 C03F     00140      DEFW     3FFFH-63       ;LOWER LEFT
4A06 FF3F     00150      DEFW     3FFFH          ;LOWER RIGHT
00BF         00160 CHAR EQU      191            ;DEFINE CHAR
0008         00170 BLANK DEFS    4A10H-4A07H-1   ;LEAVE SOME MEMDRY
              00180 ;START OF PROGRAM
4A10 CDC901   00190 DRAW CALL    01C9H          ;CLEAR THE SCREEN
4A13 3EBF     00200      LD A,CHAR              ;SET CHARACTER
4A15 32004A   00210      LD (4A00H),A          ;PRINT UPPER LEFT
4A18 32024A   00220      LD (4A02H),A          ;PRINT UPPER RIGHT
4A1B 32044A   00230      LD (4A04H),A          ;PRINT LOWER LEFT
4A1E 32064A   00240      LD (4A06H),A          ;PRINT LOWER RIGHT
4A21 18FE     00250 LODP JR LOOP                ;LOOP TO SELF
4A10         00260      END      DRAW
00000 TOTAL ERRDRS

LODP 4A21
DRAW 4A10
BLANK 4A08
CHAR 00BF

```



memory for string messages. It would be possible to build a message table using DEFB pseudo-ops, but using a single DEFM line is much simpler. Consider the following comparison:

```
DEFB 'H'  
DEFB 'E'  
DEFB 'L'  
DEFB 'L'  
DEFB 'O'
```

Now, if you write some programming that calls the 5 bytes of data memory generated by that sequence of DEFB operations and prints them in sequence on the screen, you will end up printing HELLO. But using a single DEFW operation allows exactly the same thing to happen in a simpler fashion:

```
DEFW 'HELLO'
```

This will set up the message, HELLO, in some data memory locations, and the characters can be called to the video screen in sequence.

Programs 10-9 and 10-10 illustrate an application of the DEFM pseudo-op. Program 10-9 is the assembly listing as it appears after entering it, and Program 10-10 is the assembled version.

The program prints three messages: HELLO, NOW YOU SEE IT, and NOW YOU DON'T. You can find the HELLO designated with a DEFM pseudo-op in line 110 of Program 10-9, and NOW YOU SEE IT is defined in a similar fashion in line 130. Note that those messages are enclosed in apostrophes (as opposed to string-defining quotes as used in BASIC).

The fact that the NOW YOU DON'T message contains an apostrophe thus brings up a minor problem: The assembler will interpret the apostrophe between the N and T in that message as a closing apostrophe for the message, and it thinks you want to spell out NOW YOU DON.

Lines 150 through 170 show how to get around that particular difficulty. Line 150 uses a DEFM operation to spell out the message up to the point where the apostrophe is to occur. Line 160 then uses a DEFB pseudo-op to designate the ASCII code for an apostrophe. Finally, line 170 uses a DEFB to attach the character T to the end of the message.

Also notice that each of the three messages concludes with a DEFB 0 operation. The idea is to mark the end of each message with some sort of unique character or character code. The zero, or null character, is used here because it is the same one used to mark the end of text statements in BASIC. (Recall the discussion of BASIC text in Chapter 4.)

**Program 10-9. Source listing for a program that uses a message table built by means of the DEFM operations.**

```

00100          ORC 4A00H          ;SET THE ORIGIN
00110          DEFM 'HELLO'      ;DEFINE MESS1
00120          DEFB 0           ;SET END MARKER
00130          DEFM 'NOW YOU SEE IT' ;DEFINE MESS2
00140          DEFB 0           ;SET END MARKER
00150          DEFM 'NOW YOU DON ' ;DEF PART OF MESS3
00160          DEFB 39          ;DEF APOSTROPHE
00170          DEFB 'I'         ;DEF 'I'
00180          DEFB 0           ;SET END MARKER
00190 SKIP1    DEFS 8           ;LEAVE SOME MEMORY
00200 CLS      EQU 01C9H        ;DEFINE CLS
00210 TOP      EQU 3C00H        ;DEF START OF VIDEO
00220 ;TIME DELAY SUBROUTINE
00230 TIME      LD C,OFFH        ;SET LSB OF TIME
00240 DECC      DEC C           ;COUNT LSB OF TIME
00250          JR NZ,DECC        ;IF NOT DONE,COUNT MORE
00260          DEC B             ;COUNT MSB OF TIME
00270          JR NZ,TIME        ;IF NOT DONE,COUNT MORE
00280          RET              ;ELSE RETURN
00290 ;MESSAGE SUBROUTINE
00300 CMESS      LD A,(DE)        ;FETCH CHARACTER
00310          CP 0              ;IS IT NULL?
00320          RET Z              ;IF SO,RETURN
00330          LD (HL),A          ;ELSE PRINT CHARACTER
00340          INC DE             ;GO TO NEXT CHARACTER
00350          INC HL             ;NEXT VIDEO LOCATION
00360          JR CMESS          ;AND GET NEXT CHARACTER
00370          DEFS 8            ;SKIP SOME MEMORY
00380 ;BECINNINC OF MAINLINE PROGRAM
00390 BLINK      CALL CLS        ;CLEAR THE SCREEN
00400          LD HL,TOP          ;SET VIDEO POINTER
00410          LD DE,4A0DH        ;POINT TO MESS1
00420          CALL CMESS         ;PRINT IT
00430          LD A,OFFH         ;SET UP LONG DELAY
00440 TACN      LD B,OFFH
00450          CALL TIME          ;DO DELAY
00460          DEC A              ;COUNT A
00470          JR NZ,TAGN         ;IF NOT DONE,DO AGAIN
00480 TON      CALL CLS          ;CLEAR THE SCREEN
00490          LD HL,TOP          ;SET VIDEO POINTER
00500          LD (HL),191        ;TURN ON LIGHT
00510          INC HL             ;SPACE
00520          INC HL             ;SPACE
00530          LD DE,4A06H        ;POINT TO MESS2
00540          CALL CMESS         ;PRINT IT
00550          LD B,ODFH          ;SET SHORT DELAY
00560          CALL TIME          ;DO DELAY
00570          CALL CLS          ;CLEAR THE SCREEN
00580          LD HL,TOP          ;SET VIDEO POINTER
00590          INC HL             ;SPACE
00600          INC HL             ;SPACE
00610          LD DE,4A15H        ;POINT TO MESS3
00620          CALL CMESS         ;PRINT IT
00630          LD B,ODFH          ;SET SHORT DELAY
00640          CALL TIME          ;DO DELAY
00650          JR TON            ;AND REPEAT
00660          END BLINK

```

**Program 10-10. Assembled version of the message-printing program listing  
in Program 10-9.**

```

4A00          00100      ORG      4A00H      ;SET ORIGIN
4A00 48        00110      DEFM      'HELLO' ;DEFINE MESS1
4A01 45
4A02 4C
4A03 4C
4A04 4F
4A05 00        00120      DEFB      0          ;SET END MARKER
4A06 4E        00130      DEFM      'NOW YOU SEE IT' ;DEFINE MESS2
4A07 4F
4A08 57
4A09 20
4A0A 59
4A0B 4F
4A0C 55
4A0D 20
4A0E 53
4A0F 45
4A10 45
4A11 20
4A12 49
4A13 54
4A14 00        00140      DEFB      0          ;SET END MARKER
4A15 4E        00150      DEFM      'NOW YOU DON' ;DEF PART OF MESS3
4A16 4F
4A17 57
4A18 20
4A19 59
4A1A 4F
4A1B 55
4A1C 20
4A1D 44
4A1E 4F
4A1F 4E
4A2D 27        00160      DEFB      39          ;DEF APOSTROPHE
4A21 54        00170      DEFB      'T'          ;DEF 'T'
4A22 00        00180      DEFB      0          ;SET END MARKER
0008          00190      SKIP1 DEFS      8          ;LEAVE SOME MEMORY
01C9          00200      CLS      EQU      01C9H      ;DEFINE CLS
3C00          00210      TOP      EQU      3C00H      ;DEF START OF VIDEO
          00220      ;TIME DELAY SUBROUTINE
4A2B 0EFF      00230      TIME      LD C,OFFH      ;SET LSB OF TIME
4A2D 0D        00240      DECC      DEC C          ;COUNT LSB OF TIME
4A2E 20FD      00250      JR NZ,DEC      ;IF NOT DONE,COUNT MORE
4A3D 05        00260      DEC B          ;COUNT MSB OF TIME
4A31 20F8      00270      JR NZ,TIME      ;IF NOT DONE,COUNT MORE
4A33 C9        00280      RET          ;ELSE RETURN
          00290      ;MESSAGE SUBROUTINE
4A34 1A        00300      CMES      LD A,(DE)      ;FETCH CHARACTER
4A35 FE00      00310      CP 0          ;IS IT NULL?
4A37 CB        00320      RET Z          ;IF SO,RETURN
4A38 77        00330      LD (HL),A      ;ELSE PRINT CHARACTER
4A39 13        00340      INC DE          ;GOT TO NEXT CHARACTER
4A3A 23        00350      INC HL          ;NEXT VIDEO LOCATION
4A3B 1BF7      00360      JR CMES      ;AND GET NEXT CHARACTER
0008          00370      DEFS B          ;SKIP SOME MEMORY
          00380      BEGINNING OF MAINLINE PROGRAM
4A45 CDC901    00390      BLINK      CALL CLS      ;CLEAR THE SCREEN
4A4B 21003C    00400      LD HL,TOP      ;SET VIDEO POINTER
4A4B 11D04A    00410      LD DE,4A00H      ;POINT TO MESS1
4A4E CD344A    00420      CALL CMES      ;PRINT IT
4A51 3E0F      00430      LD A,OFFH      ;SET UP LONG DELAY
4A53 06FF      00440      TAGN      LD B,OFFH
4A55 CD2B4A    00450      CALL TIME      ;DO DELAY
4A5B 3D        00460      DEC A          ;COUNT A
4A59 20FB      00470      JR NZ,TAGN      ;IF NOT DONE,DO AGAIN
4A5B CDC901    00480      TON      CALL CLS      ;CLEAR THE SCREEN
4A5E 21003C    00490      LD HL,TOP      ;SET VIDEO POINTER
4A61 36BF      00500      LD (HL),191      ;TURN ON LIGHT
4A63 23        00510      INC HL          ;SPACE

```

```

4A64 23 00520 INC HL ;SPACE
4A65 11064A 00530 LD DE,4A06H ;PD INT TO MESS 2
4A68 CD344A 00540 CALL CMESS ;PRINT IT
4A6B 06DF 00550 LD B,0DFH ;SET SHDRT DELAY
4A6D CD2B4A 00560 CALL TIME ;DO DELAY
4A7D CDC901 00570 CALL CLS ;CLEAR THE SCREEN
4A73 21003C 00580 LD HL,TDP ;SET VIDEO PD INTER
4A76 23 00590 INC HL ;SPACE
4A77 23 00600 INC HL ;SPACE
4A78 11154A 00610 LD DE,4A15H ;PD INT TO MESS 3
4A7B CD344A 00620 CALL CMESS ;PRINT IT
4A7E 06DF 00630 LD B,0DFH ;SET SHDRT DELAY
4A80 CD2B4A 00640 CALL TIME ;DO DELAY
4A83 18D6 00650 JR ION ;AND REPEAT
4A45 00660 END BLINK
00000 TOTAL ERRORS

TDN 4A5B
TAGN 4A53
BLINK 4A45
CMESS 4A34
OECC 4A2D
TIME 4A2B
TDP 3C00
CLS 01C9
SKIP1 4A23

```

The “message table” concludes at line 190 by leaving 8 bytes of unused memory—perhaps for future expansion of that table.

The remainder of the program runs a routine that calls these string messages. Upon executing it, the screen clears (line 390) and HELLO appears in the upper left-hand corner of the screen. That message remains on the screen for about 30 seconds.

When this HELLO phase is done, the program shows a rectangle of light and the NOW YOU SEE IT message. That one remains in place for about 1 second. Then the rectangle of light disappears and the message changes to NOW YOU DON’T. And that message remains on the screen for about 1 second.

After that, the program alternates between the two latter messages—at 1-second intervals—until you interrupt it by working the RESET push button.

The program uses two subroutines: TIME (beginning at line 230) does the time-delay operations, and CMESS (beginning at line 300) calls the messages from the message table and prints them onto the screen. The main program—the one that controls the overall operations and calls the subroutines—begins at BLINK (line 390) and runs through the end of the listing.

The program thus begins with the message table, followed by the two subroutines and, finally, the mainline program. The entry point is at BLINK in line 390. How does the computer know it is to begin running operations from that point? The END BLINK operation in line 660 defines that entry point.

So when this program is loaded under the SYSTEM command, it begins actual operation from the BLINK label. That first line

clears the screen, then the following lines set the video pointer—the HL register pair—to the TOP, video address 3C00H. Line 410 sets the DE register pair to the beginning address of the first message, HELLO, and line 420 calls the CMESS subroutine to print the message onto the screen.

CMESS, beginning at line 300, fetches the first character in the message, checks to see whether or not it is the end-marking null character, and returns to the mainline program if, indeed, it sees that zero character. But if there is more to the message, line 330 prints the character, and the next two lines increment both the message pointer (the DE pair) and the video pointer (the HL pair). Line 360 then calls for repeating the message-printing operation.

In short, subroutine CMESS continues printing characters onto the screen until it finds the end marker—the null character. On returning to the mainline program the next few lines set up and call the time delay routine, TIME.

The same general procedure applies to printing the NOW YOU SEE IT and NOW YOU DON'T messages. The trick, as far as the programmer is concerned, is figuring out where the messages begin in the message table.

The HELLO message certainly begins at 4A00H, because that is where the ORG operation begins loading. But where does NOW YOU SEE IT begin? You can count the characters in the message, add an extra count for the end marker, and figure the address of the next message. This is tedious and offers a fine opportunity to make a serious mistake.

A better way to determine the starting addresses of messages in a message table is to do an assembly operation immediately after the table has been entered into the assembler. After typing in lines 100 through 180, for example, do an A/WE. You will get a NO END STATEMENT message, but that isn't important. What is important is the way the message table is assembled. See the assembled table in lines 100 through 180 in Program 10-10.

In this assembled listing, you can see that HELLO is assembled as hexadecimal ASCII codes in addresses 4A00H through 4A04H. The null character then appears in address 4A50H, and NOW YOU SEE IT begins at 4A06H. The latter address is the starting point for printing out that particular message.

Then at address 4A15H, you find the beginning of the NOW YOU DON'T message.

So the entry points for printing the messages are thus identified for you. It is therefore a good idea to enter such a table first and assemble it before writing the portions of the program that must call the message strings. This way, you know the exact starting addresses for all the messages.

## CHAPTER 11

# Putting It All Together

Only the simplest BASIC and machine-language programs are developed as a one-shot operation. All but the most modest programs ought to be built in a piecemeal fashion, with the loading and testing of critical sections before integrating them into an increasingly complex and useful master program.

This chapter demonstrates some approaches to composing relatively complicated programs. Although the examples, themselves, hardly reach staggering proportions, they illustrate the overall spirit of devising programs of any size.

The whole TRS-80 ROM system is built around a lot of general-purpose routines. A number of examples cited earlier in this book take advantage of some of those routines.

In many instances, however, a TRS-80 ROM routine cannot be run until some important parameters are loaded into specified RAM locations. As the ROM routine is executed, it calls those parameters as they are needed to transform a general set of operations into a very specific one. That is the approach featured in this chapter.

The examples in this chapter are all oriented toward video graphics, but the general ideas apply equally well to any sort of i/o medium.

### **BUILDING A GENERAL-PURPOSE FILL ROUTINE**

One of the most useful graphics routines is one that draws blocks of characters on the screen. It should allow the programmer to specify the position and dimensions of the graphic as well as the character type it is built from. Once those parameters are entered

into some well-defined RAM locations, the general-purpose FILL routine uses them to draw the graphic. The routine can draw any graphic—the RAM parameters determine its nature.

So when you want to draw some figures on the screen, first build a RAM table of all the parameters and then call the FILL routine.

Program 11-1 is an assembled version of a FILL routine. It includes the essential elements of a general-purpose routine, including some goofproofing operations that prevent the system from going crazy in the event of a programming error.

The program allocates a 5-byte RAM table for the critical parameters:

4A00	CHARACTER TYPE
4A01	LSB OF STARTING ADDRESS
4A02	MSB OF STARTING ADDRESS
4A03	NUMBER OF CHARACTERS PER LINE
4A04	NUMBER OF LINES

CHARACTER TYPE determines the alphanumeric or graphic character to be used for drawing the figure on the screen. It can be any hexadecimal number between 20H and FFH. If no specific CHARACTER TYPE is designated by the programmer, the routine automatically sets up 191 (decimal). See line 120 in the listing.

The FILL routine builds the graphic by scanning one character space at a time, from left to right, across the screen. When one line is done, it begins the next line directly below the starting point of the previous one. This scheme means it is necessary to designate a starting point, a line length, and the number of lines.

The START ADDRESS is loaded into the parameter table at addresses 4A01 and 4A02. It can be any 2-byte hexadecimal number but, of course, it ought to be within the video memory space—3C00H through 3FFFH. Since any attempt to “draw” a figure outside the video memory space can cause some serious blowups of existing programs, you will find that the FILL routine includes a goofproofing feature to prevent such a disaster. But more about that later.

As shown in line 130, the START ADDRESS defaults to 3C00H. Unless the programmer designates a different one, the routine begins drawing the graphic in the upper left-hand corner of the screen.

With the CHARACTER TYPE and START ADDRESS thus specified, all that remains in the parameter table is NO. OF LINES and CHAR PER LINE. These are both 1-byte hexadecimal numbers located at addresses 4A03H and 4A04H, respectively. If no parameters are entered by the programmer, lines 140 and 150 in the listing show that the default NO. OF LINES is 2, and CHAR PER LINE is 4.

# Program 11-1. Assembled listing for FILL2.

```

00100 ;FILL ROUTINE -- VER 2: FILENAME FILL2
4A00      00110      ORG      4A00H
4A00      00120      DEFB     191      ;CHARACTER TYPE
4A01 003C      00130      DEFW     3C00H ;START ADDRESS
4A02      02      00140      DEFB     2      ;NO. OF LINES
4A03      04      00150      DEFB     4      ;CHAR PER LINE
4A05      46      00160      DEFM     'FILL OFLOW' ;OVERFLOW MESSAGE
4A06      49
4A07      4C
4A08      4C
4A09      20
4A0A      4F
4A0B      46
4A0C      4C
4A0D      4F
4A0E      57
4A0F      00      00170      DEFB     0      ;ALWAYS 'O'
4A10      2A014A      00180 FILL LD HL,(4A01H) ;FETCH START ADDR
4A13      C03B4A      00190      CALL OFLO ;VALID POINTEH?
4A16      3A034A      00200      LD A,(4A03H) ;FETCH NO. OF LINES
4A19      47      00210      LD B,A ;NO. OF LINES TO B
4A1A      3A044A      00220 LINE LD A,(4A04H) ;FETCH CHAR/LINE
4A1D      4F      00230      LD C,A ;CHAR/LINE TO C
4A1E      3A004A      00240      LD A,(4A00H) ;FETCH CHAR TYPE
4A21      54      00250      LD D,H ;SET LINE START
4A22      5D      00260      LD E,L
4A23      77      00270 PLOT LD (HL),A ;PLOT CHARACTER
4A24      0D      00280      DEC C ;DECREMENT CHAR COUNT
4A25      CA2F4A      00290      JP Z,NLINE ;IF DONE, DO NEXT LINE
4A28      23      00300      INC HL ;ELSE INCREMENT POINTEH
4A29      CD3B4A      00310      CALL OFLO ;CHECK FOR OVERFLOW
4A2C      18F5      00320      JR PLOT ;AND PLOT AGAIN
0001      00330      DEFS 1
4A2F      05      00340 NLINE DEC B ;DECREMENT LINE COUNT
4A30      C8      00350      RET Z ;RETURN IF DONE
4A31      214000      00360      LD HL,64 ;SET LINEFEED
4A34      19      00370      ADD HL,DE ;START OF NEW LINE
4A35      CD3B4A      00380      CALL OFLO ;CHECK FOR OVERFLOW
4A38      18E0      00390      JR LINE ;PLOT NEW LINE
0001      00400      DEFS 1
4A38      3E3F      00410 OFLO LD A,3FH ;SET MAX MSB OF VIDEO
4A3D      BC      00420      CP H ;LESS THAN MSB OF VIDEO?
4A3E      3808      00430      JR C,FLOMES ;IF SO,PRINT MESSAGE
4A40      7C      00440      LD A,H ;GET MSB OF VIDEO
4A41      FE3C      00450      CP 3CH ;UNDERFLOW?
4A43      3806      00460      JR C,FLOMES ;IF SO,PRINT OFLO
4A45      3A004A      00470      LD A,(4A00H) ;ELSE RESTORE CHAR TYPE
4A48      C9      00480      RET ;AND RETURN
0002      00490      DEFS 2
4A4B      21003C      00500 FLOMES LD HL,3C00H ;ELSE SET VIDEO POINTER
4A4E      11054A      00510      LD DE,4A05H ;POINT TO MESSAGE
4A51      1A      00520 MESCHR LD A,(DE) ;GET MESSAGE CHAR
4A52      FE00      00530      CP 0 ;DONE?
4A54      2805      00540      JR Z,LOOP ;IF SO,LOOP TO SELF
4A56      77      00550      LD (HL),A ;PRINT CHAR
4A57      23      00560      INC HL ;ELSE NEXT MESS CHAR
4A58      13      00570      INC DE
4A59      18F6      00580      JR MESCHR
4A5B      18FE      00590 LOOP JR LOOP ;LOOP TO SELF
4A10      00600      END      FILL
00000 TOTAL ERRORS

```



LOOP	4A5B
MESCHR	4A51
FLOMES	4A4B
NLINE	4A2F
PLOT	4A23
LINE	4A1A
OFLO	4A3B
FILL	4A10

The point of this part of the discussion is to show the following principles in action:

1. Most general-purpose routines require a table of working parameters that can be called as the routine is executed.
2. The items in the parameter table must have well-defined and documented addresses (a programmer must know the address and purpose of each item in order to preset them).
3. The program should be written so that the items in the parameter table have valid values from the outset. These determine the default parameters, and leaving them undefined can cause serious problems.
4. Each item in the parameter table should be documented in terms of allowable values. Either that, or the routine should be written so that invalid parameters will not cause a total blowup of the program.

Before running the FILL routine, then, the programmer should preset the items in the parameter table. They can, for instance, be preset by means of POKE statements under BASIC, or they can be set directly under T-BUG. When the FILL routine is finally used as part of a larger graphics program, however, the parameters will be set by means of some LD instructions to the appropriate addresses.

Line 160 sets up a FILL OFLOW message that is called whenever the FILL routine attempts to plot a character outside the video memory range. As described later, that message appears in the upper left corner of the screen whenever such an overflow (or underflow) occurs.

The entry point of the FILL routine is at line 180, address 4A01H. That address must be part of the routine's documentation; any programmer must know where the routine begins its execution.

That first working instruction loads the START ADDRESS into the HL register pair. It does that by fetching the 2-byte address beginning at 4A01H. The next instruction checks that parameter to see whether or not it is a valid one—somewhere within the video

memory space. This is done by calling a subroutine labeled OFLO (see line 410).

Assuming the START ADDRESS is valid, the program resumes operation at line 200. The instruction in that line fetches the NO. OF LINES parameter from the parameter table, and the next instruction loads it to register B. Register B then keeps track of the number of lines yet to be drawn.

Line 220 marks the beginning of the line-drawing part of the routine. It begins by fetching the CHAR PER LINE parameter from the parameter table, and then loading it into register C. Then line 240 fetches the CHAR TYPE to register A.

The current-line starting address is saved in the DE register pair by the instructions in lines 250 and 260. The routine must “remember” the starting address of each line so that it can set the start address of the next one later on.

Finally, the plotting operation begins at line 270. The CHAR TYPE is loaded to the screen at an address determined by the content of the HL register pair. Then CHAR/LINE is decremented and checked for its end. If it so happens that register C is not decremented to zero, operations pick up at line 300.

The INC HL in line 300 sets the video pointer to the next character space. CALL OFLO checks to see whether or not the operation will run the pointer out of video memory; if not, operations loop back to PLOT, and another character is loaded to the screen.

Operations continue looping in the PLOT sequence until one of two things happen: either register C decrements to zero, indicating a line is done, or the HL pointer increments out of video memory space. In the latter case the system prints FILL OFLOW and goes into a harmless loop. In the former case—the more desirable one—line 290 jumps operations to NLINE.

NLINE begins a series of instructions that reset the drawing operation to the beginning of the next line to be drawn. The NO. OF LINES register, register B, is first decremented (line 340 in the listing). If B is not yet zero, 64 (decimal) is added to the starting address of the previous line. CALL OFLO in line 380 checks for a possible overflow out of video memory, and then line 390 returns operations to the LINE routine.

Assuming that no overflows occur, the routine runs until all lines have been drawn: until register B is decremented to zero and the RET Z in line 350 is satisfied. At that time, control returns to the program that called FILL in the first place.

The OFLO routine, beginning at line 410, checks the current value of the video pointer (the HL register pair) for an address that is outside the video memory space. If that fault condition occurs, operations pick up at FLOMES in line 500.

FLOMES and MESCHR print the FILL OFLOW message stored in 4A05H through 4A0FH, using methods already described in some earlier examples. (See Chapter 10.) And once the overflow message is printed, LOOP in line 590 causes the program to “buzz” on that one instruction until the programmer operates the RESET push button. Having the program latch up in this safe fashion is far better than letting it wander off into undefined portions of memory, thus running the risk of losing a lot of programming.

If you want to work through the examples in this chapter on a first-hand basis, load the listing in Program 11-1 under EDTASM, check for errors, and then save the object code with the file name FILL2. The source version can then be saved on tape as SFILL2.

Chart 11-1 illustrates the sort of documentation required for general-purpose routines and FILL2 specifically. Such information is just as valuable as the recorded program, itself.

**Chart 11-1. Overall Documentation for the FILL2 Routine**

<b>Memory Location:</b>	4A00H-4A5BH (86 bytes)
<b>Entry Point:</b>	4A10H
<b>Parameter Table:</b>	
4A00H	CHARACTER TYPE (1 byte)
4A01H	START ADDRESS (2 bytes)
4A03H	NUMBER OF LINES (1 byte)
4A04H	NUMBER OF CHARACTERS PER LINE (1 byte)

Aside from the use of a parameter table, there is little involved in the construction of this FILL routine that hasn't been discussed to some extent in earlier chapters.

The only “problem” with FILL, as it is shown in Program 11-1, is that it cannot be executed alone. The RET Z instruction in line 350 marks the end of the routine, and it spells out a return to a calling program. But what calling program? There is none. You have to write a calling program in order to use a general-purpose routine of this sort.

A calling routine for FILL is shown in Program 11-2. The program is so short that you might be ahead of the game by loading it under T-BUG—it is intended to work in conjunction with T-BUG anyway. So you could do this:

1. Load FILL2 under system.
2. Respond to \*? at the end of the loading of FILL2 by operating the RESET push button.
3. Load T-BUG under system.
4. Respond to \*? by entering a slash (/). That will bring up T-BUG.

### Program 11-2. A test routine for FILL2.

```

                                00100 ;FILL TEST ROUTINE -- VER 1 -- FILENAME FLTST1
                                00110 ORG      7000H
7000 CDC901 00120 FLTST1 CALL 01C9H ;CLEAR THE SCREEN
7003 CD104A 00130 CALL 4A10H ;CALL FILL2
7006 3A4038 00140 KB LD A,(3840H) ;LOOK AT KEYBOARD
7009 FE01 00150 CP 1 ;IS IT ENTER?
700B C20670 00160 JP NZ,KB ;IF NOT,LOOK AGAIN
700E C3A043 00170 JP 17312 ;ELSE JUMP TO T-BUG
7000 00180 END FLTST1
00000 TOTAL ERRORS
```

Now FILL2 and T-BUG are both resident in the system. Load the object code in Program 11-2, using the “M” function of T-BUG.

Next, try out the default parameters with FILL2 by doing a J 7000 from T-BUG. Referring to Program 11-2, this will clear the screen, execute the FILL2 routine, and then hold the graphic on the screen until you strike the ENTER key (see lines 140 through 160). On striking the ENTER key, line 170 does a jump back to T-BUG. From there, you can make any corrections in FILL2; and if FILL2 did, indeed, draw a white rectangle in the upper-left corner of the screen, you are in a position to use T-BUG for setting some of FILL2’s parameters to suit your own specifications.

While running from T-BUG, you can alter the parameter table according to the description in Chart 11-1. And when you’ve set up some of your own parameters, do another J 7000 to bring up the FLTST1 routine. It, again, will do the FILL operation you specified, and it will hold the figure on the screen until you strike the ENTER key to get back to T-BUG.

This notion of using FILL2, T-BUG, and the FLTST1 testing routine together leads to a powerful programming aid. It lets you tinker with the parameters to come up with any sort of graphic you want. The present scheme allows only one graphic to be displayed at any given time, but you can keep track of the parameters for any number of graphics. As demonstrated in the next section, they can be all worked together into a single program.

The testing routine, incidentally, can also be written under EDTASM. Saving its object code as FLTST1 makes it available for future use. Assuming, then, that you have FLTST1 available on tape, you can set up the test scheme as follows:

1. Load FILL2 under SYSTEM.
2. Respond to \*? at the end of loading by operating the RESET push button.
3. Load FLTST1 under SYSTEM.
4. Respond to \*? at the end of its loading by operating the RESET push button.

5. Load T-BUG under SYSTEM.
6. Respond to \*? at the end of its loading by entering a slash (/).

Now you can run the tests and determine FILL2 parameters as described earlier.

## APPLYING AND REFINING FILL2

Table 11-1 shows the FILL parameters for drawing the four segments of a rectangular border figure. They were each generated by means of the technique just described—using Program 11-2 in conjunction with FILL2 and T-BUG. Now the trick is to work out yet another program that loads these parameters into the FILL2 parameter table and to draw them onto the screen. The idea is to run FILL2 four times in succession, loading the parameters for each of the rectangle elements.

**Table 11-1. Memory Map for the FILL Parameter Table**

Graphic	Parameter Addresses				
	4A00	4A01	4A02	4A03	4A04
Top line	83	00	3C	01	40
Bottom line	B0	C0	3F	01	40
Left line	BF	00	3C	10	01
Right line	BF	3F	3C	10	01

At this point the most straightforward way to build a rectangle-drawing program is under T-BUG. The program could, of course, be assembled under EDTASM, but T-BUG is already resident and the program is rather short. So why not try it with the tools already at hand? If it works, then it can be documented and assembled under EDTASM, perhaps with some refinements to FILL2.

A hand-assembled version of the rectangle-drawing program is shown as Program 11-3. With FILL2 loaded into the system, enter this program under T-BUG.

The program is divided into three basic sections: SET UP FILL ROUTINE, FIGURE TABLE, and DRAWING ROUTINE. FIGURE TABLE contains the four sets of parameters for the rectangle figure. Those parameters are to be loaded into FILL2's parameter table.

Referring to the DRAWING ROUTINE, it begins by clearing the screen, then pointing to the address of the first byte of data to be transferred from the FIGURE TABLE to FILL2's parameter table. The third line in the DRAWING ROUTINE then calls SET UP FILL ROUTINE.

**Program 11-3. Assembled listing for a rectangle-drawing routine.**

```

KB      7006
FLTST1 7000
;SET UP FILL ROUTINE
4A60 11 00 4A      LD DE,4A00H      ;BEGINNING OF DESTINATION
4A63 01 05 00      LD BC,5H        ;NO. OF BYTES TO LOAD
4A66 ED B0          LDIR            ;LOAD FILL PARAMETERS
4A68 CD 10 4A      CALL FILL2       ;DO FILL
4A6B C9            RET              ;RETURN TO CALLING PROGRAM

;FIGURE TABLE
4A70 83 00 3C 01 40
4A75 B0 C0 3F 01 40
4A7A BF 00 3C 10 01
4A7F BF 3F 3C 10 01

;DRAWING ROUTINE
7100 CD C9 01      CALL 01C9H      ;CLEAR THE SCREEN
7103 21 70 4A      LD HL,4A70H      ;BEGINNING OF SOURCE
7106 CD 60 4A      CALL 4A60H      ;DRAW TOP LINE
7109 21 75 4A      LD HL,4A75H      ;SET UP BOTTOM LINES
710C CD 60 4A      CALL 4A60H      ;DRAW BOTTOM LINE
710F 21 7A 4A      LD HL,4A7AH      ;SET UP LEFT LINE
7112 CD 60 4A      CALL 4A60H      ;DRAW LEFT LINE
7115 21 7F 4A      LD HL,4A7FH      ;SET UP RIGHT LINE
7118 CD 60 4A      CALL 4A60H      ;DRAW RIGHT LINE
711B 18 FE          LOOP          JR LOOP      ;LOOP TO SELF

```

SET UP FILL ROUTINE first points to the beginning of FILL2's parameter table, loads the number of bytes to be transferred, and then uses an LDIR instruction to make the transfer. By the time the program reaches the CALL FILL2 instruction, the parameters for the top line of the rectangle reside in FILL2's parameter table. Doing the CALL FILL2 then causes the system to draw that top line.

On returning to the DRAWING ROUTINE the system picks up the starting address for the second sequence of parameters—those required for drawing the bottom portion of the rectangle. And then it does the SET UP FILL ROUTINE again.

The scheme continues cycling in this fashion until the drawing is done. The final instruction in the DRAWING ROUTINE brings things to a halt by looping to itself. It is thus necessary to operate the RESET push button to get out of the loop; and returning to T-BUG is a matter of entering /I7312 under the SYSTEM command.

So enter the instructions listed as Program 11-3. With both T-BUG and FILL2 also resident in the system, execute the program from T-BUG by doing a J 7100.

It should be apparent that any program using FILL2 to draw moderately complex figures should use the SET UP FILL ROUTINE

and a FIGURE TABLE. The DRAWING ROUTINE will be different for every drawing program you devise, as will the FIGURE TABLE. But the SET UP FILL ROUTINE will remain the same, no matter what sort of figure you want to draw.

This leads to the notion that the SET UP FILL ROUTINE ought to be fit right into the FILL routine. It would also be nice if the FIGURE TABLE could be permanently defined at the end of that revised FILL routine. Things are tightened up this way in Program 11-4.

FILL3 is a revised version of the original FILL program. It has the T-BUG-tested SET UP FILL edited into the opening phase (lines 180 through 200) and a set of default FILL parameters defining the beginning of the FIGURE TABLE at the end (lines 640 through 680).

Using FILL3 is thus a matter of building an appropriate FIGURE TABLE, beginning at address 4A70H, and writing a calling program at some higher address location. Chart 11-2 documents the entry point and memory map for FILL3.

**Chart 11-2. Overall Documentation for the FILL3 Routine**

<b>Memory Location:</b>	4A00H-4A64H (101 bytes)
<b>Entry Point:</b>	4A10H
<b>Figure Table:</b>	4A70H-4AFFH (160 bytes)

NOTE: Prior to calling FILL3, the HL register pair must point to the first of 5 bytes of the FILL parameters.

The notion of writing one program, using it and testing it for a while, and then revising the original program is the usual procedure for building up truly useful machine-language programs. The various sections can be written under either EDTASM or T-BUG, depending on how extensive they are; more complex routines are better written under EDTASM. But they are tested and coordinated under T-BUG in order to provide instant feedback between tests and small modifications.

With the various sections thus written and debugged, they can be merged into a single program under EDTASM. In the previous example, FILL3 was generated by first loading FILL2 under EDTASM and then using EDTASM's editing features to add in the new sections. The result is a single program and an object code tape that loads the entire FILL3 routine.

Program 11-5 is an example of a program that calls the FILL3 routine to draw a rectangle on the screen. It begins by filling the FIGURE TABLE with the parameters for the rectangular figure (the same ones used in Program 11-3). The DEFS pseudo-op in

# Program 11-4. Assembled listing for FILL3.

```

00100 ;FILL ROUTINE -- VER 3: FILENAME FILL3
4A00      00110      ORC      4A00H
4A00 BF 00120      DEFB      191      ;CHARACTER TYPE
4A01 003C 00130      OEFW      3C00H      ;START ADDRESS
4A03 02 00140      OEFB      2      ;ND. OF LINES
4A04 04 00150      OEFB      4      ;CHAR PER LINE
4A05 46 00160      DEFM      'FILL OFLOW'      ;OFLDW MESSAGE
4A06 49
4A07 4C
4A08 4C
4A09 20
4A0A 4F
4A0B 46
4A0C 4C
4A0D 4F
4A0E 57
4A0F 00 00170      DEFB      0      ;ALWAYS 'O'
4A10 11U04A 00180 SETUP      LD DE,4A00H      ;POINT TO PARAM TABLE
4A13 010500 00190      LD BC,5H      ;SET ND. BYTES TO LOAD
4A16 EDB0 00200      LDIR      ;XFER FROM FIC. TD PAR. TABLE
      00210 ;DO THE FILL ROUTINE
4A18 2A014A 00220 FILL      LD HL,(4A01H)      ;FETCH START ADDR
4A1B CD434A 00230      CALL OFLO      ;VALID POINTER?
4A1E 3A034A 00240      LO A,(4A03H)      ;FETCH NO. OF LINES
4A21 47 00250      LO B,A      ;NO. DF LINES TO B
4A22 3A044A 00260 LINE      LD A,(4A04H)      ;FETCH CHAR/LINE
4A25 4F 00270      LD C,A      ;CHAR/LINE TD C
4A26 3A004A 00280      LD A,(4A00H)      ;FETCH CHAR TYPE
4A29 54 00290      LD D,H      ;SET LINE START
4A2A 5D 00300      LD E,L
4A2B 77 00310 PLOT      LD (HL),A      ;PLOT CHARACTER
4A2C 0D 00320      DEC C      ;OECREMENT CHAR COUNT
4A2D CA374A 00330      JP Z,NLINE      ;IF DONE, DO NEXT LINE
4A30 23 00340      INC HL      ;ELSE INCREMENT POINTER
4A31 CD434A 00350      CALL DFLO      ;CHECK FOR OVERFLOW
4A34 18F5 00360      JR PLOT      ;AND PLOT AGAIN
0001      00370      OEFS 1
4A37 05 00380 NLINE      DEC B      ;DECREMENT LINE COUNT
4A38 C8 00390      RET Z      ;RETURN IF DONE
4A39 214000 00400      LO HL,64      ;SET LINEFEED
4A3C 19 00410      ADD HL,OE      ;START DF NEW LINE
4A3D CD434A 00420      CALL OFLD      ;CHECK FDR OVERFLOW
4A40 1BE0 00430      JR LINE      ;PLOT NEW LINE
0001      00440      DEFS 1
4A43 3E3F 00450 OFLO      LD A,3FH      ;SET MAX MSB OF VIDE0
4A45 BC 00460      CP H      ;LESS THAN MSB OF VIDEO?
4A46 380B 00470      JR C,FLDMES      ;IF SO,PRINT MESSAGE
4A48 7C 00480      LD A,H      ;CET MSB OF VIDE0
4A49 FE3C 00490      CP 3CH      ;UNDERFLOW?
4A4B 3806 00500      JR C,FLOMES      ;IF SO, PRINT OFLO
4A4D 3A004A 00510      LD A,(4A00H)      ;ELSE RESTORE CHAR TYPE
4A50 C9 00520      RET      ;AND RETURN
0002      00530      DEFS 2
4A53 21003C 00540 FLOMES      LD HL,3C00H      ;ELSE SET VIOEO POINTER
4A56 11054A 00550      LD DE,4A05H      ;POINT TO MESSAGE
4A59 1A 00560 MESCJ      LD A,(DE)      ;CET MESSAGE CHAR
4A5A FE00 00570      CP 0      ;DDNE?
4A5C 2805 00580      JR Z,LOOP      ;IF SO,LODP
4A5E 77 00590      LD (HL),A      ;PRINT CHAR
4A5F 23 00600      INC HL      ;ELSE NEXT MESS CHAR
4A60 13 00610      INC DE
4A61 18F6 00620      JR MESCHR
4A63 18FE 00630 LODP      JR LOOP      ;LOOP TD SELF
000B      00640      OEFS 4A70H-4A64H-1 ;START DF FIC. TABLE

```



4A70	BF	00650	DEFB	191
4A71	003C	00660	DEFW	3C00H
4A73	02	00670	DEFB	2
4A74	04	00680	DEFB	4
4A18		00690	END	FILL
00000	TOTAL ERRORS			
LOOP	4A63			
MESCHR	4A59			
FLOMES	4A53			
NLINE	4A37			
PLOT	4A2B			
LINE	4A22			
OFLO	4A43			
FILL	4A18			
SETUP	4A10			

line 330 then skips over a lot of unused FIGURE TABLE memory, thus beginning the actual program at address 7000H. With this REC1 and FILL3 resident in the system, executing from address 7000H causes the rectangle figure to appear on the screen.

The object codes for FILL3 and REC1 can be saved separately on cassette tape. Getting the program running is then a matter of going through the following procedure:

1. Load FILL3 under the SYSTEM command.
2. Answer the \*? at the end of the loading operation by operating the RESET push button.
3. Load REC1 under the SYSTEM command.
4. Answer the \*? at the end of the loading operation by ENTERING a slash or doing /28672.

## BUILDING A GENERAL-PURPOSE MOVE ROUTINE

Most animated graphics programs can be built around general-purpose FILL and MOVE routines. The FILL routine has just been described, and this section takes up the matter of a MOVE routine. In this case the idea is to move simple figures in any desired fashion on the screen. See the MOV1 routine in Program 11-6.

MOV1 has the same basic elements as the general-purpose FILL routines—a parameter table (4B00H through 4B04H) and a routine that uses those parameters (4B10 through 4B3E).

The parameter table consists of an address indicating the current video point on the screen, another address indicating the next video address, and a 1-byte motion code that indicates the direction of motion. Generally speaking, the program fetches the current video address (CPNT) and the motion code (MCODE) from the param-

### Program 11-5. Assembled listing for REC1.

```

00100 ;RECTANGLE-DRAWING SUBROUTINE -- VER 1:FILENAME REC1
4A70      00110      DRG      4A70H
          00120 ;LOAD  FIGURE  TABLE
4A70 83      00130      DEFB      83H
4A71 00      00140      DEFB      00
4A72 3C      00150      DEFB      3CH
4A73 01      00160      DEFB      1
4A74 40      00170      DEFB      40H
4A75 B0      00180      DEFB      0B0H
4A76 C0      00190      DEFB      0C0H
4A77 3F      00200      DEFB      3FH
4A78 01      00210      DEFB      1
4A79 40      00220      DEFB      40H
4A7A BF      00230      DEFB      0BFH
4A7B 00      00240      DEFB      0
4A7C 3C      00250      DEFB      3CH
4A7D 10      00260      DEFB      10H
4A7E 01      00270      DEFB      1
4A7F BF      00280      DEFB      0BFH
4A80 3F      00290      DEFB      3FH
4A81 3C      00300      DEFB      3CH
4A82 10      00310      DEFB      10H
4A83 01      00320      DEFB      1
257C      00330      DEFS      7000H-4A83H-1
4A10      00340 SETUP  EQU      4A10H
4A70      00350 TDP    EQU      4A70H
4A75      00360 BDT    EQU      4A75H
4A7A      00370 LEFT   EQU      4A7AH
4A7F      00380 RIGHT  EQU      4A7FH
7000 CDC901 00390 REC1  CALL    01C9H      ;CLEAR THE SCREEN
7003 21704A 00400      LD HL,TDP      ;SET UP TOP
7006 CD104A 00410      CALL SETUP      ;DRAW IT
7009 21754A D042D      LD HL,BDT      ;SET UP BOTTOM
700C CD104A 00430      CALL SETUP      ;DRAW IT
700F 217A4A 00440      LD HL,LEFT      ;SET UP LEFT
7012 CD104A 00450      CALL SETUP      ;DRAW IT
7015 217F4A 00460      LD HL,RIGHT     ;SET UP RIGHT
7018 CD104A 00470      CALL SETUP      ;DRAW IT
701B 18FE   00480 LODP  JR LODP      ;LOOP TD SELF
7000      00490      END      REC1
00000 TDTAL ERRDRS

LOOP 701B
REC1 7000
RIGHT 4A7F
LEFT 4A7A
BDT 4A75
TDP 4A70
SETUP 4A10

```

enter table. Then it generates a new video point (NPNT) based on the values of CPNT and MCODE. The NPNT thus developed is loaded to the parameter table before the program returns to the calling, or controlling, routine.

Chart 11-3 documents MOV1. Note the motion-code parameters for stepping the video address from a current to a new point. Each time MOV1 is called and MCODE is not a STOP code, NPNT takes on a value representing a single step in the indicated direction. Using

### Program 11-6. Assembled listing for MOV1.

```

00100 ;MOTION ROUTINE -- VER 1:FILENAME MOV1
4B00 00110 ORG 4B00H
00120 ;MOTION TABLE
4B00 003C 00130 OEFW 3C00H ;CURRENT VIDEO POINT
4B02 003C 00140 OEFW 3C00H ;NEW VIDEO POINT
4B04 00 00150 OEFB 0 ;MOTION COOE
000B 00160 OEFB 4B10H-4B04H-1
4B00 00170 CPNT EQU 4B00H
4B02 00180 NPNT EQU 4B02H
4B04 00190 MCODE EQU 4B04H
4B10 2A004B 00200 MOV1 LO HL,(CPNT) ;FETCH CURRENT VIDEO
4B13 3A044B 00210 LO A,(MCOOE) ;FETCH MOTION COOE
4B16 E603 00220 ANO 3 ;ISOLATE HOR BITS
4B18 FE01 00230 CP 1 ;IS IT RIGHT?
4B1A 2807 00240 JR Z,RT ;IF SO,MOVE RIGHT
4B1C FE02 00250 CP 2 ;IS IT LEFT?
4B1E 2004 00260 JR NZ,UO ;IF NOT, CHECK U/O
4B20 2B 00270 OEC HL ;ELSE MOVE LEFT
4B21 1801 00280 JR UO ;ANO CHECK U/O
4B23 23 00290 RT INC HL ;MOVE RIGHT
4B24 114000 00300 UO LO OE,64 ;SET UP VERT MOVE
4B27 3A044B 00310 LO A,(MCOOE) ;FETCH MOTION COOE
4B2A E6C0 00320 ANO 0COH ;ISOLATE U/O BITS
4B2C FE04 00330 CP 4 ;IS IT UP?
4B2E 2807 00340 JR Z,UP ;IF SO,SET UP
4B30 FE08 00350 CP 8 ;IS IT DOWN?
4B32 2006 00360 JR NZ,OUT ;IF NOT, GET OUT
4B34 19 00370 AOD HL,OE ;ELSE MOVE DOWN
4B35 1803 00380 JR OUT ;ANO GET OUT
4B37 AF 00390 UP XOR A ;ZERO CY FLAG
4B38 E052 00400 SBC HL,OE ;MOVE UP
4B3A 22024B 00410 OUT LO (NPNT),HL ;SAVE NEW VIDEO POINT
4B3D C9 00420 RET ;RETURN
4B10 00430 END MOV1
00000 TOTAL ERRORS

OUT 4B3A
UP 4B37
UO 4B24
RT 4B23
MOV1 4B10
MCOOE 4B04
NPNT 4B02
CPNT 4B00

```

this routine to create the effect of motion on the screen is thus a matter of calling it a number of times in succession.

Like most other kinds of general-purpose routines, MOV1 cannot be executed alone. It must be used in conjunction with a calling program that sets the CPNT and MCODE parameters and uses the NPNT parameter to draw the moving figure into its new place on the screen. See one such program, BNCEL, in Program 11-7.

BNCEL calls MOV1 to create the impression of a white rectangle of light bouncing back and forth across the screen. It begins by initializing the CPNT parameter at 3D00H—a point at the left side of the screen, just a bit above center. The instruction in line 200

### Chart 11-3. Overall Documentation for the MOV1 Routine

<b>Memory Location:</b>	4B00H–4B3E (63 bytes)		
<b>Entry Point:</b>	4B10H		
<b>Parameter Table:</b>			
4B00H	CURRENT VIDEO POINT (CPNT), 2 bytes		
4B02H	NEW VIDEO POINT (NPNT), 2 bytes		
4B04H	MOTION CODE (MCODE), 1 byte		
<b>Motion Codes (MCODE Values):</b>			
00H	STOP	08H	DOWN
01H	RIGHT	09H	DOWN/RIGHT
02H	LEFT	0AH	DOWN/LEFT
03H	STOP	0BH	DOWN
04H	UP	0CH	STOP
05H	UP/RIGHT	0DH	RIGHT
06H	UP/LEFT	0EH	LEFT
07H	UP	0FH	STOP

sets NPNT one character space to the right of CPNT. Then line 210 sets the initial MCODE for right motion. Those first steps, lines 160 through 220, initialize the parameter table for MOV1.

Label TRY marks the beginning of the actual animation sequence. The first set of steps determines whether or not the figure's NPNT value will carry it beyond the left or right extremes of the screen. Note that the decision is based on the value of NPNT that is fetched from the MOV1 parameter table at line 230.

If NPNT is going to put the figure beyond the right side of the screen, as determined by instructions in lines 270 through 280, the routine calls SETL to load the MCODE byte for left motion—to a value of 2. But if the instructions in lines 240 and 250 determine that NPNT is about to carry the figure beyond the left extreme edge of the screen, the scheme calls SETR to adjust the MCODE for right motion. Of course, no adjustments in MCODE are made if the figure is not reaching either extreme.

In any case, the routine finds its way to MOVIT at line 380, and that set of instructions calls MOV1 to set up the next NPNT value.

The figure is first erased from its current position and drawn in its new one by the instructions under label PLOT (lines 410 through 460). In those instructions the current video point (CPNT) is loaded from the parameter table of MOV1 to the BC register pair (line 410). A hexadecimal 20—ASCII code for a space—is then loaded to that point to clear the image from the screen. Line 440 fetches the new video point (NPNT) from MOV1's parameter table, and the remaining lines in PLOT print the TRS-80 graphic BFH (a solid rectangle) at that new point. The figure that appears to be moving on the screen is thus determined by the data byte that

# Program 11-7. Assembled listing for BNCE1.

```

00100 ;BOUNCE ROUTINE -- VER 1: FILENAME BNCE1
7000      00110      ORG      7000H
4B10      00120 MOV1  EQU     4B10H
4B00      00130 CPNT  EQU     4B00H
4B02      00140 NPNT  EQU     4B02H
4B04      00150 MCODE EQU     4B04H
7000 CDC901 00160 BNCE1 CALL 01C9H ;CLEAR THE SCREEN
7003 21003D 00170      LD HL,3D00H ;SET INIIAL CURRENT POINT
7006 22004B 00180      LD (CPNT),HL ;LOAD TO MOV1 PARAMETER TABLE
7009 23      00190      INC HL      ;SET INITIAL NEW POINT
700A 22024B 00200      LD (NPNT),HL ;LOAD TO MOV1 MOTION TABLE
700D 3E01    00210      LD A,1      ;SET UP INITIAL MOTION CODE
700F 32044B 00220      LD (MCODE),A ;LOAD IT TO PARAMETER TABLE
7012 2A024B 00230 TRY   LD HL,(NPNT) ;GET NEW POINT FOR TESTING
7015 7C      00240      LD A,H      ;LOOK AT LEFT POSIIION
7016 FE3D    00250      CP 3DH      ;IS IT AT LEFT EXTREME?
7018 3814    00260      JR C,SETR   ;IF SO,SET UP RIGHT MOTION
701A 3E3F    00270      LD A,3FH    ;SET UP RIGHT EXTREME
701C BD      00280      CP L        ;IS IT THERE?
701D 3808    00290      JR C,SETL   ;IF SO, SET UP LEFT MOTION
701F CD4070 00300      CALL PLOT    ;ELSE PLOT THE NEW POINT
7022 22004B 00310      LD (CPNT),HL ;AND SET AS CURRENT
7025 180C    00320      JR MOVIT    ;AND DO THE NEXT STEP
7027 3E02    00330 SETL  LD A,2      ;SET UP LEFT MOTION
7029 32044B 00340      LD (MCODE),A ;AND LOAD TO PARAMETER TABLE
702C 1805    00350      JR MOVIT    ;AND MOVE
702E 3E01    00360 SETR  LD A,1      ;SET UP RIGHT MOIION
7030 32044B 00370      LD (MCODE),A ;AND LOAD IT TO PARAMETER TABLE
7033 CD104B 00380 MOVIT CALL MOV1   ;DO THE MOVE
7036 18DA    00390      JR TRY      ;AND CHECK IT OUT
0008      00400      DEFS 8        ;LEAVE SOME MEMORY
7040 ED4B004B 00410 PLOT LD BC,(CPNT) ;FETCH CURRENT POINT
7044 3E20    00420      LD A,20H    ;SET UP ERASE
7046 02      00430      LD (BC),A   ;ERASE CURRENT POINT
7047 ED4B004B 00440      LD BC,(NPNT) ;FETCH NEW POINT
704B 3EBF    00450      LD A,0BFH   ;SET UP CHARACTER
704D 02      00460      LD (BC),A   ;PLOT NEW POINT
      00470 ;DO A TIME DELAY
704E 060C    00480      LD B,0CH    ;SET MSB OF DELAY
7050 0EFF    00490 SETC  LD C,0FFH  ;SET LSB OF DELAY
7052 0D      00500 DECC  DEC C      ;COUNTDOWN LSB
7053 20FD    00510      JR NZ,DECC  ;IF NOT DONE, COUNT MORE
7055 05      00520      DEC B      ;COUNTDOWN MSB
7056 20F8    00530      JR NZ,SETC  ;IF NOT DONE, COUNT MORE
7058 C9      00540      RET
7000      00550      END      BNCE1
00000 TOTAL ERRORS

DECC 7052
SETC 7050
MOVIT 7033
PLOT 7040
SETL 7027
SETR 702E
TRY 7012
BNCE1 7000
MCODE 4B04
NPNT 4B02
CPNT 4B00
MOV1 4B10

```

is loaded to the A register in line 450. Change that byte, and the moving character will change its appearance.

The PLOT sequence is called from line 300. Looking over that portion of the routine, you can see that PLOT is called only if the next position on the screen is a valid one—not beyond either extreme edge.

Since machine-language routines execute so rapidly, it is usually necessary to insert a time delay between successive PLOTtings. A time-delay routine thus follows the PLOT routine. Using the values loaded to registers B and C, the rectangle of light requires roughly two seconds to cross the screen in either direction.

BNCEL sets up an endless looping action. The animation thus cycles endlessly, and the only way out of it is by operating the RESET push button.

Thus far in this section we have generated two programs: MOV1 and BNCEL. MOV1 is a general-purpose, motion-generating routine, and BNCEL uses MOV1 to achieve one particular sort of animation. A number of other routines can be written—all calling MOV1—to create complex sets of animated figures.

To run the suggested program, generate the MOV1 and BNCEL object codes separately and under EDTASM. Then use the following procedure to get it going:

1. Load the MOV1 object code under SYSTEM.
2. Answer the \*? at the end of the loading by operating the RESET push button.
3. Load BNCEL under the SYSTEM command.
4. Start the program by answering the \*? with a slash or a /28672.

## **APPENDIX A**

# **Number System Base Conversions**

Just about any computer (certainly the TRS-80) is essentially a binary machine; the Z-80 microprocessor does all its control, arithmetic, and logic operations in the base-2, or binary, number system. And it so happens that the Z-80 works with 8-bit binary numbers—a full byte of 8 bits.

People do not think and work with binary numbers very well, however. Such numbers, being made up of 1s and 0s, are very cumbersome. One alternative to purely binary representations of numbers is hexadecimal numbers. The hexadecimal (base-16) number system looks at binary numbers in groups of four; every group of four binary figures (sometimes called a *nibble*) can be represented by a single hexadecimal character. So, instead of having to work with strings of eight 1s and 0s in binary, it is possible to work with just two hexadecimal characters.

While, indeed, many machine-language programmers can learn to work with hexadecimal numbers with great proficiency, the general population still prefers the usual decimal number system. TRS-80 engineers were aware of that fact, and TRS-80 BASIC is built around the decimal number system exclusively.

As long as one works with TRS-80 BASIC in its most elementary fashion—doing no special addressing or machine-language work—there is no need to be aware of hexadecimal or binary numbers. But hexadecimal numbers become helpful when POKEing and PEEKing in memory, and they become quite necessary when doing extensive machine-language programming.

Then, too, binary numbering becomes important when attempting to write machine-language program for controlling custom circuitry connected to the TRS-80's output connector.

Thus, programmers working deeper and deeper into the TRS-80 system will find themselves having to make conversions between decimal and hexadecimal numbers and, eventually, between binary and hexadecimal numbers. The purpose of this appendix is to make such conversion tasks as simple as possible.

There are many ways to approach the conversions between these three different number systems; but, in this writer's opinion, the following are the most straightforward.

## HEXADECIMAL-TO-DECIMAL CONVERSIONS

In the TRS-80/Z-80 system, data is carried as a 1-byte (two-hexadecimal-character) code, and addresses are carried as a 2-byte (four-hexadecimal-character) code. Table A-1 can be very helpful for translating hexadecimal numbers into their decimal counterparts. This sort of situation often arises when one is writing programs in both BASIC and machine language.



**Table A-1. Hexadecimal to Decimal Conversion**

4		3		2		1	
Hex	Decimal	Hex	Decimal	Hex	Decimal	Hex	Decimal
0	0	0	0	0	0	0	0
1	4096	1	256	1	16	1	1
2	8192	2	512	2	32	2	2
3	12288	3	768	3	48	3	3
4	16384	4	1024	4	64	4	4
5	20480	5	1280	5	80	5	5
6	24576	6	1536	6	96	6	6
7	28672	7	1792	7	112	7	7
8	32768	8	2048	8	128	8	8
9	36864	9	2304	9	144	9	9
A	40960	A	2560	A	160	A	10
B	45056	B	2816	B	176	B	11
C	49152	C	3072	C	192	C	12
D	53248	D	3328	D	208	D	13
E	57344	E	3584	E	224	E	14
F	61440	F	3840	F	240	F	15

The table can be used for converting up to four hexadecimal places to their decimal counterpart. Note that there are four major columns, labeled 1 through 4. These column numbers represent the relative positions of the hexadecimal characters as they are usually written, with the lsb on the left and the msb on the right.

To see how the table works, suppose that you want to convert the hexadecimal number 1A3F into decimal. The first character on the left takes a decimal equivalent shown in column 4—4096. The second character from the left takes the value from column 3—2560. The last two figures get their decimal equivalents from columns 2 and 1—48 and 15, respectively. Then, to get the true decimal value, add those decimal equivalents:  $4096+2560+48+15$ . That comes out to 6719. In other words, 1A3F (hexadecimal) is equal to 6719 (decimal).

If you are converting a two-place hexadecimal number, just use columns 2 and 1. Hexadecimal C3, for instance, is equal to  $192+3$ , or decimal 195.

Table A-1 is adequate for hexadecimal-to-decimal conversions for all the usual sort of work on the TRS-80/Z-80 system.

### DECIMAL-TO-HEXADECIMAL CONVERSIONS

When working back and forth between BASIC and machine language, it is often necessary to convert decimal data and addresses into hexadecimal notation. Table A-1 comes to the rescue here. The procedure is a rather straightforward one, but it involves several steps.

Suppose, for example, you want to convert decimal 65 into its hexadecimal counterpart. First, find the decimal number (in a decimal column) that is equal to, or less than, the desired decimal number—65 is the decimal number, and its closest, and less, value is 64. The 64 is equivalent to a hexadecimal 4 in column 2. Thus the most significant character in the hexadecimal representation is a 4.

Next, subtract that 64 from the number you are working with:  $65 - 64 = 1$ . Now look up the hexadecimal value of the 1 in the next-lower column—the 1 column in this case. The hexadecimal version of that number is 1.

Putting together those two hexadecimal characters, you get a 41. Indeed, decimal 65 translates into hexadecimal 41.

By way of a somewhat more involved example, suppose you have to convert decimal 19314 into hexadecimal notation.

Looking through the columns of DEC numbers, you find that 16384 is the next-lower value; it translates into hexadecimal 4 in column 4. So you are going to end up with a four-character hexadecimal number, with the digit on the left being a 4.

To get the next place value, subtract 16384 from 19314:  $19314 - 16384 = 2930$ . The next-lower decimal value in this case is 2816 from column 3; and that turns up a B as the next hexadecimal character. So far, the number is 4B.

Now do this:  $2930 - 2816 = 114$ . The next-lower decimal value from column 2 is 112, and its hexadecimal counterpart is 7. And to this point, the hexadecimal number is 4B7.

Finally, do this:  $114 - 112 = 2$ . From column 1, decimal 2 is the same as hexadecimal 2; so the final hexadecimal character is 2.

Putting this all together, it turns out that decimal 19314 is the same as hexadecimal 4B72. Fig. A-1 summarizes the operation.

## **DECIMAL ADDRESS TO 2-BYTE DECIMAL FORMAT**

When POKEing addresses as 2-byte numbers into memory, it is necessary to convert the address to be loaded into a 2-byte format. In decimal, this isn't easy, but it is all a part of setting up address locations in decimal-oriented BASIC.

By way of an example, suppose you are to load a 2-byte version of decimal address 1234 into memory addresses 16787 and 16788. That number to be stored, 1234, has to be broken up into two parts: one for each of the places it is to be stored.

Before a decimal number can be divided into a 2-byte version, it must be converted into a hexadecimal form. Using the decimal to hexadecimal conversion described in the previous section, it turns out that 1234 decimal is equal to 04D2 in hexadecimal.

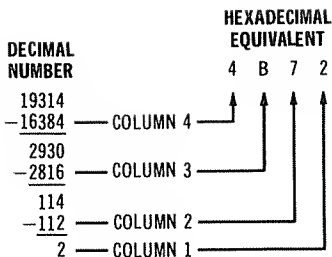


Fig. A-1. Converting 19314 decimal to 4B72 hexadecimal.

19314 (DECIMAL) = 4B72 (HEXADECIMAL)

Next, divide that hexadecimal version into two bytes: the most significant byte (msb) is 04, and the least significant byte is D2. Divided this way, you get the combinations 04 D2.

Finally, convert those two sets of hexadecimal numbers into their decimal equivalents, treating them as two separate hexadecimal numbers. Thus 04 hexadecimal is 4 decimal, and D2 hexadecimal is 210.

The 2-byte version of decimal 1234 is thus 4 and 210, with 4 being the msb and 210 being the lsb.

That takes care of the conversion of an ordinary decimal number into a 2-byte version—also in decimal. This is the main point of the discussion at hand. But the example calls for POKEing these numbers into locations 16787 and 16788—both in decimal.

Now, almost without exception, you will find that the lsb of the 2-byte number must go into the lower-numbered address; so the BASIC operation for satisfying the requirements for the example looks like this:

POKE 16787,210:POKE 16788,4

No, it isn't a simple operation, but it's the price that must be paid for working with a binary/hexadecimal-oriented computer system in a decimal, BASIC-oriented language.

## 2-BYTE DECIMAL TO CONVENTIONAL DECIMAL

Suppose you are disassembling a machine-language program from BASIC. Under that condition, a 2-byte address will appear as a set of two decimal numbers; and if you want to get that pair of numbers into a conventional decimal format, you have to play with the numbers a bit.

Consider a case where 223 turns up as the lsb in decimal, and 104 is the msb. First, convert both sets of numbers into their hexa-

decimal counterparts: 223 decimal = DF hexadecimal, and 104 decimal = 68 hexadecimal.

That means DF is the lsb and 68 is the msb. Putting them together in the conventional order (msb first), the hexadecimal equivalent of the number under consideration is 68DF.

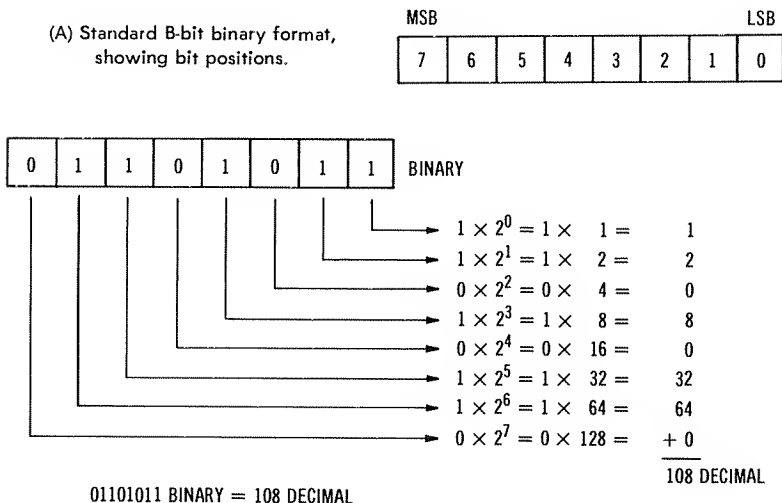
All that remains to be done is to convert that into a full decimal number:  $68DF = 24567 + 2048 + 208 + 15 = 26849$  decimal. That's it—the conventional decimal version of the 2-byte decimal combination 104 223 is 26849. For one reason or another, the disassembled program is saving address 26849.

## BINARY-TO-DECIMAL CONVERSION

In practice, most binary-to-decimal conversions are carried out with 1-byte (or 8-bit) binary numbers, although there are occasions when it is necessary to do the conversion from 2-byte (16-bit) numbers.

Fig. A-2A shows the breakdown of an 8-bit binary number. The positions are labeled 0 through 7, with zero indicating the least significant bit (lsb) position. Each of those eight bit locations will contain either a 1 or a 0.

To appreciate the relevance of this whole thing, suppose you want to POKE a decimal number into some address location, and select a number such that it looks like this in binary form: 01101011.



(B) Converting binary word 01101011 to decimal.

**Fig. A-2. Power-of-2 place values of an 8-bit binary word.**

So you want bit positions 0, 1, 3, 5, and 6 to be a 1, and the rest 0. But you have to POKE a decimal version from BASIC. Here's how to go about determining that decimal version.

First, multiply the 1 or 0 in each bit location times  $2^n$ , where  $n$  is the bit's position value in each case. Then simply add the results. See the example in Fig. A-2B. That particular number in binary is equal to 108 decimal.

The same idea applies to converting 16-bit binary to a decimal equivalent. The place values run from 0 to 15 in that case, and Table A-2 can help you determine those larger powers of 2.

**Table A-2. Powers of 2**

n	$2^n$
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
10	1 024
11	2 048
12	4 096
13	8 192
14	16 384
15	32 768

## **BINARY-TO-HEXADECIMAL CONVERSION**

Converting a binary number into a hexadecimal format is perhaps the simplest of all the conversion operations. All you have to do is group the binary number into sets of 4 bits each, beginning with the lsb, and then find the hexadecimal value for each group. Table A-3 helps with the latter operation.

Suppose the binary number is 10011101. Grouping this number into sets of 4 bits (or nibbles), it looks like this: 1001 1101. The hexadecimal equivalents for each group can be found in Table A-3: 9 D. The hexadecimal version of that 8-bit binary number is thus 9D.

The same procedure works equally well for 16-bit numbers, the only difference being that you end up with four hexadecimal characters instead of just two of them.

Incidentally, there is an algorithm for converting directly from binary to hexadecimal (without using the table), but its complexity far exceeds that of doing the job with Table A-3.

**Table A-3. Binary to Hexadecimal Conversion**

Binary	Hexadecimal
0000	0
0001	1
0010	2
0011	3
0100	4
0101	5
0110	6
0111	7
1000	8
1001	9
1010	A
1011	B
1100	C
1101	D
1110	E
1111	F

### **HEXADECIMAL-TO-BINARY CONVERSION**

Converting a hexadecimal number to binary is a simple matter of applying Table A-3 to convert each hexadecimal character into the appropriate groups of 4 binary bits.

Example: Convert address 403D into binary. According to Table A-3, that hexadecimal number can be represented as

0100 0000 0011 1101

Nothing more, nothing less. That's all there is to it.

### **DECIMAL-TO-BINARY CONVERSION**

There is a nice algorithm (several of them, in fact) for mathematically converting any decimal number into its binary format. But it is simpler in the long run, and probably more accurate as well, to use a two-step procedure.

The general idea is to convert the decimal number into its hexadecimal counterpart as described earlier in this appendix. Then convert the hexadecimal characters into their binary versions as described in the previous section. Besides, it is often handy to have the hexadecimal version available for later use.

Example: Convert 1234 decimal into binary. First, as described earlier, calculate the hexadecimal version of decimal 1234—that comes out to be 04D2. And that hexadecimal number, expressed in binary (from Table A-3) is 0000 0100 1101 0010. Thus 1234 in decimal is equal to 10011010010 in binary. You may retain the leading zeros if you wish.

## **APPENDIX B**

### **Z-80 Instruction Set: Object and Source Codes**

# ADD WITH CARRY INSTRUCTIONS

```

8E          ADC A,(HL)
DD 8E byte  ADC A,(IX+indx)
FD 8E byte  ADC A,(IY+indx)

```

```

8F          ADC A,A
88          ADC A,B
89          ADC A,C
8A          ADC A,D
8B          ADC A,E
8C          ADC A,H
8D          ADC A,L

```

```
CE byte    ADC A,data
```

```

ED 4A      ADC HL,BC
ED 5A      ADC HL,DE
ED 6A      ADC HL,HL
ED 7A      ADC HL,SP

```

# ADD INSTRUCTIONS

```

86          ADD A,(HL)
DD 86 byte  ADD A,(IX+indx)
FD 86 byte  ADD A,(IY+indx)

```

```

87          ADD A,A
80          ADD A,B
81          ADD A,C
82          ADD A,D
83          ADD A,E
84          ADD A,H
85          ADD A,L

```

```
C6 byte    ADD A,data
```

```

09          ADD HL,BC
19          ADD HL,DE
29          ADD HL,HL
39          ADD HL,SP

```

```

DD 09      ADD IX,BC
DD 19      ADD IX,DE
DD 29      ADD IX,IX
DD 39      ADD IX,SP
FD 09      ADD IY,BC
FD 19      ADD IY,DE
FD 29      ADD IY,IY
FD 39      ADD IY,SP

```

# LOC IC-AND INSTRUCTIONS

```

A6          AND (HL)
DD A6 byte  AND (IX+indx)
FD A6 byte  AND (IY+indx)

```

```

A7          AND A
AD          AND B
A1          AND C
A2          AND D
A3          AND E
A4          AND H
A5          AND L

```

```
E6 byte    AND data
```

# BIT-TEST INSTRUCTIONS

```

CB 46      BIT 0,(HL)
DD CB byte 46  BIT 0,(IX+indx)
FD CB byte 46  BIT 0,(IY+indx)

```

```

CB 47      BIT 0,A
CB 40      BIT 0,B
CB 41      BIT 0,C
CB 42      BIT 0,D
CB 43      BIT 0,E
CB 44      BIT 0,H
CB 45      BIT 0,L

```

```

CB 4E      BIT 1,(HL)
DD CB byte 4E  BIT 1,(IX+indx)
FD CB byte 4E  BIT 1,(IY+indx)

```

```

CB 4F      BIT 1,A
CB 48      BIT 1,B
CB 49      BIT 1,C
CB 4A      BIT 1,D
CB 4B      BIT 1,E
CB 4C      BIT 1,H
CB 4D      BIT 1,L

```

```

CB 56      BIT 2,(HL)
DD CB byte 56  BIT 2,(IX+indx)
FD CB byte 56  BIT 2,(IY+indx)

```

```

CB 57      BIT 2,A
CB 5D      BIT 2,B
CB 51      BIT 2,C
CB 52      BIT 2,D
CB 53      BIT 2,E
CB 54      BIT 2,H
CB 55      BIT 2,L
CB 5E      BIT 3,(HL)
DD CB byte 5E  BIT 3,(IX+indx)
FD CB byte 5E  BIT 3,(IY+indx)

```

```

CB 5F      BIT 3,A
CB 58      BIT 3,B
CB 59      BIT 3,C
CB 5A      BIT 3,D
CB 5B      BIT 3,E
CB 5C      BIT 3,H
CB 5D      BIT 3,L

```



CB 66	BIT 4,(HL)
DD CB byte 66	BIT 4,(IX+indx)
FD CB byte 66	BIT 4,(IX+indx)
CB 67	BIT 4,A
CB 60	BIT 4,B
CB 61	BIT 4,C
CB 62	BIT 4,D
CB 63	BIT 4,E
CB 64	BIT 4,H
CB 65	BIT 4,L
CB 6E	BIT 5,(HL)
DD CB byte 6E	BIT 5,(IX+indx)
FD CB byte 6E	BIT 5,(IX+indx)
CB 6F	BIT 5,A
CB 68	BIT 5,B
CB 69	BIT 5,C
CB 6A	BIT 5,D
CB 6B	BIT 5,E
CB 6C	BIT 5,H
CB 6D	BIT 5,L
CB 76	BIT 6,(HL)
DD CB byte 76	BIT 6,(IX+indx)
FD CB byte 76	BIT 6,(IX+indx)
CB 77	BIT 6,A
CB 70	BIT 6,B
CB 71	BIT 6,C
CB 72	BIT 6,D
CB 73	BIT 6,E
CB 74	BIT 6,H
CB 75	BIT 6,L
CB 7E	BIT 7,(HL)
DD CB byte 7E	BIT 7,(IX+indx)
FD CB byte 7E	BIT 7,(IX+indx)
CB 7F	BIT 7,A
CB 78	BIT 7,B
CB 79	BIT 7,C
CB 7A	BIT 7,D
CB 7B	BIT 7,E
CB 7C	BIT 7,H
CB 7D	BIT 7,L

#### CALL INSTRUCTIONS

DC byte byte	CALL C,addr
FC byte byte	CALL M,addr
4C byte byte	CALL NC,addr
C4 byte byte	CALL NZ,addr
F4 byte byte	CALL P,addr
EC byte byte	CALL PE,addr
E4 byte byte	CALL PO,addr
CC byte byte	CALL Z,addr
CD byte byte	CALL addr

#### COMPLEMENT CARRY FLAG

3F	CCF
----	-----

#### COMPARE ACCUMULATOR INSTRUCTIONS

BE	CP (HL)
DD BE byte	CP (IX+indx)
FD BE byte	CP (IX+indx)
BF	CP A
B8	CP B
B9	CP C
BA	CP D
BB	CP E
BC	CP H
BD	CP L

#### COMPARE AND INCREMENT/DECREMENT

ED A9	CPD
ED B9	CPDR
ED A1	CPI
ED B1	CPIR

#### COMPLEMENT ACCUMULATOR

2F	CPL
----	-----

#### DECIMAL ADJUST THE ACCUMULATOR

27	DAA
----	-----

#### DECREMENT INSTRUCTIONS

35	DEC (HL)
DD 35 byte	DEC (IX+indx)
FD 35 byte	DEC (IX+indx)
3D	DEC A
05	DEC B
0D	DEC C
15	DEC D
1D	DEC E
25	DEC H
2D	DEC L
0B	DEC BC
1B	DEC DE
2B	DEC HL
DD 2B	DEC IX
FD 2B	DEC IX
3B	DEC SP



FD 77 byte	LD (IY+indx),A	57	LD D,A
FD 70 byte	LD (IY+indx),B	50	LD D,B
FD 71 byte	LD (IY+indx),C	51	LD D,C
FD 72 byte	LD (IY+indx),D	52	LD D,D
FD 73 byte	LD (IY+indx),E	53	LD D,E
FD 74 byte	LD (IY+indx),H	54	LD D,H
FD 75 byte	LD (IY+indx),L	55	LD D,L
FD 36 byte byte	LD (IY+indx),data	16 byte	LD D,data
32 byte byte	LD (addr),A	ED 5B byte byte	LD DE,(addr)
ED 43 byte byte	LD (addr),BC	11 byte byte	LD DE,data data
ED 53 byte byte	LD (addr),DE	5E	LD E,(HL)
22 byte byte	LD (addr),HL	DD 5E byte	LD E,(IX+indx)
DD 22 byte byte	LD (addr),IX	FD 5E byte	LD E,(IY+indx)
FD 22 byte byte	LD (addr),IY		
ED 73 byte byte	LD (addr),SP	5F	LD E,A
0A	LD A,(BC)	58	LD E,B
1A	LD A,(DE)	59	LD E,c
7E	LD A,(HL)	5A	LD E,D
DD 7E byte	LD A,(IX+indx)	5B	LD E,E
FD 7E byte	LD A,(IY+indx)	5C	LD E,H
3A byte byte	LD A,(addr)	5D	LD E,L
		1E byte	LD E,data
7F	LD A,A		
78	LD A,B	66	LD H,(HL)
79	LD A,C	DD 66 byte	LD H,(IX+indx)
7A	LD A,D	FD 66 byte	LD H,(IY+indx)
7B	LD A,E		
7C	LD A,H	67	LD H,A
7D	LD A,L	60	LD H,B
ED 57	LD A,I	61	LD H,C
ED 5F	LD A,R	63	LD H,D
3E byte	LD A,data	64	LD H,E
		65	LD H,L
46	LD B,(HL)	26 byte	LD H,data
DD 46 byte	LD B,(IX+indx)		
FD 46 byte	LD B,(IY+indx)	2A byte byte	LD HL,(addr)
47	LD 8,A	21 byte byte	LD HL,data data
40	LD B,B		
41	LD B,C	6E	LD L,(HL)
42	LD B,D	DD 6E byte	LD L,(IX+indx)
43	LD B,E	FD 6E byte	LD L,(IY+indx)
44	LD B,H		
45	LD B,L	6F	LD L,A
06 byte	LD B,data	68	LD L,B
		69	LD L,C
ED 4B byte byte	LD BC,(addr)	6A	LD L,D
01 byte byte	LD BC,data data	6B	LD L,E
		6C	LD L,H
4E	LD C,(HL)	6D	LD L,L
DD 4E byte	LD C,(IX+indx)	2E byte	LD L,data
FD 4E byte	LD C,(IY+indx)		
4F	LD C,A	ED 47	LD I,A
48	LD C,B	ED 4F	LD R,A
49	LD C,C	DD 2A byte byte	LD IX,(addr)
4A	LD C,D	FD 2A byte byte	LD IY,(addr)
4B	LD C,E	DD 21 byte byte	LD IX,data data
4C	LD C,H	FD 21 byte byte	LD IY,data data
4D	LD C,L	ED 7B byte byte	LD SP,(addr)
0E byte	LD C,data	F9	LD SP,HL
		DD F9	LD SP,IX
		FD F9	LD SP,IY
		31 byte byte	LD SP,data data
56	LD D,(HL)		
DD 56 byte	LD D,(IX+indx)		
FD 56 byte	LD D,(IY+indx)		

# LOAD WITH INCREMENT/DECREMENT

```
ED A8      LDD
ED B8      LDDR
ED A0      LDI
ED B0      LDIR
```

```
ED BB      OTRD
ED B3      OTIR
ED AB      OUTD
ED A3      OUTI
```

# NO OPERATION

```
00      NOP
```

# NEGATE ACC. (2'S COMPLEMENT)

```
ED 44      NEC
```

# PUSH AND POP INSTRUCTIONS

```
F1      POP AF
C1      POP BC
D1      POP DE
E1      POP HL
DD E1   POP IX
FD E1   POP IY
```

```
F5      PUSH AF
C5      PUSH BC
D5      PUSH DE
E5      PUSH HL
DD 55   PUSH IX
FD 55   PUSH IY
```

# LOGIC-OR INSTRUCTIONS

```
B6      OR (HL)
DD B6 byte OR (IX+indx)
FD B6 byte OR (IY+indx)
```

# RESET BII INSTRUCTIONS

```
B7      OR A
B0      OR B
B1      OR C
B2      OR D
B3      OR E
B4      OR H
B5      OR L
F6 byte OR data
```

```
CB 86      RES (HL)
DD CB byte 86 RES (IX+indx)
FD CB byte 86 RES (IY+indx)
```

```
CB 87      RES 0,A
CB 80      RES 0,B
CB 81      RES 0,C
CB 82      RES 0,D
CB 83      RES 0,E
CB 84      RES 0,H
CB 85      RES 0,L
```

# EXCLUSIVE-OR INSTRUCTIONS

```
AE      XOR (HL)
DD AE byte XOR (IX+indx)
FD AE byte XOR (IY+indx)
```

```
CB 8E      RES 1,(HL)
DD CB byte 8E RES 1,(IX+indx)
FD CB byte 8E RES 1,(IY+indx)
```

```
AF      XOR A
A8      XOR B
A9      XOR C
AA      XOR D
AB      XOR E
AC      XOR H
AD      XOR L
EE byte XOR dataa
```

```
CB 8F      RES 1,A
CB 86      RES 1,B
CB 89      RES 1,C
CB 8A      RES 1,D
CB 8B      RES 1,E
CB 8C      RES 1,H
CB 8D      RES 1,L
```

```
CB 96      RES 2,(HL)
DD CB byte 96 RES 2,(IX+indx)
FD CB byte 96 RES 2,(IY+indx)
```

# OUTPUT TO PORT INSTRUCTIONS

```
ED 79      OUT (C),A
ED 41      OUT (C),B
ED 49      OUT (C),C
ED 51      OUT (C),D
ED 59      OUT (C),E
ED 61      OUT (C),H
ED 69      OUT (C),L
D3 byte OUT (port),A
```

```
CB 97      RES 2,A
CB 90      RES 2,B
CB 91      RES 2,C
CB 92      RES 2,D
CB 93      RES 2,E
CB 94      RES 2,H
CB 95      RES 2,L

CB 9E      RES 3,(HL)
DD CB byte 9E RES 3,(IX+indx)
FD CB byte 9E RES 3,(IY+indx)
```

CB 9F	RES 3,A
CB 98	RES 3,B
CB 99	RES 3,C
CB 9A	RES 3,D
CB 9B	RES 3,E
CB 9C	RES 3,H
CB 9D	RES 3,L

CB A6	RES 4,(HL)
DD CB byte A6	RES 4,(IX+indx)
FD CB byte A6	RES 4,(IY+indx)

CB A7	RES 4,A
CB A0	RES 4,B
CB A1	RES 4,C
CB A2	RES 4,D
CB A3	RES 4,E
CB A4	RES 4,H
CB A5	RES 4,L

CB AE	RES 5,(HL)
DD CB byte AF	RES 5,(IX+indx)
FD CB byte AF	RES 5,(IY+indx)

CB AF	RES 5,A
CB A8	RES 5,B
CB A9	RES 5,C
CB AA	RES 5,D
CB AB	RES 5,E
CB AC	RES 5,H
CB AD	RES 5,L

CB B6	RES 6,(HL)
DD CB byte B6	RES 6,(IX+indx)
FD CB byte B6	RES 6,(IY+indx)

CB B7	RES 6,A
CB B0	RES 6,B
CB B1	RES 6,C
CB B2	RES 6,D
CB B3	RES 6,E
CB B4	RES 6,H
CB B5	RES 6,L

CB BE	RES 7,(HL)
DD CB byte BE	RES 7,(IX+indx)
FD CB byte BE	RES 7,(IY+indx)

CB BF	RES 7,A
CB B8	RES 7,B
CB B9	RES 7,C
CB BA	RES 7,D
CB BB	RES 7,E
CB BC	RES 7,H
CB BD	RES 7,L

#### RETURN INSTRUCTIONS

C9	RET
D8	REI C
F8	RET M
D0	REI NC
C0	REI NZ
F0	REI P
E8	RET PE
E0	RET PO
C8	RET Z

ED 4D	RETI
ED 45	RETN

#### ROTATE THROUGH CARRY INSTRUCTIONS

17	RLA
CB 16	RL (HL)
DD CB byte 16	RL (IX+indx)
FD CB byte 16	RL (IY+indx)

CB 17	RL A
CB 10	RL B
CB 11	RL C
CB 12	RL D
CB 13	RL E
CB 14	RL H
CB 15	RL L

1F	RRA
CB 1E	RR (HL)
DD CB byte 1E	RR (IX+indx)
FD CB byte 1E	RR (IY+indx)

CB 1F	RR A
CB 18	RR B
CB 19	RR C
CB 1A	RR D
CB 1B	RR E
CB 1C	RR H
CB 1D	RR L

#### ROTATE CIRCULAR INSTRUCTIONS

07	RLCA
CB 06	RLC (HL)
DD CB byte 06	RLC (IX+indx)
FD CB byte 06	RLC (IY+indx)

CB 07	RLC A
CB 00	RLC B
CB 01	RLC C
CB 02	RLC D
CB 03	RLC E
CB 04	RLC H
CB 05	RLC L

0F	RRCA
CB 0E	RRC (HL)
DD CB byte 0E	RRC (IX+indx)
FD CB byte 0E	RRC (IY+indx)

CB 0F	RRC A
CB 08	RRC B
CB 09	RRC C
CB 0A	RRC D
CB 0B	RRC E
CB 0C	RRC H
CB 0D	RRC L

ROTATE ACC. AND (HL) INSTRUCTIONS				
-----			CB C1	SET 0,C
			CB C2	SET 0,D
			CB C3	SET 0,E
ED 6F	RLD		CB C4	SET 0,H
ED 67	RRD		CB C5	SET 0,L
			CB CE	SET 1,(HL)
			DD CB byte CE	SET 1,(IX+indx)
			FD CB byte CE	SET 1,(IY+indx)
RESTART INSTRUCTIONS				
-----			CB CF	SET 1,A
			CB CB	SET 1,B
			CB C9	SET 1,C
C7	RST 00H		CB CA	SET 1,D
CF	RST 0BH		CB CB	SET 1,E
D7	RST 10H		CB CC	SET 1,H
DF	RST 18H		CB CD	SET 1,L
E7	RST 20H			
EF	RST 2BH		CB D6	SET 2,(HL)
F7	RST 30H		DD CB byte D6	SET 2,(IX+indx)
FF	RST 3BH		FD CB byte D6	SET 2,(IY+indx)
			CB D7	SET 2,A
			CB D0	SET 2,B
			CB D1	SET 2,C
			CB D2	SET 2,D
			CB D3	SET 2,E
			CB D4	SET 2,H
			CB D5	SET 2,L
SUBTRACT WITH CARRY INSTRUCTIONS				
-----			CB DE	SET 3,(HL)
9E	SBC A,(HL)		DD CB byte DE	SET 3,(IX+indx)
DD 9E byte	SBC A,(IX+indx)		FD CB byte DE	SET 3,(IY+indx)
FD 9E byte	SBC A,(IY+indx)			
9F	SBC A,A		CB DF	SET 3,A
9B	SBC A,B		CB DB	SET 3,B
99	SBC A,C		CB D9	SET 3,C
9A	SBC A,D		CB DA	SET 3,D
9B	SBC A,E		CB DB	SET 3,E
9C	SBC A,H		CB DC	SET 3,H
9D	SBC A,L		CB DD	SET 3,L
DE byte	SBC A,data			
ED 42	SBC HL,BC		CB E6	SET 4,(HL)
ED 52	SBC HL,DE		DD CB byte E6	SET 4,(IX+indx)
ED 62	SBC HL,HL		FD CB byte E6	SET 4,(IY+indx)
ED 72	SBC HL,SP		CB E7	SET 4,A
			CB E0	SET 4,B
			CB E1	SET 4,C
			CB E2	SET 4,D
			CB E3	SET 4,E
			CB E4	SET 4,H
			CB E5	SET 4,L
SET CARRY FLAG				
-----			CB EE	SET 5,(HL)
37	SCF		DD CB byte EE	SET 5,(IX+indx)
			FD CB byte EE	SET 5,(IY+indx)
			CB EF	SET 5,A
			CB E8	SET 5,B
			CB E9	SET 5,C
			CB EA	SET 5,D
			CB EB	SET 5,E
			CB EC	SET 5,H
			CB ED	SET 5,L
SET BIT INSTRUCTIONS				
-----			CB F6	SET 6,(HL)
CB C6	SET 0,(HL)		DD CB byte F6	SET 6,(IX+indx)
DD CB byte C6	SET 0,(IX+indx)		FD CB byte F6	SET 6,(IY+indx)
FD CB byte C6	SET 0,(IY+indx)			
CB C7	SET 0,A			
CB C0	SET 0,B			

CB F7	SET 6,A
CB F0	SET 6,B
CB F1	SET 6,C
CB F2	SET 6,D
CB F3	SET 6,E
CB F4	SET 6,H
CB F5	SET 6,L

CB FE	SET 7,(HL)
DD CB byte FE	SET 7,(IX+indx)
FD CB byte FE	SET 7,(IY+indx)
CB FF	SET 7,A
CB F8	SET 7,B
CB F9	SET 7,C
CB FA	SET 7,D

#### SHIFT ARITHMETIC INSTRUCTIONS

CB 26	SLA (HL)
DD CB byte 26	SLA (IX+indx)
FD CB byte 26	SLA (IY+indx)

CB 27	SLA A
CB 28	SLA B
CB 29	SLA C
CB 2A	SLA D
CB 2B	SLA E
CB 2C	SLA H
CB 2D	SLA L

CB 2E	SRA (HL)
DD CB byte 2E	SRA (IX+indx)
FD CB byte 2E	SRA (IY+indx)

CB 2F	SRA A
CB 28	SRA B
CB 29	SRA C
CB 2A	SRA D
CB 2B	SRA E
CB 2C	SRA H
CB 2D	SRA L

#### SHIFT LOGICAL INSTRUCTIONS

CB 3E	SRL (HL)
DD CB byte 3E	SRL (IX+indx)
FD CB byte 3E	SRL (IY+indx)

CB 3F	SRL A
CB 38	SRL B
CB 39	SRL C
CB 3A	SRL D
CB 3B	SRL E
CB 3C	SRL H
CB 3D	SRL L

#### SUBTRACT INSTRUCTIONS

96	SUB (HL)
DD 96 byte	SUB (IX+indx)
FD 96 byte	SUB (IY+indx)

97	SUB A
90	SUB B
91	SUB C
92	SUB D
93	SUB E
94	SUB H
95	SUB L
D6 byte	SUB data

## **APPENDIX C**

### **TRS-80 ASCII Character Set: Decimal and Hexadecimal Codes**



Decimal	Hexadecimal	Character	Decimal	Hexadecimal	Character
32	20	space	80	50	P
33	21	!	81	51	Q
34	22	"	82	52	R
35	23	#	83	53	S
36	24	\$	84	54	T
37	25	%	85	55	U
38	26	&	86	56	V
39	27	'	87	57	W
40	28	(	88	58	X
41	29	)	89	59	Y
42	2A	*	90	5A	Z
43	2B	+	91	5B	↑
44	2C	,	92	5C	↓
45	2D	-	93	5D	←
46	2E	.	94	5E	→
47	2F	/	95	5F	—
48	30	0	96	60	@
49	31	1	97	61	
50	32	2	98	62	b
51	33	3	99	63	c
52	34	4	100	64	d
53	35	5	101	65	e
54	36	6	102	66	f
55	37	7	103	67	g
56	38	8	104	68	h
57	39	9	105	69	i
58	3A	:	106	6A	j
59	3B	;	107	6B	k
60	3C	<	108	6C	l
61	3D	=	109	6D	m
62	3E	>	110	6E	n
63	3F	?	111	6F	o
64	40	@	112	70	p
65	41	A	113	71	q
66	42	B	114	72	r
67	43	C	115	73	s
68	44	D	116	74	t
69	45	E	117	75	u
70	46	F	118	76	v
71	47	G	119	77	w
72	48	H	120	78	x
73	49	I	121	79	y
74	4A	J	122	7A	z
75	4B	K	123	7B	↑
76	4C	L	124	7C	↓
77	4D	M	125	7D	←
78	4E	N	126	7E	→
79	4F	O	127	7F	—

## **APPENDIX D**

### **TRS-80 Graphics Character Set: Decimal and Hexadecimal Codes**



0A0H  
160D



0A1H  
161D



0A2H  
162D



0A3H  
163D



0A4H  
164D



0A5H  
165D



0A6H  
166D



0A7H  
167D



0A8H  
168D



0A9H  
169D



0AAH  
170D



0ABH  
171D



0ACH  
172D



0ADH  
173D



0AEH  
174D



0AFH  
175D



0B0H  
176D



0B1H  
177D



0B2H  
178D



0B3H  
179D



0B4H  
180D



0B5H  
181D



0B6H  
182D



0B7H  
183D



0B8H  
184D



0B9H  
185D



0BAH  
186D



0BBH  
187D



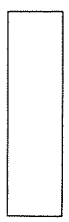
0BCH  
188D



0BDH  
189D



0BEH  
190D



0BFH  
191D



80H  
128D



81H  
129D



82H  
130D



83H  
131D



84H  
132D



85H  
133D



86H  
134D



87H  
135D



88H  
136D



89H  
137D



8AH  
138D



8BH  
139D



8CH  
140D



8DH  
141D



8EH  
142D



8FH  
143D



90H  
144D



91H  
145D



92H  
146D



93H  
147D



94H  
148D



95H  
149D



96H  
150D



97H  
151D



98H  
152D



99H  
153D



9AH  
154D



9BH  
155D



9CH  
156D



9DH  
157D



9EH  
158D



9FH  
159D

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